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D. Funaro

FORTRAN Routines for Spectral Methods

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Daniele Funaro

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Introduction

Spectral methods are an efficient tool to recover accurate approximated solutions to ordinary and partial differential equations, in terms of high-degree trigonometric or algebraic polynomials. A lot has been done to study the numerical implementation and the theoretical analysis of convergence of these techniques. A survey of the main results is given for instance in [1], [2], [3], [5], [7].

When building a numerical code using spectral methods, a preliminary part has to be devoted to a series of basic algorithms, not often readily available in the usual program libraries. This requires the user to be a little acquainted with the properties of orthogonal functions. The initialization is however quite standard and the collection of FORTRAN routines presented here is intended to provide the user with a ready software product, in order to allow a smooth start in the developement of a more extensive code.

The *double precision* subroutines are listed in this manual together with a detailed description. Style and notations are those adopted in [4], where *tau* and *collocation* approximations by algebraic polynomials are studied for one-dimensional differential equations. As a matter of fact, this text has been conceived to complement the material contained in [4].

The name of each routine consists of six letters. The third and the fourth denote the kind of polynomial basis considered: JA stands for Jacobi, LE for Legendre, CH for Chebyshev, LA for Laguerre, HE for Hermite. For the routines related with a certain set of collocation nodes, the last two letters denote the type of integration formula originating from these nodes: GA stands for Gauss, GL for Gauss-Lobatto, GR for Gauss-Radau.

In order to fit the mathematical notations, many vectors are dimensioned starting from the component zero. The user is invited to check the vector dimensioning of his main program with the help of the tables provided at the end of the manual.

Some of the routines related to the set of Chebyshev collocation nodes have a second version using the algorithm of the Fast Fourier Transform. In this case, external auxiliary subroutines of the NAG (*Numerical Algorithms Group*) FORTRAN Library are needed.

The routines have been tested when the various parameters fall in the range of standard applications. Of course, for high values of the polynomial degrees, round off errors may occur, and a special care is required in the computations involving Laguerre or Hermite polynomials. The user is suggested to double check the validity of the output results in those situations expected to be critical.

Purposes

GAMMAF: evaluates the real gamma function at a given point

VAJAPO: computes the value and the derivatives of the Jacobi polynomials at a given point

VALEPO: computes the value and the derivatives of the Legendre polynomials at a given point

VACHPO: computes the value and the derivatives of the Chebyshev polynomials at a given point

VALAPO: computes the value and the derivatives of the Laguerre polynomials at a given point

VAHEPO: computes the value and the derivatives of the Hermite polynomials at a given point

VALASF: computes the value and the derivative of the scaled Laguerre functions at a given point

VAHESF: computes the value and the derivative of the scaled Hermite functions at a given point

ZEJAGA: finds the zeroes of the Jacobi polynomials

ZELEGA: finds the zeroes of the Legendre polynomials

ZECHGA: finds the zeroes of the Chebyshev polynomials

ZELAGA: finds the zeroes of the Laguerre polynomials

ZEHEGA: finds the zeroes of the Hermite polynomials

ZEJAGL: finds the nodes of the Jacobi Gauss-Lobatto integration formula

ZELEGL: finds the nodes of the Legendre Gauss-Lobatto integration formula

ZECHGL: finds the nodes of the Chebyshev Gauss-Lobatto integration formula

- ZELAGR:** finds the nodes of the Laguerre Gauss-Radau integration formula
- WEJAGA:** finds the weights of the Jacobi Gauss integration formula
- WELEGA:** finds the weights of the Legendre Gauss integration formula
- WECHGA:** finds the weights of the Chebyshev Gauss integration formula
- WELAGA:** finds the weights of the Laguerre Gauss integration formula
- WEHEGA:** finds the weights of the Hermite Gauss integration formula
- WEJAGL:** finds the weights of the Jacobi Gauss-Lobatto integration formula
- WELEGL:** finds the weights of the Legendre Gauss-Lobatto integration formula
- WECHGL:** finds the weights of the Chebyshev Gauss-Lobatto integration formula
- WELAGR:** finds the weights of the Laguerre Gauss-Radau integration formula
- WECHCC:** finds the weights of the Clenshaw-Curtis integration formula
- INJAGA:** evaluates at a given point the value of a polynomial given at the Jacobi zeroes
- INLEGA:** evaluates at a given point the value of a polynomial given at the Legendre zeroes
- INCHGA:** evaluates at a given point the value of a polynomial given at the Chebyshev zeroes
- INLAGA:** evaluates at a given point the value of a polynomial given at the Laguerre zeroes
- INHEGA:** evaluates at a given point the value of a polynomial given at the Hermite zeroes
- INJAGL:** evaluates at a given point the value of a polynomial given at the Jacobi Gauss-Lobatto nodes
- INLEGL:** evaluates at a given point the value of a polynomial given at the Legendre Gauss-Lobatto nodes
- INCHGL:** evaluates at a given point the value of a polynomial given at the Chebyshev Gauss-Lobatto nodes
- INLAGR:** evaluates at a given point the value of a polynomial given at the Laguerre Gauss-Radau nodes
- NOLEGA:** evaluates the $L^2(-1,1)$ norm and the discrete maximum norm of a polynomial given at the Legendre zeroes

NOCHGA: evaluates the $L^2(-1,1)$ norm (with and without weight function) and the discrete maximum norm of a polynomial given at the Chebyshev zeroes

NOJAGL: evaluates the weighted $L^2(-1,1)$ norm, the quadrature norm, and the discrete maximum norm of a polynomial given at the Jacobi Gauss-Lobatto nodes

NOLEGL: evaluates the $L^2(-1,1)$ norm, the quadrature norm, and the discrete maximum norm of a polynomial given at the Legendre Gauss-Lobatto nodes

NOCHGL: evaluates the $L^2(-1,1)$ norm (with and without weight function), the quadrature norm, and the discrete maximum norm of a polynomial given at the Chebyshev Gauss-Lobatto nodes

COJAGA: evaluates the Jacobi Fourier coefficients of a polynomial given at the Jacobi zeroes

COLEGA: evaluates the Legendre Fourier coefficients of a polynomial given at the Legendre zeroes

COCHGA: evaluates the Chebyshev Fourier coefficients of a polynomial given at the Chebyshev zeroes (without using FFT)

COLAGA: evaluates the Laguerre Fourier coefficients of a polynomial given at the Laguerre zeroes

COHEGA: evaluates the Hermite Fourier coefficients of a polynomial given at the Hermite zeroes

COJAGL: evaluates the Jacobi Fourier coefficients of a polynomial given at the Jacobi Gauss-Lobatto nodes

COLEGL: evaluates the Legendre Fourier coefficients of a polynomial given at the Legendre Gauss-Lobatto nodes

COCHGL: evaluates the Chebyshev Fourier coefficients of a polynomial given at the Chebyshev Gauss-Lobatto nodes (without using FFT)

COLAGR: evaluates the Laguerre Fourier coefficients of a polynomial given at the Laguerre Gauss-Radau nodes

PVJAEX: computes the value and the derivatives at a given point of a polynomial, from its Jacobi Fourier coefficients

PVLEEX: computes the value and the derivatives at a given point of a polynomial, from its Legendre Fourier coefficients

PVCHEX: computes the value and the derivatives at a given point of a polynomial, from its Chebyshev Fourier coefficients

- PVLAEX:** computes the value and the derivatives at a given point of a polynomial, from its Laguerre Fourier coefficients
- PVHEEX:** computes the value and the derivatives at a given point of a polynomial, from its Hermite Fourier coefficients
- NOJAEX:** evaluates the weighted $L^2(-1, 1)$ norm of a polynomial, from its Jacobi Fourier coefficients
- NOLEEX:** evaluates the $L^2(-1, 1)$ norm of a polynomial, from its Legendre Fourier coefficients
- NOCHEX:** evaluates the $L^2(-1, 1)$ norm (with and without weight function) of a polynomial, from its Chebyshev Fourier coefficients
- NOLAEX:** evaluates the weighted $L^2(0, +\infty)$ norm of a polynomial, from its Laguerre Fourier coefficients
- NOHEEX:** evaluates the weighted $L^2(\mathbf{R})$ norm of a polynomial, from its Hermite Fourier coefficients
- COJADE:** computes the Jacobi Fourier coefficients of the derivatives of a polynomial, from its Jacobi Fourier coefficients
- COLEDE:** computes the Legendre Fourier coefficients of the derivatives of a polynomial, from its Legendre Fourier coefficients
- COCHDE:** computes the Chebyshev Fourier coefficients of the derivatives of a polynomial, from its Chebyshev Fourier coefficients
- COLADE:** computes the Laguerre Fourier coefficients of the derivatives of a polynomial, from its Laguerre Fourier coefficients
- COHEDE:** computes the Hermite Fourier coefficients of the derivatives of a polynomial, from its Hermite Fourier coefficients
- DEJAGA:** computes the derivative at the Jacobi zeroes of a polynomial given at the Jacobi zeroes
- DELAGA:** computes the derivative at the Laguerre zeroes of a polynomial given at the Laguerre zeroes
- DEHEGA:** computes the derivative at the Hermite zeroes of a polynomial given at the Hermite zeroes
- DEJAGL:** computes the derivative at the Jacobi Gauss-Lobatto nodes of a polynomial given at the Jacobi Gauss-Lobatto nodes
- DELEGL:** computes the derivative at the Legendre Gauss-Lobatto nodes of a polynomial given at the Legendre Gauss-Lobatto nodes

- DECHGL:** computes the derivative at the Chebyshev Gauss-Lobatto nodes of a polynomial given at the Chebyshev Gauss-Lobatto nodes (without using FFT)
- DELAGR:** computes the derivative at the Laguerre Gauss-Radau nodes of a polynomial given at the Laguerre Gauss-Radau nodes
- DMJAGL:** gives the entries of the derivative matrix relative to the Jacobi Gauss-Lobatto nodes
- DMLEGL:** gives the entries of the derivative matrix relative to the Legendre Gauss-Lobatto nodes
- DMCHGL:** gives the entries of the derivative matrix relative to the Chebyshev Gauss-Lobatto nodes
- DMLAGR:** gives the entries of the derivative matrix relative to the Laguerre Gauss-Radau nodes
- FCCHGA:** evaluates the Chebyshev Fourier coefficients of a polynomial given at the Chebyshev zeroes (using FFT).
- FCCHGL:** evaluates the Chebyshev Fourier coefficients of a polynomial given at the Chebyshev Gauss-Lobatto nodes (using FFT).
- FVCHGL:** computes the values at the Chebyshev Gauss-Lobatto nodes of a polynomial, from its Chebyshev Fourier coefficients (using FFT).
- FDCHGL:** computes the derivative at the Chebyshev Gauss-Lobatto nodes of a polynomial given at the Chebyshev Gauss-Lobatto nodes (using FFT).

**DESCRIPTIONS
AND
LISTINGS**

GAMMAF

The routine computes the value, at a given point $x \in]10^{-75}, 57[$, of the Gamma function

$$\Gamma(x) = \int_0^{+\infty} t^{x-1} e^{-t} dt.$$

One is reduced to work with $x < 1$ by using the relations

$$\Gamma(x+1) = x\Gamma(x), \quad x > 0 \quad \text{and} \quad \Gamma(n+1) = n!, \quad n \in \mathbf{N}.$$

The computation is further restricted to the interval $]0.5, 1[$ by virtue of the formula

$$\Gamma(x)\Gamma(1-x) = \frac{\pi}{\sin \pi x}, \quad 0 < x < 1.$$

Finally, the value of $\Gamma(x)$, $x \in]0.5, 1[$, is obtained by a truncation of a series expansion as suggested in [6], Vol.1, p.30.

<i>Input variable</i>	<i>Output variable</i>
X , the argument x	GX , the value of Γ in x

```

SUBROUTINE GAMMAF(X,GX)
*****
*   COMPUTES THE GAMMA FUNCTION AT A GIVEN POINT
*   X = ARGUMENT GREATER THAN 1.E-75 AND SMALLER THAN 57.
*   GX= VALUE OF GAMMA IN X
*****
  IMPLICIT DOUBLE PRECISION (A-H,O-Z)
  DIMENSION C(11)

  IF (X .LE. 1.D-75 .OR. X .GE. 57.DO) THEN
    WRITE(*,*) 'ARGUMENT OUT OF RANGE IN SUBROUTINE GAMMAF'
    RETURN
  ENDIF

```

```
      PI = 3.14159265358979323846D0
      EPS = 1.D-14
      XX = X
      GX = 1.0D0

1     IF (DABS(XX-1.D0) .LT. EPS) RETURN
      IF (XX .GE. 1.D0) THEN
        XX = XX-1.D0
        GX = GX*XX
        GOTO 1
      ENDIF
      IND = 0
      IF (XX .LT. .5D0) THEN
        IND = 1
        GX = GX*PI/DSIN(PI*XX)
        XX = 1.D0-XX
      ENDIF

      PR = 1.D0
      S = 0.426401432711220868D0
      C(1) = -0.524741987629368444D0
      C(2) = 0.116154405493589130D0
      C(3) = -0.765978624506602380D-2
      C(4) = 0.899719449391378898D-4
      C(5) = -0.194536980009534621D-7
      C(6) = 0.199382839513630987D-10
      C(7) = -0.204209590209541319D-11
      C(8) = 0.863896817907000175D-13
      C(9) = 0.152237501608472336D-13
      C(10) = -0.82572517527771995D-14
      C(11) = 0.29973478220522461D-14

      DO 2 K=1,11
        PR = PR*(XX-DFLOAT(K))/(XX+DFLOAT(K-1))
        S = S+C(K)*PR
2     CONTINUE
      G = S*DEXP(1.D0-XX)*(XX+4.5D0)**(XX-.5D0)
      IF (IND .EQ. 1) THEN
        GX = GX/G
      ELSE
        GX = GX*G
      ENDIF

      RETURN
      END
```

VAJAPO

The routine computes the value of the Jacobi polynomial $P_n^{(\alpha,\beta)}$ of degree $n \in \mathbf{N}$ ($\alpha > -1$ and $\beta > -1$ are real parameters) at the point x , according to the recursion formula

$$\begin{cases} P_0^{(\alpha,\beta)}(x) = 1, \\ P_1^{(\alpha,\beta)}(x) = \frac{1}{2}(\alpha + \beta + 2)x + \frac{1}{2}(\alpha - \beta), \\ P_n^{(\alpha,\beta)}(x) = c_1^{-1}[(c_2x + c_3)P_{n-1}^{(\alpha,\beta)}(x) - c_4P_{n-2}^{(\alpha,\beta)}(x)], \quad n \geq 2, \end{cases}$$

where

$$c_1 = 2n(n + \alpha + \beta)(2n + \alpha + \beta - 2),$$

$$c_2 = (2n + \alpha + \beta - 1)(2n + \alpha + \beta - 2)(2n + \alpha + \beta),$$

$$c_3 = (2n + \alpha + \beta - 1)(\alpha^2 - \beta^2),$$

$$c_4 = 2(n + \alpha - 1)(2n + \alpha + \beta)(n + \beta - 1).$$

Simultaneously, first and second derivatives of $P_n^{(\alpha,\beta)}$ at the point x are computed by taking the derivatives of the terms of the above recursion formula.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of $P_n^{(\alpha,\beta)}$ in x
A , the parameter α	DY , the value of $\frac{d}{dx}P_n^{(\alpha,\beta)}$ in x
B , the parameter β	$D2Y$, the value of $\frac{d^2}{dx^2}P_n^{(\alpha,\beta)}$ in x
X , the argument x	

```

      SUBROUTINE VAJAP0(N,A,B,X,Y,DY,D2Y)
*****
*   COMPUTES THE VALUE OF THE JACOBI POLYNOMIAL OF DEGREE N
*   AND ITS FIRST AND SECOND DERIVATIVES AT A GIVEN POINT
*   N   = DEGREE OF THE POLYNOMIAL
*   A   = PARAMETER >-1
*   B   = PARAMETER >-1
*   X   = POINT IN WHICH THE COMPUTATION IS PERFORMED
*   Y   = VALUE OF THE POLYNOMIAL IN X
*   DY  = VALUE OF THE FIRST DERIVATIVE IN X
*   D2Y = VALUE OF THE SECOND DERIVATIVE IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      Y   = 1.DO
      DY  = 0.DO
      D2Y = 0.DO
      IF (N .EQ. 0) RETURN

      AB  = A+B
      Y   = .5D0*(AB+2.DO)*X+.5D0*(A-B)
      DY  = .5D0*(AB+2.DO)
      D2Y = 0.DO
      IF(N .EQ. 1) RETURN

      YP  = 1.DO
      DYP = 0.DO
      D2YP = 0.DO
      DO 1 I=2,N
      DI  = DFLOAT(I)
      CO  = 2.DO*DI+AB
      C1  = 2.DO*DI*(DI+AB)*(CO-2.DO)
      C2  = (CO-1.DO)*(CO-2.DO)*CO
      C3  = (CO-1.DO)*(A-B)*AB
      C4  = 2.DO*(DI+A-1.DO)*CO*(DI+B-1.DO)
      YM  = Y
      Y   = ((C2*X+C3)*Y-C4*YP)/C1
      YP  = YM
      DYM = DY
      DY  = ((C2*X+C3)*DY-C4*DYP+C2*YP)/C1
      DYP = DYM
      D2YM = D2Y
      D2Y = ((C2*X+C3)*D2Y-C4*D2YP+2.DO*C2*DYP)/C1
      D2YP = D2YM
1      CONTINUE

      RETURN
      END

```

VALEPO

The routine computes the value of the Legendre polynomial P_n of degree $n \in \mathbf{N}$ and its first and second derivatives at the point x , according to the recursion formulas

$$\begin{cases} P_0(x) = 1, \\ P_1(x) = x, \\ P_n(x) = \frac{1}{n}[(2n-1)xP_{n-1}(x) - (n-1)P_{n-2}(x)] \quad n \geq 2, \end{cases}$$

$$\begin{cases} P'_0(x) = 0, \\ P'_1(x) = 1, \\ P'_n(x) = \frac{1}{n}[(2n-1)(xP'_{n-1}(x) + P_{n-1}(x)) - (n-1)P'_{n-2}(x)] \quad n \geq 2, \end{cases}$$

$$\begin{cases} P''_0(x) = 0, \\ P''_1(x) = 0, \\ P''_n(x) = \frac{1}{n}[(2n-1)(xP''_{n-1}(x) + 2P'_{n-1}(x)) - (n-1)P''_{n-2}(x)] \quad n \geq 2. \end{cases}$$

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of P_n in x
X , the argument x	DY , the value of P'_n in x
	$D2Y$, the value of P''_n in x

```

      SUBROUTINE VALEPO(N,X,Y,DY,D2Y)
*****
*   COMPUTES THE VALUE OF THE LEGENDRE POLYNOMIAL OF DEGREE N
*   AND ITS FIRST AND SECOND DERIVATIVES AT A GIVEN POINT
*   N = DEGREE OF THE POLYNOMIAL
*   X = POINT IN WHICH THE COMPUTATION IS PERFORMED
*   Y = VALUE OF THE POLYNOMIAL IN X
*   DY = VALUE OF THE FIRST DERIVATIVE IN X
*   D2Y= VALUE OF THE SECOND DERIVATIVE IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)

      Y   = 1.DO
      DY  = 0.DO
      D2Y = 0.DO
      IF (N .EQ. 0) RETURN

      Y   = X
      DY  = 1.DO
      D2Y = 0.DO
      IF(N .EQ. 1) RETURN

      YP  = 1.DO
      DYP = 0.DO
      D2YP = 0.DO
      DO 1 I=2,N
      C1 = DFLOAT(I)
      C2 = 2.DO*C1-1.DO
      C4 = C1-1.DO
      YM = Y
      Y   = (C2*X*Y-C4*YP)/C1
      YP = YM
      DYM = DY
      DY  = (C2*X*DY-C4*DYP+C2*YP)/C1
      DYP = DYM
      D2YM = D2Y
      D2Y = (C2*X*D2Y-C4*D2YP+2.DO*C2*DYP)/C1
      D2YP = D2YM
1      CONTINUE

      RETURN
      END

```

VACHPO

The routine computes the value of the Chebyshev polynomial T_n of degree $n \in \mathbf{N}$ and its first and second derivatives at the point x , according to the recursion formulas

$$\begin{cases} T_0(x) = 1, \\ T_1(x) = x, \\ T_n(x) = 2xT_{n-1}(x) - T_{n-2}(x) \quad n \geq 2, \end{cases}$$

$$\begin{cases} T'_0(x) = 0, \\ T'_1(x) = 1, \\ T'_n(x) = 2xT'_{n-1}(x) + 2T_{n-1}(x) - T'_{n-2}(x) \quad n \geq 2, \end{cases}$$

$$\begin{cases} T''_0(x) = 0, \\ T''_1(x) = 0, \\ T''_n(x) = 2xT''_{n-1}(x) + 4T'_{n-1}(x) - T''_{n-2}(x) \quad n \geq 2. \end{cases}$$

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of T_n in x
X , the argument x	DY , the value of T'_n in x
	$D2Y$, the value of T''_n in x


```
      SUBROUTINE VACHPO(N,X,Y,DY,D2Y)
*****
*   COMPUTES THE VALUE OF THE CHEBYSHEV POLYNOMIAL OF DEGREE N
*   AND ITS FIRST AND SECOND DERIVATIVES AT A GIVEN POINT
*   N = DEGREE OF THE POLYNOMIAL
*   X = POINT IN WHICH THE COMPUTATION IS PERFORMED
*   Y = VALUE OF THE POLYNOMIAL IN X
*   DY = VALUE OF THE FIRST DERIVATIVE IN X
*   D2Y= VALUE OF THE SECOND DERIVATIVE IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)

      Y   = 1.DO
      DY  = 0.DO
      D2Y = 0.DO
      IF (N .EQ. 0) RETURN

      Y   = X
      DY  = 1.DO
      D2Y = 0.DO
      IF (N .EQ. 1) RETURN

      YP  = 1.DO
      DYP = 0.DO
      D2YP = 0.DO
      DO 1 K=2,N
      YM  = Y
      Y   = 2.DO*X*Y-YP
      YP  = YM
      DYM = DY
      DY  = 2.DO*X*DY+2.DO*YP-DYP
      DYP = DYM
      D2YM= D2Y
      D2Y = 2.DO*X*D2Y+4.DO*DYP-D2YP
      D2YP= D2YM
1      CONTINUE

      RETURN
      END
```

VALAPO

The routine computes the value of the Laguerre polynomial $L_n^{(\alpha)}$ of degree $n \in \mathbf{N}$ ($\alpha > -1$ is a real parameter) at the point x , according to the recursion formula

$$\begin{cases} L_0^{(\alpha)}(x) = 1, \\ L_1^{(\alpha)}(x) = 1 + \alpha - x, \\ L_n^{(\alpha)}(x) = \frac{1}{n}[(2n + \alpha - 1 - x)L_{n-1}^{(\alpha)}(x) - (n + \alpha - 1)L_{n-2}^{(\alpha)}(x)], \quad n \geq 2. \end{cases}$$

Simultaneously, first and second derivatives of $L_n^{(\alpha)}$ at the point x are computed by taking the derivatives of the terms of the above recursion formula.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of $L_n^{(\alpha)}$ in x
A , the parameter α	DY , the value of $\frac{d}{dx}L_n^{(\alpha)}$ in x
X , the argument x	$D2Y$, the value of $\frac{d^2}{dx^2}L_n^{(\alpha)}$ in x

```

      SUBROUTINE VALAPO(N,A,X,Y,DY,D2Y)
      *****
      *   COMPUTES THE VALUE OF THE LAGUERRE POLYNOMIAL OF DEGREE N
      *   AND ITS FIRST AND SECOND DERIVATIVES AT A GIVEN POINT
      *   N   = DEGREE OF THE POLYNOMIAL
      *   A   = PARAMETER >-1
      *   X   = POINT IN WHICH THE COMPUTATION IS PERFORMED
      *   Y   = VALUE OF THE POLYNOMIAL IN X
      *   DY  = VALUE OF THE FIRST DERIVATIVE IN X
      *   D2Y = VALUE OF THE SECOND DERIVATIVE IN X
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)

      Y   = 1.DO
      DY  = 0.DO
      D2Y = 0.DO
      IF (N .EQ. 0) RETURN

      Y   = 1.DO+A-X
      DY  = -1.DO
      D2Y = 0.DO
      IF (N .EQ. 1) RETURN

      YP  = 1.DO
      DYP = 0.DO
      D2YP= 0.DO
      DO 1 K=2,N
      DK = DFLOAT(K)
      B1 = (2.DO*DK+A-1.DO-X)/DK
      B2 = (DK+A-1.DO)/DK
      YM = Y
      Y  = B1*Y-B2*YP
      YP = YM
      DYM = DY
      DY  = B1*DY-YP/DK-B2*DYP
      DYP = DYM
      D2YM= D2Y
      D2Y = B1*D2Y-2.DO*DYP/DK-B2*D2YP
      D2YP= D2YM
1     CONTINUE

      RETURN
      END

```

VAHEPO

The routine computes the value of the Hermite polynomial H_n , $n \in \mathbf{N}$, and its first and second derivatives at the point x , according to the recursion formulas

$$\begin{cases} H_0(x) = 1, \\ H_1(x) = 2x, \\ H_n(x) = 2xH_{n-1}(x) - 2(n-1)H_{n-2}(x) & n \geq 2, \end{cases}$$

$$\begin{cases} H'_0(x) = 0, \\ H'_n(x) = 2nH_{n-1}(x) & n \geq 1, \end{cases}$$

$$\begin{cases} H''_0(x) = 0, \\ H''_1(x) = 0, \\ H''_n(x) = 4n(n-1)H_{n-2}(x) & n \geq 2. \end{cases}$$

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of H_n in x
X , the argument x	DY , the value of H'_n in x
	$D2Y$, the value of H''_n in x

```
      SUBROUTINE VAHEPO(N,X,Y,DY,D2Y)
*****
*   COMPUTES THE VALUE OF THE HERMITE POLYNOMIAL OF DEGREE N
*   AND ITS FIRST AND SECOND DERIVATIVES AT A GIVEN POINT
*   N = DEGREE OF THE POLYNOMIAL
*   X = POINT IN WHICH THE COMPUTATION IS PERFORMED
*   Y = VALUE OF THE POLYNOMIAL IN X
*   DY = VALUE OF THE FIRST DERIVATIVE IN X
*   D2Y= VALUE OF THE SECOND DERIVATIVE IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)

      Y   = 1.DO
      DY  = 0.DO
      D2Y = 0.DO
      IF (N .EQ. 0) RETURN

      Y   = 2.DO*X
      DY  = 2.DO
      D2Y = 0.DO
      IF (N .EQ. 1) RETURN

      YP=1.DO
      DO 1 K=2,N
        DK = DFLOAT(K-1)
        YM = Y
        Y  = 2.DO*X*Y-2.DO*DK*YP
        YPM = YP
        YP = YM
1      CONTINUE
      DN = 2.DO*DFLOAT(N)
      DNN = 2.DO*DFLOAT(N-1)
      DY  = DN*YP
      D2Y = DN*DNN*YPM

      RETURN
      END
```

VALASF

The routine computes the value at the point $x \geq 0$ of the *scaled Laguerre function* of degree $n \in \mathbf{N}$, defined for $\alpha > -1$ by

$$\begin{cases} \hat{L}_0^{(\alpha)}(x) &= L_0^{(\alpha)}(x), \\ \hat{L}_n^{(\alpha)}(x) &= \left[\frac{\Gamma(n + \alpha + 1)}{n! \Gamma(\alpha + 1)} \prod_{k=1}^n \left(1 + \frac{x}{4k} \right) \right]^{-1} L_n^{(\alpha)}(x) \quad n \geq 1. \end{cases}$$

The following recursion formula is used:

$$\begin{cases} \hat{L}_0^{(\alpha)}(x) &= 1, \\ \hat{L}_1^{(\alpha)}(x) &= \frac{4(\alpha + 1 - x)}{(\alpha + 1)(x + 4)}, \\ \hat{L}_n^{(\alpha)}(x) &= \frac{4n}{(n + \alpha)(4n + x)} \left[(2n + \alpha - 1 - x)\hat{L}_{n-1}^{(\alpha)}(x) - \frac{4(n-1)^2}{4n + x - 4}\hat{L}_{n-2}^{(\alpha)}(x) \right]. \end{cases}$$

Simultaneously, the first derivative of $\hat{L}_n^{(\alpha)}$ at the point x is computed by taking the derivative of the terms of the above recursion formula.

Scaled Laguerre functions will be used in place of Laguerre polynomials when possible, in order to improve the implementation (see the introduction).

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of $\hat{L}_n^{(\alpha)}$ in x
A , the parameter α	DY , the value of $\frac{d}{dx}\hat{L}_n^{(\alpha)}$ in x
X , the argument x	

```

      SUBROUTINE VALASF(N,A,X,Y,DY)
*****
*   COMPUTES THE VALUES OF THE SCALED LAGUERRE FUNCTION OF DEGREE N
*   AND ITS FIRST DERIVATIVE AT A GIVEN POINT
*   N   = DEGREE
*   A   = PARAMETER >-1
*   X   = POINT (NON NEGATIVE) IN WHICH THE COMPUTATION IS PERFORMED
*   Y   = VALUE OF THE FUNCTION IN X
*   DY  = VALUE OF THE FIRST DERIVATIVE IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)

      Y = 1.DO
      DY = 0.DO
      IF (N .EQ. 0) RETURN

      CO = 1.DO/(4.DO+X)
      C1 = 4.DO*CO/(A+1.DO)
      Y = C1*(A+1.DO-X)
      DY = -CO*C1*(A+5.DO)
      IF (N .EQ. 1) RETURN

      YP = 1.DO
      DYP = 0.DO
      DO 1 K=2,N
      DK = DFLOAT(K)
      DK4 = 4.DO*DK
      CO = 1.DO/(DK4+X)
      C1 = DK4+X-2.DO
      C2 = 1.DO/(C1-2.DO)
      C3 = DK4*CO/(DK+A)
      C4 = 2.DO*DK+A-1.DO
      C5 = C4-X
      C6 = C4+DK4
      C7 = C2*DFLOAT(4*(K-1)**2)
      DYM = DY
      DY = C3*(C5*DY-CO*C6*Y+C7*(2.DO*CO*C1*C2*YP-DYP))
      DYP = DYM
      YM = Y
      Y = C3*(C5*Y-C7*YP)
      YP = YM
1     CONTINUE

      RETURN
      END

```

VAHESF

The routine computes the value and the first derivative at the point x of the *scaled Hermite function* of degree $n \in \mathbf{N}$ defined by

$$\hat{H}_n(x) = \hat{L}_{n/2}^{(-1/2)}(x^2) \quad \text{if } n \text{ is even,}$$

$$\hat{H}_n(x) = x \hat{L}_{(n-1)/2}^{(1/2)}(x^2) \quad \text{if } n \text{ is odd.}$$

Here, $\hat{L}_k^{(\alpha)}$, $k \in \mathbf{N}$, $\alpha > -1$, are the scaled Laguerre functions, defined in the description of subroutine VALASF.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of \hat{H}_n in x
X , the argument x	DY , the value of \hat{H}'_n in x

Auxiliary routine: **VALASF**


```
      SUBROUTINE VAHESF(N,X,Y,DY)
*****
*   COMPUTES THE VALUES OF THE SCALED HERMITE FUNCTION
*   OF DEGREE N AND ITS FIRST DERIVATIVE AT A GIVEN POINT
*   N   = DEGREE
*   X   = POINT IN WHICH THE COMPUTATION IS PERFORMED
*   Y   = VALUE OF THE FUNCTION IN X
*   DY  = VALUE OF THE FIRST DERIVATIVE IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)

      Y = 1.DO
      DY = 0.DO
      IF (N .EQ. 0) RETURN

      XX = X*X
      M = N/2
      IF(N .EQ. 2*M) THEN
        A = -.5DO
        CALL VALASF(M,A,XX,Y,DY)
        DY = 2.DO*X*DY
      ELSE
        A = .5DO
        CALL VALASF(M,A,XX,Y,DY)
        DY = Y+2.DO*XX*DY
        Y = X*Y
      ENDIF

      RETURN
      END
```

ZEJAGA

For any $n \geq 1$, the routine computes the n zeroes $\xi_i^{(n)} \in]-1, 1[$, $1 \leq i \leq n$, of the Jacobi polynomial $P_n^{(\alpha, \beta)}$, where the parameters satisfy $-\frac{1}{2} \leq \alpha \leq \frac{1}{2}$, $-\frac{1}{2} \leq \beta \leq \frac{1}{2}$.

The Newton method is used with the initial guess

$$\xi_i^{(n)} \approx -\cos\left(\frac{4i + \alpha + \beta - 1}{2n + \alpha + \beta + 1}\right) \frac{\pi}{2} \quad 1 \leq i \leq n.$$

The values of the Jacobi polynomial and its derivative at a given point are obtained by the recursion formula given in the description of subroutine VAJAPO.

An output vector contains the zeroes in increasing order.

Another output vector contains the quantities $\frac{d}{dx} P_n^{(\alpha, \beta)}(\xi_i^{(n)})$, $1 \leq i \leq n$.

The zeroes of the Legendre polynomial P_n are obtained either by setting $\alpha = \beta = 0$ in this routine, or by calling subroutine ZELEGA. Similarly, the zeroes of the Chebyshev polynomial T_n are obtained either by setting $\alpha = \beta = -\frac{1}{2}$ in this routine, or by calling subroutine ZECHGA.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	CS , vector containing the zeroes $\xi_i^{(n)}$
A , the parameter α	DZ , vector of the values $\frac{d}{dx} P_n^{(\alpha, \beta)}(\xi_i^{(n)})$
B , the parameter β	

Auxiliary routine: **VAJAPO**

```

      SUBROUTINE ZEJAGA(N,A,B,CS,DZ)
      *****
      *   COMPUTES THE ZEROES OF THE JACOBI POLYNOMIAL OF DEGREE N
      *   N   = THE NUMBER OF ZEROES
      *   A   = PARAMETER BETWEEN -1/2 AND 1/2
      *   B   = PARAMETER BETWEEN -1/2 AND 1/2
      *   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
      *   DZ  = VECTOR OF THE DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1)
      IF (N .EQ. 0) RETURN

      AB = A+B
      CS(1) = (B-A)/(AB+2.DO)
      DZ(1) = .5DO*AB+1.DO
      IF(N .EQ. 1) RETURN

      EPS= 1.E-17
      PH = 1.57079632679489661923DO
      C  = PH/(2.DO*DFLOAT(N)+AB+1.DO)
      DO 1 I=1,N
        DI = DFLOAT(I)
        CSX = -DCOS(C*(4.DO*DI+AB-1.DO))
      DO 2 IT=1,8
        CALL VAJAP0(N,A,B,CSX,Y,DY,D2Y)
        IF(DABS(Y) .LT. EPS) GOTO 3
        CSX = CSX-Y/DY
2      CONTINUE
3      IF(DABS(CSX) .LT. EPS) CSX=0.DO
        CS(I) = CSX
        DZ(I) = DY
1      CONTINUE

      RETURN
      END

```

ZELEGA

For any $n \geq 1$, the routine computes the n zeroes $\xi_i^{(n)} \in]-1, 1[$, $1 \leq i \leq n$, of the Legendre polynomial P_n .

The Newton method is used with the initial guess

$$\xi_i^{(n)} \approx -\cos\left(\frac{4i-1}{2n+1}\right)\frac{\pi}{2} \quad 1 \leq i \leq n.$$

The values of the Legendre polynomial and its derivative at a given point are obtained by the recursion formula given in the description of subroutine VALEPO.

An output vector contains the zeroes in increasing order.

Another output vector contains the quantities $P'_n(\xi_i^{(n)})$, $1 \leq i \leq n$.

<i>Input variable</i>	<i>Output variables</i>
N , the degree n	CS , vector containing the zeroes $\xi_i^{(n)}$ DZ , vector of the values $P'_n(\xi_i^{(n)})$

Auxiliary routine: **VALEPO**

```

      SUBROUTINE ZELEGA(N,CS,DZ)
      *****
      *   COMPUTES THE ZEROES OF THE LEGENDRE POLYNOMIAL OF DEGREE N
      *   N   = THE NUMBER OF ZEROES
      *   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
      *   DZ  = VECTOR OF THE DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1)
      IF (N .EQ. .0) RETURN

      CS(1) = 0.DO
      DZ(1) = 1.DO
      IF(N .EQ. 1) RETURN

      N2 = N/2
      IN = 2*N-4*N2-1
      PH = 1.57079632679489661923D0
      C  = PH/(2.DO*DFLOAT(N)+1.DO)
      DO 1 I=1,N2
      DI  = DFLOAT(I)
      CSX = DCOS(C*(4.DO*DI-1.DO))
      DO 2 IT=1,8
      CALL VALEPO(N,CSX,Y,DY,D2Y)
      CSX = CSX-Y/DY
2     CONTINUE
      CS(I) = -CSX
      CS(N-I+1) = CSX
      DZ(I) = DY*DFLOAT(IN)
      DZ(N-I+1) = DY
1     CONTINUE

      IF(IN .EQ. -1) RETURN
      CSX = 0.DO
      CS(N2+1) = CSX
      CALL VALEPO(N,CSX,Y,DY,D2Y)
      DZ(N2+1) = DY

      RETURN
      END

```

ZECHGA

For any $n \geq 1$, the routine gives the n zeroes

$$\xi_i^{(n)} = -\cos \frac{(2i-1)\pi}{2n} \in]-1, 1[\quad 1 \leq i \leq n,$$

of the Chebyshev polynomial T_n .

An output vector contains the zeroes in increasing order.

Another output vector contains the quantities

$$T'_n(\xi_i^{(n)}) = \frac{n(-1)^{n+i}}{\sqrt{1 - [\xi_i^{(n)}]^2}} \quad 1 \leq i \leq n.$$

<i>Input variable</i>	<i>Output variables</i>
N , the degree n	CS , vector containing the zeroes $\xi_i^{(n)}$ DZ , vector of the values $T'_n(\xi_i^{(n)})$

```

SUBROUTINE ZECHGA(N,CS,DZ)
*****
*   COMPUTES THE ZEROES OF THE CHEBYSHEV POLYNOMIAL OF DEGREE N
*   N   = THE NUMBER OF ZEROES
*   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
*   DZ  = VECTOR OF THE DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1)
      IF (N .EQ. 0) RETURN

      CS(1) = 0.DO
      DZ(1) = 1.DO
      IF(N .EQ. 1) RETURN

      N2 = N/2
      IN = 1+4*N2-2*N
      PH = 1.57079632679489661923D0
      DN = DFLOAT(N)
      C  = PH/DN
      SI = -1.DO
DO 1 I=1,N2
      DI = DFLOAT(I)
      CSX = DCOS(C*(2.DO*DI-1.DO))
      CS(I) = -CSX
      CS(N-I+1) = CSX
      QX = DN/DSQRT(1.DO-CSX*CSX)
      DZ(I) = QX*SI*DFLOAT(IN)
      DZ(N-I+1) = -QX*SI
      SI = -SI
1  CONTINUE

      IF(IN .EQ. 1) RETURN
      CS(N2+1) = 0.DO
      N4 = N2/2
      IN2 = 1+4*N4-2*N2
      DZ(N2+1) = DN*DFLOAT(IN2)

      RETURN
      END

```

ZELAGA

For any $n \geq 1$, the routine computes the n zeroes $\xi_i^{(n)} \in]0, +\infty[$, $1 \leq i \leq n$, of the Laguerre polynomial $L_n^{(\alpha)}$, $\alpha > -1$. These coincide with the zeroes of the scaled Laguerre function $\hat{L}_n^{(\alpha)}$, $\alpha > -1$, introduced in the description of subroutine VALASF. A better treatment of the rounding errors is realized by using scaled Laguerre functions in the computation.

The Newton method is used with the initial guess

$$\xi_i^{(n)} \approx 2(2n + \alpha + 1)z_i - \frac{1}{6(2n + \alpha + 1)} \left[\frac{5}{4(z_i - 1)^2} + \frac{1}{z_i - 1} - 1 + 3\alpha^2 \right],$$

where

$$z_i = \left[\cos \left(\frac{y_i}{2} \right) \right]^2 \quad 1 \leq i \leq n,$$

and

$$y_i - \sin y_i = 2\pi \frac{n - i + 3/4}{2n + \alpha + 1} \quad 1 \leq i \leq n.$$

The values of the scaled Laguerre function $\hat{L}_n^{(\alpha)}$ and its derivative at a given point are obtained by the recursion formula given in the description of subroutine VALASF.

An output vector contains the zeroes in increasing order.

Another output vector contains the quantities $\frac{d}{dx} \hat{L}_n^{(\alpha)}(\xi_i^{(n)})$, $1 \leq i \leq n$.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	CS , vector containing the zeroes $\xi_i^{(n)}$
A , the parameter α	DZ , vector of the values $\frac{d}{dx} \hat{L}_n^{(\alpha)}(\xi_i^{(n)})$

Auxiliary routine: **VALASF**


```

      SUBROUTINE ZELAGA(N,A,CS,DZ)
      *****
      *   COMPUTES THE ZEROES OF THE LAGUERRE POLYNOMIAL OF DEGREE N
      *   N   = THE NUMBER OF ZEROES
      *   A   = PARAMETER >-1
      *   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
      *   DZ  = DERIVATIVES OF THE SCALED FUNCTIONS AT THE ZEROES, DZ(I), I=1,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1)
      IF (N .EQ. 0) RETURN

      A1 = A+1.D0
      CS(1) = A1
      DZ(1) = -4.D0/(A1*(A+5.D0))
      IF (N .EQ. 1) RETURN

      PI = 3.14159265358979323846D0
      DN = DFLOAT(N)
      C1 = 2.D0*DN+A1
      DO 1 M=1,N
         DM = DFLOAT(M)
         C2 = 2.D0*(DN+.75D0-DM)*PI/C1
         XN = (C2+PI)/2.D0
      DO 2 IT=1,8
         XP = XN
         XN = (DSIN(XP)-XP*DCOS(XP)+C2)/(1.D0-DCOS(XP))
2      CONTINUE
         Z = (DCOS(XN/2.D0))**2
         ZD = 1.D0/(Z-1.D0)
         CSX= 2.D0*C1*Z-((1.25D0*ZD+1.D0)*ZD-1.D0+3.D0*A*A)/(6.D0*C1)
      DO 3 IT=1,6
         CALL VALASF(N,A,CSX,Y,DY)
         CSX = CSX-Y/DY
3      CONTINUE
         CS(M) = CSX
         DZ(M) = DY
1      CONTINUE
4      IN = 0
      DO 5 M=1,N-1
         IF(CS(M) .LE. CS(M+1)) GOTO 5
         CSM = CS(M)
         CS(M) = CS(M+1)
         CS(M+1) = CSM
         DZM = DZ(M)
         DZ(M) = DZ(M+1)
         DZ(M+1) = DZM
         IN = 1
5      CONTINUE
         IF(IN .EQ. 1) GOTO 4

      RETURN
      END

```

ZEHEGA

For any $n \geq 1$, the routine computes the n zeroes $\xi_i^{(n)} \in \mathbf{R}$, $1 \leq i \leq n$, of the Hermite polynomial H_n . These are related to the zeroes of the Laguerre polynomials by the formulas

$$L_{n/2}^{(-1/2)}([\xi_i^{(n)}]^2) = 0 \quad 1 \leq i \leq n, \quad \text{if } n \text{ is even,}$$

$$\xi_i^{(n)} L_{(n-1)/2}^{(1/2)}([\xi_i^{(n)}]^2) = 0 \quad 1 \leq i \leq n, \quad \text{if } n \text{ is odd,}$$

Therefore, subroutine ZELAGA is called in this computation.

An output vector contains the zeroes in increasing order.

Another output vector contains the quantities $\hat{H}'_n(\xi_i^{(n)})$, $1 \leq i \leq n$, where \hat{H}_n is the scaled Hermite functions of degree n , introduced in the description of subroutine VAHESF.

<i>Input variable</i>	<i>Output variables</i>
N , the degree n	CS , vector containing the zeroes $\xi_i^{(n)}$ DZ , vector of the values $\hat{H}'_n(\xi_i^{(n)})$

Auxiliary routines: **ZELAGA, VALASF**

```

SUBROUTINE ZEHEGA(N,CS,DZ)
*****
*   COMPUTES THE ZEROES OF THE HERMITE POLYNOMIAL OF DEGREE N
*   N   = THE NUMBER OF ZEROES
*   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
*   DZ  = DERIVATIVES OF THE SCALED FUNCTIONS AT THE ZEROES, DZ(I), I=1,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1)
      IF (N .EQ. 0) RETURN

      CS(M+1) = 0.DO
      DZ(M+1) = 1.DO
      IF(N .EQ. 1) RETURN

      M = N/2
      IN = 2*N-4*M-1
      IF(IN .EQ. -1) THEN
        A = -.5DO
        CALL ZELAGA(M,A,CS,DZ)
      ELSE
        A = .5DO
        CALL ZELAGA(M,A,CS,DZ)
      ENDIF

      DO 1 I=1,M
        CSX = DSQRT(CS(M-I+1))
        CS(N-I+1) = CSX
        IF(IN .EQ. -1) THEN
          DZ(N-I+1) = 2.DO*CSX*DZ(M-I+1)
        ELSE
          DZ(N-I+1) = 2.DO*CSX*CSX*DZ(M-I+1)
        ENDIF
1      CONTINUE

      DO 2 I=1,M
        CS(I) = -CS(N-I+1)
        DZ(I) = DZ(N-I+1)*DFLOAT(IN)
2      CONTINUE

      RETURN
      END

```

ZEJAGL

For any $n \geq 1$, the routine computes the $n + 1$ nodes $\eta_i^{(n)} \in [-1, 1]$, $0 \leq i \leq n$, of the Jacobi Gauss-Lobatto formula

$$\sum_{i=0}^n f(\eta_i^{(n)}) \tilde{w}_i^{(n)} \approx \int_{-1}^1 f(x)(1-x)^\alpha(1+x)^\beta dx,$$

where the parameters satisfy $-\frac{1}{2} \leq \alpha \leq \frac{1}{2}$, $-\frac{1}{2} \leq \beta \leq \frac{1}{2}$. The quantities $\tilde{w}_i^{(n)}$, $0 \leq i \leq n$, are the weights of the formula.

One has $\eta_0^{(n)} = -1$, $\eta_n^{(n)} = 1$, and $\frac{d}{dx}P_n^{(\alpha,\beta)}(\eta_i^{(n)}) = 0$, $1 \leq i \leq n - 1$. In particular, the nodes in $] - 1, 1[$ satisfy $P_{n-1}^{(\alpha+1,\beta+1)}(\eta_i^{(n)}) = 0$, $1 \leq i \leq n - 1$. Hence, they are obtained with the same algorithm used for subroutine ZEJAGA.

An output vector contains the nodes in increasing order.

Another output vector contains the quantities $P_n^{(\alpha,\beta)}(\eta_i^{(n)})$, $0 \leq i \leq n$.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	ET , vector containing the nodes $\eta_i^{(n)}$
A , the parameter α	VN , vector of the values $P_n^{(\alpha,\beta)}(\eta_i^{(n)})$
B , the parameter β	

Auxiliary routine: **VAJAPO**

```

      SUBROUTINE ZEJAGL(N,A,B,ET,VN)
*****
*   COMPUTES THE NODES RELATIVE TO THE JACOBI GAUSS-LOBATTO FORMULA
*   N   = ORDER OF THE FORMULA
*   A   = PARAMETER BETWEEN -1/2 AND 1/2
*   B   = PARAMETER BETWEEN -1/2 AND 1/2
*   ET  = VECTOR OF THE NODES, ET(I), I=0,N
*   VN  = VALUES OF THE JACOBI POLYNOMIAL AT THE NODES, VN(I), I=0,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), VN(0:*)
      IF (N .EQ. 0) RETURN

      ET(0) = -1.DO
      ET(N) = 1.DO
      X = -1.DO
      CALL VAJAPO(N,A,B,X,Y,DY,D2Y)
      VN(0) = Y
      X = 1.DO
      CALL VAJAPO(N,A,B,X,Y,DY,D2Y)
      VN(N) = Y
      IF (N .EQ. 1) RETURN

      EPS= 1.E-17
      PH = 1.57079632679489661923D0
      AB = A+B
      DN = DFLOAT(N)
      C  = PH/(2.DO*DN+AB+1.DO)
      N1 = N-1
      A1 = A+1.DO
      B1 = B+1.DO
      DO 1 I=1,N1
         DI = DFLOAT(I)
         ETX = -DCOS(C*(4.DO*DI+AB+1.DO))
      DO 2 IT=1,8
         CALL VAJAPO(N1,A1,B1,ETX,Y,DY,D2Y)
         IF(DABS(Y) .LE. EPS) GOTO 3
         ETX = ETX-Y/DY
2      CONTINUE
3      IF(DABS(ETX) .LT. EPS) ETX=0.DO
         ET(I) = ETX
         VN(I) = -.5D0*DY*(1.DO-ETX*ETX)/DN
1      CONTINUE

      RETURN
      END

```

ZELEGL

For any $n \geq 1$, the routine computes the $n + 1$ nodes $\eta_i^{(n)} \in [-1, 1]$, $0 \leq i \leq n$, of the Legendre Gauss-Lobatto formula

$$\sum_{i=0}^n f(\eta_i^{(n)}) \tilde{w}_i^{(n)} \approx \int_{-1}^1 f(x) dx.$$

The quantities $\tilde{w}_i^{(n)}$, $0 \leq i \leq n$, are the weights of the formula.

One has $\eta_0^{(n)} = -1$, $\eta_n^{(n)} = 1$, and $\frac{d}{dx}P_n(\eta_i^{(n)}) = 0$, $1 \leq i \leq n - 1$. For the nodes in $] - 1, 1[$, the Newton method is used with the initial guess $\eta_i^{(n)} \approx -\cos \frac{i\pi}{n}$, $1 \leq i \leq n - 1$. The derivatives of the Legendre polynomial at a given point are obtained by the recursion formula given in the description of subroutine VALEPO.

An output vector contains the nodes in increasing order.

Another output vector contains the quantities $P_n(\eta_i^{(n)})$, $0 \leq i \leq n$.

<i>Input variable</i>	<i>Output variables</i>
N , the degree n	ET , vector containing the nodes $\eta_i^{(n)}$ VN , vector of the values $P_n(\eta_i^{(n)})$

Auxiliary routine: **VALEPO**

```

      SUBROUTINE ZELEGL(N,ET,VN)
      *****
      *   COMPUTES THE NODES RELATIVE TO THE LEGENDRE GAUSS-LOBATTO FORMULA
      *   N   = ORDER OF THE FORMULA
      *   ET  = VECTOR OF THE NODES, ET(I), I=0,N
      *   VN  = VALUES OF THE LEGENDRE POLYNOMIAL AT THE NODES, VN(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), VN(0:*)
      IF (N .EQ. 0) RETURN

      N2 = (N-1)/2
      SN = DFLOAT(2*N-4*N2-3)
      ET(0) = -1.DO
      ET(N) = 1.DO
      VN(0) = SN
      VN(N) = 1.DO
      IF (N .EQ. 1) RETURN

      ET(N2+1) = 0.DO
      X = 0.DO
      CALL VALEPO(N,X,Y,DY,D2Y)
      VN(N2+1) = Y
      IF(N .EQ. 2) RETURN

      PI = 3.14159265358979323846D0
      C = PI/DFLOAT(N)
      DO 1 I=1,N2
        ETX = DCOS(C*DFLOAT(I))
      DO 2 IT=1,8
        CALL VALEPO(N,ETX,Y,DY,D2Y)
        ETX = ETX-DY/D2Y
2      CONTINUE
        ET(I) = -ETX
        ET(N-I) = ETX
        VN(I) = Y*SN
        VN(N-I) = Y
1      CONTINUE

      RETURN
      END

```

ZECHGL

For any $n \geq 1$, the routine computes the $n + 1$ nodes $\eta_i^{(n)} \in [-1, 1]$, $0 \leq i \leq n$, of the Chebyshev Gauss-Lobatto formula

$$\sum_{i=0}^n f(\eta_i^{(n)}) \tilde{w}_i^{(n)} \approx \int_{-1}^1 f(x) \frac{dx}{\sqrt{1-x^2}}.$$

These, in increasing order, are given by

$$\eta_i^{(n)} = -\cos \frac{i\pi}{n} \quad 0 \leq i \leq n.$$

<i>Input variable</i>	<i>Output variable</i>
N , the degree n	ET , vector containing the nodes $\eta_i^{(n)}$


```
      SUBROUTINE ZECHGL(N,ET)
*****
*   COMPUTES THE NODES RELATIVE TO THE CHEBYSHEV GAUSS-LOBATTO FORMULA
*   N   = ORDER OF THE FORMULA
*   ET  = VECTOR OF THE NODES, ET(I), I=0,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*)
      IF (N .EQ. 0) RETURN

      ET(0) = -1.DO
      ET(N) = 1.DO
      IF (N .EQ. 1) RETURN

      N2 = (N-1)/2
      ET(N2+1) = 0.DO
      IF(N .EQ. 2) RETURN

      PI = 3.14159265358979323846D0
      C  = PI/DFLOAT(N)
      DO 1 I=1,N2
          ETX = DCOS(C*DFLOAT(I))
          ET(I) = -ETX
          ET(N-I) = ETX
1      CONTINUE

      RETURN
      END
```

ZELAGR

For any $n \geq 1$, the routine computes the n nodes $\eta_i^{(n)} \in [0, +\infty[$, $0 \leq i \leq n-1$, of the Laguerre Gauss-Radau formula

$$\sum_{i=0}^{n-1} f(\eta_i^{(n)}) \tilde{w}_i^{(n)} \approx \int_0^{+\infty} f(x) x^\alpha e^{-x} dx,$$

where $\alpha > -1$. The quantities $\tilde{w}_i^{(n)}$, $0 \leq i \leq n-1$, are the weights of the formula.

One has $\eta_0^{(n)} = 0$, and $\frac{d}{dx} L_n^{(\alpha)}(\eta_i^{(n)}) = 0$, $1 \leq i \leq n-1$. The nodes in $]0, +\infty[$ satisfy $L_{n-1}^{(\alpha+1)}(\eta_i^{(n)}) = 0$, $1 \leq i \leq n-1$. Hence, they are obtained with the same algorithm used for subroutine ZELAGA.

An output vector contains the nodes in increasing order.

Another output vector contains the quantities $\hat{L}_n^{(\alpha)}(\eta_i^{(n)})$, $0 \leq i \leq n-1$, where $\hat{L}_n^{(\alpha)}$ is the scaled Laguerre functions of degree n , introduced in the description of subroutine VALASF.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	ET , vector containing the nodes $\eta_i^{(n)}$
A , the parameter α	VN , vector of the values $\hat{L}_n^{(\alpha)}(\eta_i^{(n)})$

Auxiliary routine: **VALASF**

```

SUBROUTINE ZELAGR(N,A,ET,VN)
*****
*   COMPUTES THE NODES RELATIVE TO THE LAGUERRE GAUSS-RADAU FORMULA
*   N = ORDER OF THE FORMULA
*   A = PARAMETER >-1
*   ET = VECTOR OF THE NODES, ET(I), 0=1,N-1
*   VN = SCALED LAGUERRE FUNCTION AT THE NODES, VN(I), I=0,N-1
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), VN(0:*)
      IF (N .EQ. 0) RETURN

```

```

      ET(0) = 0.DO
      VN(0) = 1.DO
      IF (N .EQ. 1) RETURN

      A1 = A+1.DO
      PI = 3.14159265358979323846D0
      N1 = N-1
      DN = DFLOAT(N1)
      C1 = 2.DO*DN+A1+1.DO
      DO 1 M=1,N1
        DM = DFLOAT(M)
        C2 = 2.DO*(DN+.75D0-DM)*PI/C1
        XN = (C2+PI)/2.DO
      DO 2 IT=1,8
        XP = XN
        XN = (DSIN(XP)-XP*DCOS(XP)+C2)/(1.DO-DCOS(XP))
2      CONTINUE
        Z = (DCOS(XN/2.DO))**2
        ZD = 1.DO/(Z-1.DO)
        ETX= 2.DO*C1*Z-((1.25D0*ZD+1.DO)*ZD-1.DO+3.DO*A1*A1)/(6.DO*C1)
      DO 3 IT=1,6
        CALL VALASF(N1,A1,ETX,Y,DY)
        ETX = ETX-Y/DY
3      CONTINUE
        ET(M) = ETX
        CALL VALASF(N,A,ETX,Y,DY)
        VN(M) = Y
1      CONTINUE

4      IN = 0
      DO 5 M=1,N-2
        IF(ET(M) .LE. ET(M+1)) GOTO 5
        ETM = ET(M)
        ET(M) = ET(M+1)
        ET(M+1) = ETM
        VNM = VN(M)
        VN(M) = VN(M+1)
        VN(M+1) = VNM
        IN = 1
5      CONTINUE
        IF(IN .EQ. 1) GOTO 4

      RETURN
      END

```

WEJAGA

For any $n \geq 1$, the routine computes the n weights $w_i^{(n)}$, $1 \leq i \leq n$, of the Jacobi Gauss integration formula

$$\sum_{i=1}^n f(\xi_i^{(n)}) w_i^{(n)} \approx \int_{-1}^1 f(x)(1-x)^\alpha(1+x)^\beta dx,$$

where $\alpha > -1$, $\beta > -1$, and $\xi_i^{(n)}$, $1 \leq i \leq n$, are the zeroes of the Jacobi polynomial $P_n^{(\alpha,\beta)}$. For $1 \leq i \leq n$, the following expression is used:

$$w_j^{(n)} = 2^{\alpha+\beta+1} \frac{\Gamma(n+\alpha+1) \Gamma(n+\beta+1)}{n! \Gamma(n+\alpha+\beta+1)} \frac{1}{1-[\xi_j^{(n)}]^2} \left[\frac{d}{dx} P_n^{(\alpha,\beta)}(\xi_j^{(n)}) \right]^{-2}.$$

Both the zeroes and the quantities $\frac{d}{dx} P_n^{(\alpha,\beta)}(\xi_i^{(n)})$, $1 \leq i \leq n$, can be computed by calling subroutine ZEJAGA.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n A , the parameter α B , the parameter β CS , vector of the zeroes $\xi_i^{(n)}$ DZ , the values $\frac{d}{dx} P_n^{(\alpha,\beta)}(\xi_i^{(n)})$	WE , vector containing the weights $w_i^{(n)}$

Auxiliary routine: **GAMMAF**

```

      SUBROUTINE WEJAGA(N,A,B,CS,DZ,WE)
*****
*   COMPUTES THE WEIGHTS RELATIVE TO THE JACOBI GAUSS FORMULA
*   N   = ORDER OF THE FORMULA
*   A   = PARAMETER > -1
*   B   = PARAMETER > -1
*   CS  = ZEROES OF THE JACOBI POLYNOMIAL, CS(I), I=1,N
*   DZ  = VECTOR OF THE DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
*   WE  = VECTOR OF THE WEIGHTS, WE(I), I=1,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1), WE(1)
      IF (N .EQ. 0) RETURN

      AB = A+B+2.DO
      A2 = A+2.DO
      B2 = B+2.DO
      CALL GAMMAF(A2,GA2)
      CALL GAMMAF(B2,GB2)
      CALL GAMMAF(AB,GAB)
      C  = .5DO*(2.DO**AB)*GA2*GB2/GAB
      DO 1 M=2,N
          DM = DFLOAT(M)
          C  = C*(DM+A)*(DM+B)/(DM*(DM+A+B))
1      CONTINUE

      DO 2 I=1,N
          X  = CS(I)
          DY = DZ(I)
          WE(I) = C/((1.DO-X*X)*DY*DY)
2      CONTINUE

      RETURN
      END

```

WELEGA

For any $n \geq 1$, the routine computes the n weights $w_i^{(n)}$, $1 \leq i \leq n$, of the Legendre Gauss integration formula

$$\sum_{i=1}^n f(\xi_i^{(n)}) w_i^{(n)} \approx \int_{-1}^1 f(x) dx,$$

where $\xi_i^{(n)}$, $1 \leq i \leq n$, are the zeroes of the Legendre polynomial P_n .

The following expression is used:

$$w_i^{(n)} = \frac{2}{1 - [\xi_i^{(n)}]^2} [P'_n(\xi_i^{(n)})]^{-2}, \quad 1 \leq i \leq n.$$

Both the zeroes and the quantities $P'_n(\xi_i^{(n)})$, $1 \leq i \leq n$, can be computed by calling subroutine ZELEGA.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n CS , vector of the zeroes of P_n DZ , values of P'_n at the zeroes	WE , vector of the weights $w_i^{(n)}$

```
      SUBROUTINE WELEGA(N,CS,DZ,WE)
*****
*   COMPUTES THE WEIGHTS RELATIVE TO THE LEGENDRE GAUSS FORMULA
*   N   = ORDER OF THE FORMULA
*   CS  = ZEROES OF THE LEGENDRE POLYNOMIAL, CS(I), I=1,N
*   DZ  = VECTOR OF THE DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
*   WE  = VECTOR OF THE WEIGHTS, WE(I), I=1,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1), WE(1)
      IF (N .EQ. 0) RETURN

      N2 = N/2
      DO 1 I=1,N2
        X   = CS(I)
        DY  = DZ(I)
        WEX = 2.DO/((1.DO-X*X)*DY*DY)
        WE(I) = WEX
        WE(N-I+1) = WEX
1      CONTINUE

      IF(N .EQ. 2*N2) RETURN
        DY = DZ(N2+1)
        WE(N2+1) = 2.DO/(DY*DY)

      RETURN
      END
```

WECHGA

For any $n \geq 1$, the routine computes the n weights $w_i^{(n)}$, $1 \leq i \leq n$, of the Chebyshev Gauss integration formula

$$\sum_{i=1}^n f(\xi_i^{(n)}) w_i^{(n)} \approx \int_{-1}^1 f(x) \frac{dx}{\sqrt{1-x^2}},$$

where $\xi_i^{(n)}$, $1 \leq i \leq n$, are the zeroes of the Chebyshev polynomial T_n .

The weights are explicitly given by

$$w_i^{(n)} = \frac{\pi}{n} \quad 1 \leq i \leq n.$$

<i>Input variable</i>	<i>Output variable</i>
N , the degree n	WE , vector containing the weights $w_i^{(n)}$


```
      SUBROUTINE WECHGA(N,WE)
*****
*   COMPUTES THE WEIGHTS RELATIVE TO THE CHEBYSHEV GAUSS FORMULA
*   N   = ORDER OF THE FORMULA
*   WE  = VECTOR OF THE WEIGHTS, WE(I), I=1,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION WE(1)
      IF (N .EQ. 0) RETURN

      PI = 3.14159265358979323846D0
      C  = PI/DFLOAT(N)
      DO 1 I=1,N
        WE(I) = C
1     CONTINUE

      RETURN
      END
```

WELAGA

For any $n \geq 1$, the routine computes the n weights $w_i^{(n)}$, $1 \leq i \leq n$, of the Laguerre Gauss integration formula

$$\sum_{i=1}^n f(\xi_i^{(n)}) w_i^{(n)} \approx \int_0^{+\infty} f(x) x^\alpha e^{-x} dx,$$

where $\alpha > -1$, and $\xi_i^{(n)}$, $1 \leq i \leq n$, are the zeroes of the Laguerre polynomial $L_n^{(\alpha)}$, which can be computed by subroutine ZELAGA.

The following expression is used:

$$w_i^{(n)} = \frac{\Gamma(n + \alpha + 1) \xi_i^{(n)}}{(n + 1)! (n + 1)} \left[L_{n+1}^{(\alpha)}(\xi_i^{(n)}) \right]^{-2}, \quad 1 \leq i \leq n.$$

<i>Input variables</i>	<i>Output variable</i>
N , the degree n A , the parameter α CS , vector of the zeroes of $L_n^{(\alpha)}$	WE , vector of the weights $w_i^{(n)}$

Auxiliary routines: **GAMMAF, VALAPO**

```
      SUBROUTINE WELAGA(N,A,CS,WE)
*****
*   COMPUTES THE WEIGHTS RELATIVE TO THE LAGUERRE GAUSS FORMULA
*   N   = ORDER OF THE FORMULA
*   A   = PARAMETER > -1
*   CS  = ZEROES OF THE LAGUERRE POLYNOMIAL, CS(I), I=1,N
*   WE  = VECTOR OF THE WEIGHTS, WE(I), I=1,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), WE(1)
      IF (N .EQ. 0) RETURN

      A1 = A+1.DO
      CALL GAMMAF(A1,GA1)
      N1 = N+1
      DN = DFLOAT(N1)
      C  = GA1/DN
      DO 1 M=1,N
          DM = DFLOAT(M)
          C  = C*(DM+A)/(DM+1.DO)
1      CONTINUE

      DO 2 I=1,N
          X =CS(I)
          CALL VALAPO(N1,A,X,Y,DY,D2Y)
          WE(I) = C*X/(Y*Y)
2      CONTINUE

      RETURN
      END
```

WEHEGA

For any $n \geq 1$, the routine computes the n weights $w_i^{(n)}$, $1 \leq i \leq n$, of the Hermite Gauss integration formula

$$\sum_{i=1}^n f(\xi_i^{(n)}) w_i^{(n)} \approx \int_{-\infty}^{+\infty} f(x) e^{-x^2} dx,$$

where $\xi_i^{(n)}$, $1 \leq i \leq n$, are the zeroes of the Hermite polynomial H_n , which can be computed by subroutine ZEHEGA.

For $1 \leq i \leq n$, the following expression is used:

$$w_i^{(n)} = \frac{\sqrt{\pi}}{n} 2^{n-1} (n-1)! [H_{n-1}(\xi_i^{(n)})]^{-2} = \frac{\sqrt{\pi}}{n} [\tilde{H}_{n-1}(\xi_i^{(n)})]^{-2},$$

where $\tilde{H}_k = (2^k k!)^{-1/2} H_k$, $k \in \mathbf{N}$, is obtained by recursion

$$\begin{cases} \tilde{H}_0(x) = 1, \\ \tilde{H}_1(x) = \sqrt{2}x, \\ \tilde{H}_k(x) = \sqrt{\frac{2}{k}}x\tilde{H}_{k-1}(x) - \sqrt{\frac{k-1}{k}}\tilde{H}_{k-2}(x), \quad k \geq 2. \end{cases}$$

<i>Input variables</i>	<i>Output variable</i>
N , the degree n CS , vector of the zeroes of H_n	WE , vector of the weights $w_i^{(n)}$

```

      SUBROUTINE WEHEGA(N,CS,WE)
*****
*   COMPUTES THE WEIGHTS RELATIVE TO THE HERMITE GAUSS FORMULA
*   N   = ORDER OF THE FORMULA
*   CS  = ZEROES OF THE HERMITE POLYNOMIAL, CS(I), I=1,N
*   WE  = VECTOR OF THE WEIGHTS, WE(I), I=1,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), WE(1)
      IF (N .EQ. 0) RETURN

      PR = 1.77245385090551588D0
      R2 = 1.41421356237309515D0
      C  = PR/DFLOAT(N)
      N2 = N/2
DO 1 I=1,N2
      X  = CS(I)
      YP = 1.DO
      Y  = R2*X
DO 2 K=2,N-1
      DK = DFLOAT(K)
      RK = DSQRT(DK)
      QK = DSQRT(DK-1.DO)
      YM = Y
      Y  = (R2*X*Y-QK*YP)/RK
      YP = YM
2   CONTINUE
      WEX = C/(Y*Y)
      WE(I) = WEX
      WE(N-I+1) = WEX
1   CONTINUE

      IF(N .EQ. 2*N2) RETURN
      Y = 1.DO
DO 3 K=2,N-1,2
      DK = DFLOAT(K)
      Y  = Y*DSQRT((DK-1.DO)/DK)
3   CONTINUE
      WE(N2+1) = C/(Y*Y)

      RETURN
      END

```

WEJAGL

For any $n \geq 1$, the routine computes the $n + 1$ weights $\tilde{w}_j^{(n)}$, $0 \leq j \leq n$, of the Jacobi Gauss-Lobatto formula of order n , introduced in the description of subroutine ZEJAGL.

For $n = 1$, one has $\tilde{w}_0^{(n)} = 2^{\alpha+\beta+1} \frac{\Gamma(\alpha+2) \Gamma(\beta+1)}{\Gamma(\alpha+\beta+3)}$, $\tilde{w}_n^{(n)} = 2^{\alpha+\beta+1} \frac{\Gamma(\alpha+1) \Gamma(\beta+2)}{\Gamma(\alpha+\beta+3)}$.

For $n \geq 2$, the following expression is used:

$$\tilde{w}_j^{(n)} = \begin{cases} K_n(\alpha, \beta) \sum_{m=1}^{n-1} (1 - \eta_m^{(n)}) & \text{if } j = 0, \\ \tilde{K}_n(\alpha, \beta) \left[P_n^{(\alpha, \beta)}(\eta_j^{(n)}) \frac{d}{dx} P_{n-1}^{(\alpha, \beta)}(\eta_j^{(n)}) \right]^{-1} & \text{if } 1 \leq j \leq n-1, \\ K_n(\beta, \alpha) \sum_{m=1}^{n-1} (1 + \eta_m^{(n)}) & \text{if } j = n, \end{cases}$$

with

$$K_n(\alpha, \beta) = 2^{\alpha+\beta} \frac{\Gamma(\beta+1) \Gamma(\beta+2) (2n + \alpha + \beta) (n-2)! \Gamma(n + \alpha)}{\Gamma(n + \alpha + \beta + 2) \Gamma(n + \beta + 1)},$$

$$\tilde{K}_n(\alpha, \beta) = -2^{\alpha+\beta} \frac{(2n + \alpha + \beta) \Gamma(n + \alpha) \Gamma(n + \beta)}{(n + \alpha + \beta + 1) n! \Gamma(n + \alpha + \beta)},$$

where $\alpha > -1$, $\beta > -1$, and $\eta_j^{(n)}$, $0 \leq j \leq n$, are the Gauss-Lobatto nodes, i.e., $\eta_0^{(n)} = -1$, $\eta_n^{(n)} = 1$, and $\frac{d}{dx} P_n^{(\alpha, \beta)}(\eta_j^{(n)}) = 0$, $1 \leq j \leq n-1$, which can be obtained by calling subroutine ZEJAGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n A , the parameter α B , the parameter β ET , vector of the nodes $\eta_j^{(n)}$	WT , vector containing the weights $\tilde{w}_i^{(n)}$

Auxiliary routines: **GAMMAF, VAJAPO**

```

      SUBROUTINE WEJAGL(N,A,B,ET,WT)
      *****
      *   COMPUTES THE WEIGHTS RELATIVE TO THE JACOBI GAUSS-LOBATTO FORMULA
      *   N   = ORDER OF THE FORMULA
      *   A   = PARAMETER > -1
      *   B   = PARAMETER > -1
      *   ET  = JACOBI GAUSS-LOBATTO NODES, ET(I), I=0,N
      *   WT  = VECTOR OF THE WEIGHTS, WT(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), WT(0:*)
      IF (N .EQ. 0) RETURN

      A1 = A+1.DO
      B1 = B+1.DO
      AB = A+B
      AB2 = A1+B1
      CALL GAMMAF(A1,GA1)
      CALL GAMMAF(B1,GB1)
      CALL GAMMAF(AB2,GAB2)
      C = (2.DO**AB)*GA1*GB1/GAB2
      WT(0) = 2.DO*C*A1/AB2
      WT(N) = 2.DO*C*B1/AB2
      IF (N .EQ. 1) RETURN

      N1 = N-1
      DN = DFLOAT(N)
      C = C*(2.DO*DN+AB)/(DN+AB+1.DO)
      C1 = C*A1/((B+2.DO)*AB2)
      C2 = C*B1/((A+2.DO)*AB2)
      C3 = .5DO*C*A1*B1
      DO 1 K=1,N-2
      DK = DFLOAT(K)
      C1 = C1*(DK+A1)*DK/((DK+AB2)*(DK+B+2.DO))
      C2 = C2*(DK+B1)*DK/((DK+AB2)*(DK+A+2.DO))
      C3 = C3*(DK+A1)*(DK+B1)/((DK+2.DO)*(DK+AB+1.DO))
1     CONTINUE

      SU = 0.DO
      DO 2 M=1,N1
      SU = SU+ET(M)
2     CONTINUE
      WT(0) = C1*(DN-1.DO-SU)
      WT(N) = C2*(DN-1.DO+SU)
      DO 3 I=1,N1
      X = ET(I)
      CALL VAJAPO(N,A,B,X,Y,DY,D2Y)
      C4 = -C3/Y
      CALL VAJAPO(N1,A,B,X,Y,DY,D2Y)
      WT(I) = C4/DY
3     CONTINUE
      RETURN
      END

```

WELEGL

For any $n \geq 1$, the routine computes the $n + 1$ weights $\tilde{w}_j^{(n)}$, $0 \leq j \leq n$, of the Legendre Gauss-Lobatto formula of order n , introduced in the description of subroutine ZELEGL.

The following expression is used:

$$\tilde{w}_j^{(n)} = \frac{2}{n(n+1)} [P_n(\eta_j^{(n)})]^{-2}, \quad 0 \leq j \leq n,$$

where $\eta_j^{(n)}$, $0 \leq j \leq n$, are the Gauss-Lobatto nodes, i.e., $\eta_0^{(n)} = -1$, $\eta_n^{(n)} = 1$, and $\frac{d}{dx}P_n(\eta_j^{(n)}) = 0$, $1 \leq j \leq n-1$. Both the nodes and the quantities $P_n(\eta_j^{(n)})$, $0 \leq j \leq n$, can be computed by calling subroutine ZELEGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n ET , vector of the nodes $\eta_j^{(n)}$ VN , values of P_n at the nodes	WT , vector containing the weights $\tilde{w}_j^{(n)}$


```
      SUBROUTINE WELEGL(N,ET,VN,WT)
*****
*   COMPUTES THE WEIGHTS RELATIVE TO THE LEGENDRE GAUSS-LOBATTO FORMULA
*   N   = ORDER OF THE FORMULA
*   ET  = LEGENDRE GAUSS-LOBATTO NODES, ET(I), I=0,N
*   VN  = VALUES OF THE LEGENDRE POLYNOMIAL AT THE NODES, VN(I), I=0,N
*   WT  = VECTOR OF THE WEIGHTS, WT(I), I=0,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), VN(0:*), WT(0:*)
      IF (N .EQ. 0) RETURN

      N2 = (N-1)/2
      DN = DFLOAT(N)
      C  = 2.DO/(DN*(DN+1.DO))
DO 1 I=0,N2
      X = ET(I)
      Y = VN(I)
      WTX = C/(Y*Y)
      WT(I) = WTX
      WT(N-I) = WTX
1  CONTINUE

      IF(N-1 .EQ. 2*N2) RETURN
      X = 0.DO
      Y = VN(N2+1)
      WT(N2+1) = C/(Y*Y)

      RETURN
      END
```

WECHGL

For any $n \geq 1$, the routine computes the $n + 1$ weights $\tilde{w}_j^{(n)}$, $0 \leq j \leq n$, of the Chebyshev Gauss-Lobatto formula of order n , introduced in the description of subroutine ZECHGL.

These are given by

$$\tilde{w}_0^{(n)} = \tilde{w}_n^{(n)} = \frac{\pi}{2n}, \quad \tilde{w}_j^{(n)} = \frac{\pi}{n} \quad 1 \leq j \leq n - 1.$$

<i>Input variable</i>	<i>Output variable</i>
N , the degree n	WT , vector containing the weights $\tilde{w}_j^{(n)}$

```
      SUBROUTINE WECHGL(N,WT)
*****
*   COMPUTES THE WEIGHTS RELATIVE TO THE CHEBYSHEV GAUSS-LOBATTO FORMULA
*   N   = ORDER OF THE FORMULA
*   WT  = VECTOR OF THE WEIGHTS, WT(I), I=0,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION WT(0:*)
      IF (N .EQ. 0) RETURN

      PI = 3.14159265358979323846D0
      C  = PI/DFLOAT(N)
      C2 = .5D0*C
      WT(0) =C2
      WT(N) =C2
      IF (N .EQ. 1) RETURN

      DO 1 I=1,N-1
          WT(I) = C
1      CONTINUE

      RETURN
      END
```

WELAGR

For any $n \geq 1$, the routine computes the n weights $\tilde{w}_j^{(n)}$, $0 \leq j \leq n-1$, of the Laguerre Gauss-Radau formula of order n , introduced in the description of subroutine ZELAGR.

The following expression is used:

$$\tilde{w}_j^{(n)} = \begin{cases} \frac{(\alpha+1) \Gamma^2(\alpha+1) (n-1)!}{\Gamma(n+\alpha+1)} & \text{if } j = 0, \\ \frac{\Gamma(n+\alpha)}{n!} \left[L_n^{(\alpha)}(\eta_j^{(n)}) \frac{d}{dx} L_{n-1}^{(\alpha)}(\eta_j^{(n)}) \right]^{-1} & \text{if } 1 \leq j \leq n-1, \end{cases}$$

where $\alpha > -1$, and $\eta_j^{(n)}$, $0 \leq j \leq n-1$, are the Gauss-Radau nodes, i.e., $\eta_0^{(n)} = 0$ and $\frac{d}{dx} L_n^{(\alpha)}(\eta_j^{(n)}) = 0$, $1 \leq j \leq n-1$, which can be computed by calling subroutine ZELAGR.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n A , the parameter α ET , vector of the nodes $\eta_j^{(n)}$	WT , vector containing the weights $\tilde{w}_j^{(n)}$

Auxiliary routines: **GAMMAF, VALAPO**

```

      SUBROUTINE WELAGR(N,A,ET,WT)
      *****
      *   COMPUTES THE WEIGHTS RELATIVE TO THE LAGUERRE GAUSS-RADAU FORMULA
      *   N   = ORDER OF THE FORMULA
      *   A   = PARAMETER > -1
      *   ET  = LAGUERRE GAUSS-RADAU NODES, ET(I), I=0,N-1
      *   WT  = VECTOR OF THE WEIGHTS, WT(I), I=0,N-1
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), WT(0:*)
      IF (N .EQ. 0) RETURN

      A1 = A+1.DO
      CALL GAMMAF(A1,GA1)
      C1 = GA1
      WT(0) = C1
      IF (N .EQ. 1) RETURN

      N1 = N-1
      C2 = GA1
      DO 1 K=1,N1
         DK = DFLOAT(K)
         C1 = C1*DK/(DK+A1)
         C2 = C2*(DK+A)/(DK+1.DO)
1      CONTINUE
         WT(0) = C1
      DO 2 I=1,N1
         X = ET(I)
         CALL VALAPO(N,A,X,Y,DY,D2Y)
         C3 = C2/Y
         CALL VALAPO(N1,A,X,Y,DY,D2Y)
         WT(I) = C3/DY
2      CONTINUE

      RETURN
      END

```

WECHCC

For any $n \geq 1$, the routine computes the $2n + 1$ weights $\chi_j^{(2n)}$, $0 \leq j \leq 2n$, of the Clenshaw-Curtis formula of order $2n$

$$\int_{-1}^1 f(x) dx \approx \sum_{j=0}^{2n} f\left(\cos \frac{j\pi}{2n}\right) \chi_j^{(2n)}.$$

The following expression is used:

$$\chi_j^{(2n)} = \begin{cases} \frac{1}{4n^2 - 1} & \text{if } j = 0, \\ \frac{1}{n} \left[1 - \frac{(-1)^j}{4n^2 - 1} + \sum_{k=1}^{n-1} \frac{2}{1 - 4k^2} \cos \frac{kj\pi}{n} \right] & \text{if } 1 \leq j \leq n, \\ \chi_{2n-j}^{(2n)} & \text{if } n + 1 \leq j \leq 2n. \end{cases}$$

<i>Input variable</i>	<i>Output variable</i>
N , the parameter n	WK , vector containing the weights $\chi_j^{(2n)}$

```

      SUBROUTINE WECHCC(N,WK)
      *****
      *   COMPUTES THE WEIGHTS OF THE CLENSHAW-CURTIS FORMULA OF ORDER 2*N
      *   N   = INTEGER PARAMETER
      *   WK  = VECTOR OF THE WEIGHTS, WK(I), I=0,2*N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION WK(0:*)
      IF (N .EQ. 0) RETURN

      PI = 3.14159265358979323846D0
      M  = 2*N
      DN = DFLOAT(N)
      DM = DFLOAT(M)
      C  = 1.DO/(DM*DM-1.DO)
      WK(0) = C
      WK(M) = C
      WK(N) = 1.33333333333333333333D0
      IF (N .EQ. 1) RETURN

      DO 1 J=1,N-1
        DJ = DFLOAT(J)
        SU = 1.DO-((-1.DO)**J)*C
      DO 2 K=1,N-1
        DK = 2.DO*DFLOAT(K)
        SU = SU+2.DO*DCOS(DJ*DK*PI/DM)/(1.DO-DK*DK)
2     CONTINUE
        WK(J) = SU/DN
        WK(M-J) = SU/DN
1     CONTINUE

        SU = 1.DO-((-1.DO)**N)*C
      DO 3 K=1,N-1
        DK = 2.DO*DFLOAT(K)
        SU = SU+2.DO*((-1.DO)**K)/(1.DO-DK*DK)
3     CONTINUE
        WK(N) = SU/DN

      RETURN
      END

```

INJAGA

For any $n \geq 1$, the routine computes the value at the point x of a polynomial q of degree at most $n-1$, determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the Jacobi polynomial $P_n^{(\alpha, \beta)}$.

One has

$$q(x) = \sum_{j=1}^n q(\xi_j^{(n)}) l_j^{(n)}(x),$$

where for $1 \leq j \leq n$

$$l_j^{(n)}(x) = \begin{cases} \left[\frac{d}{dx} P_n^{(\alpha, \beta)}(\xi_j^{(n)}) \right]^{-1} \frac{P_n^{(\alpha, \beta)}(x)}{x - \xi_j^{(n)}} & \text{if } x \neq \xi_j^{(n)}, \\ 1 & \text{if } x = \xi_j^{(n)}. \end{cases}$$

The zeroes and the values $\frac{d}{dx} P_n^{(\alpha, \beta)}(\xi_j^{(n)})$, $1 \leq j \leq n$, can be obtained by calling subroutine ZJAGA.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n A , the parameter α B , the parameter β CS , vector of the zeroes $\xi_j^{(n)}$ DZ , the values $\frac{d}{dx} P_n^{(\alpha, \beta)}(\xi_j^{(n)})$ QZ , values of q at the zeroes X , the point x	QX , the value of q in x

Auxiliary routine: **VAJAPO**


```

      SUBROUTINE INJAGA(N,A,B,CS,DZ,QZ,X,QX)
*****
*   COMPUTES THE VALUE AT A GIVEN POINT OF A POLYNOMIAL INDIVIDUATED
*   BY THE VALUES ATTAINED AT THE ZEROES OF THE JACOBI POLYNOMIAL
*   N   = THE NUMBER OF ZEROES
*   A   = PARAMETER > -1
*   B   = PARAMETER > -1
*   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
*   DZ  = JACOBI POLYNOMIAL DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
*   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
*   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
*   QX  = VALUE OF THE POLYNOMIAL IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1), QZ(1)
      IF (N .EQ. 0) RETURN

      EPS = 1.D-14
      CALL VAJAP0(N,A,B,X,Y,DY,D2Y)
      QX = 0.DO
      DO 1 J=1,N
          ED = X-CS(J)
          IF(DABS(ED) .LT. EPS) THEN
              QX = QZ(J)
          RETURN
          ELSE
              QX = QX+QZ(J)*Y/(DZ(J)*ED)
          ENDIF
1      CONTINUE

      RETURN
      END

```

INLEGA

For any $n \geq 1$, the routine computes the value at the point x of a polynomial q of degree at most $n - 1$, determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the Legendre polynomial P_n .

One has

$$q(x) = \sum_{j=1}^n q(\xi_j^{(n)}) l_j^{(n)}(x),$$

where for $1 \leq j \leq n$

$$l_j^{(n)}(x) = \begin{cases} [P_n'(\xi_j^{(n)})]^{-1} \frac{P_n(x)}{x - \xi_j^{(n)}} & \text{if } x \neq \xi_j^{(n)}, \\ 1 & \text{if } x = \xi_j^{(n)}. \end{cases}$$

The zeroes and the values $P_n'(\xi_j^{(n)})$, $1 \leq j \leq n$, can be obtained by calling subroutine ZELEGA.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n CS , vector of the zeroes $\xi_j^{(n)}$ DZ , the values $P_n'(\xi_j^{(n)})$ QZ , values of q at the zeroes X , the point x	QX , the value of q in x

Auxiliary routine: **VALEPO**

```
      SUBROUTINE INLEGA(N,CS,DZ,QZ,X,QX)
*****
*   COMPUTES THE VALUE AT A GIVEN POINT OF A POLYNOMIAL INDIVIDUATED
*   BY THE VALUES ATTAINED AT THE ZEROES OF THE LEGENDRE POLYNOMIAL
*   N   = THE NUMBER OF ZEROES
*   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
*   DZ  = LEGENDRE POLYNOMIAL DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
*   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
*   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
*   QX  = VALUE OF THE POLYNOMIAL IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1), QZ(1)
      IF (N .EQ. 0) RETURN

      EPS = 1.D-14
      CALL VALEPO(N,X,Y,DY,D2Y)
      QX = 0.DO
      DO 1 J=1,N
          ED = X-CS(J)
          IF(DABS(ED) .LT. EPS) THEN
              QX = QZ(J)
          RETURN
          ELSE
              QX = QX+QZ(J)*Y/(DZ(J)*ED)
          ENDIF
1      CONTINUE

      RETURN
      END
```

INCHGA

For any $n \geq 1$, the routine computes the value at the point x of a polynomial q of degree at most $n - 1$, determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the Chebyshev polynomial T_n .

One has

$$q(x) = \sum_{j=1}^n q(\xi_j^{(n)}) l_j^{(n)}(x),$$

where for $1 \leq j \leq n$

$$l_j^{(n)}(x) = \begin{cases} [T_n'(\xi_j^{(n)})]^{-1} \frac{T_n(x)}{x - \xi_j^{(n)}} & \text{if } x \neq \xi_j^{(n)}, \\ 1 & \text{if } x = \xi_j^{(n)}. \end{cases}$$

The values $T_n'(\xi_j^{(n)})$, $1 \leq j \leq n$, can be obtained by calling subroutine ZECHGA.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n DZ , the values $T_n'(\xi_j^{(n)})$ QZ , values of q at the zeroes X , the point x	QX , the value of q in x

Auxiliary routine: **VACHPO**

```

      SUBROUTINE INCHGA(N,DZ,QZ,X,QX)
      *****
      *   COMPUTES THE VALUE AT A GIVEN POINT OF A POLYNOMIAL INDIVIDUATED
      *   BY THE VALUES ATTAINED AT THE ZEROES OF THE CHEBYSHEV POLYNOMIAL
      *   N   = THE NUMBER OF ZEROES
      *   DZ  = CHEBYSHEV POLYNOMIAL DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
      *   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
      *   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
      *   QX  = VALUE OF THE POLYNOMIAL IN X
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION DZ(1), QZ(1)
      IF (N .EQ. 0) RETURN

      EPS = 1.D-14
      PH = 1.57079632679489661923D0
      DN = DFLOAT(N)
      C  = PH/DN
      CALL VACHPO(N,X,Y,DY,D2Y)
      QX = 0.DO
      DO 1 J=1,N
          ED = X+DCOS(C*(2.DO*DFLOAT(J)-1.DO))
          IF(DABS(ED) .LT. EPS) THEN
              QX = QZ(J)
          RETURN
          ELSE
              QX = QX+QZ(J)*Y/(DZ(J)*ED)
          ENDIF
      1  CONTINUE

      RETURN
      END

```

INLAGA

For any $n \geq 1$, the routine computes the value at the point x of a polynomial q of degree at most $n - 1$ determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the Laguerre polynomial $L_n^{(\alpha)}$.

One has

$$q(x) = \sum_{j=1}^n q(\xi_j^{(n)}) l_j^{(n)}(x),$$

where for $1 \leq j \leq n$

$$l_j^{(n)}(x) = \begin{cases} \left[\frac{d}{dx} \hat{L}_n^{(\alpha)}(\xi_j^{(n)}) \right]^{-1} \frac{\hat{L}_n^{(\alpha)}(x)}{x - \xi_j^{(n)}} \prod_{k=1}^n \frac{4k + x}{4k + \xi_j^{(n)}} & \text{if } x \neq \xi_j^{(n)}, \\ 1 & \text{if } x = \xi_j^{(n)}. \end{cases}$$

Here, $\hat{L}_n^{(\alpha)}$ is the scaled Laguerre function of degree n , introduced in the description of subroutine VALASF. The zeroes and the values $\frac{d}{dx} \hat{L}_n^{(\alpha)}(\xi_j^{(n)})$, $1 \leq j \leq n$, can be obtained by calling subroutine ZELAGA.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n	QX , the value of q in x
A , the parameter α	
CS , vector of the zeroes $\xi_j^{(n)}$	
DZ , the values $\frac{d}{dx} \hat{L}_n^{(\alpha)}(\xi_j^{(n)})$	
QZ , values of q at the zeroes	
X , the point x	

Auxiliary routine: **VALASF**

```

      SUBROUTINE INLAGA(N,A,CS,DZ,QZ,X,QX)
      *****
      *   COMPUTES THE VALUE AT A GIVEN POINT OF A POLYNOMIAL INDIVIDUATED
      *   BY THE VALUES ATTAINED AT THE ZEROES OF THE LAGUERRE POLYNOMIAL
      *   N   = THE NUMBER OF ZEROES
      *   A   = PARAMETER > -1
      *   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
      *   DZ  = SCALED LAGUERRE DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
      *   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
      *   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
      *   QX  = VALUE OF THE POLYNOMIAL IN X
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1), QZ(1)
      IF (N .EQ. 0) RETURN

      EPS = 1.D-14
      CALL VALASF(N,A,X,Y,DY)
      QX = 0.DO
      DO 1 J=1,N
        ED = X-CS(J)
        IF(DABS(ED) .LT. EPS) THEN
          QX = QZ(J)
        RETURN
      ELSE
        PR = 1.DO
        DO 2 K=1,N
          DK = 4.DO*DFLOAT(K)
          PR = PR*(DK+X)/(DK+CS(J))
2      CONTINUE
        QX = QX+QZ(J)*Y*PR/(DZ(J)*ED)
      ENDIF
1      CONTINUE

      RETURN
      END

```

INHEGA

For any $n \geq 1$, the routine computes the value at the point x of a polynomial q of degree at most $n-1$, determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the Hermite polynomial H_n .

One has

$$q(x) = \sum_{j=1}^n q(\xi_j^{(n)}) l_j^{(n)}(x),$$

where for $1 \leq j \leq n$

$$l_j^{(n)}(x) = \begin{cases} [\hat{H}'_n(\xi_j^{(n)})]^{-1} \frac{\hat{H}_n(x)}{x - \xi_j^{(n)}} \prod_{k=1}^m \frac{4k + x^2}{4k + [\xi_j^{(n)}]^2} & \text{if } x \neq \xi_j^{(n)}, \\ 1 & \text{if } x = \xi_j^{(n)}. \end{cases}$$

Here, \hat{H}_n is the scaled Hermite function of degree n introduced in the description of subroutine VAHESF, and $m = n/2$ if n is even, $m = (n-1)/2$ if n is odd.

The zeroes and the values $\hat{H}'_n(\xi_j^{(n)})$, $1 \leq j \leq n$, can be obtained by calling subroutine ZEHEGA.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n	QX , the value of q in x
CS , vector of the zeroes $\xi_j^{(n)}$	
DZ , the values $\hat{H}'_n(\xi_j^{(n)})$	
QZ , values of q at the zeroes	
X , the point x	

Auxiliary routine: **VAHESF**


```

      SUBROUTINE INHEGA(N,CS,DZ,QZ,X,QX)
      *****
      *   COMPUTES THE VALUE AT A GIVEN POINT OF A POLYNOMIAL INDIVIDUATED
      *   BY THE VALUES ATTAINED AT THE ZEROES OF THE HERMITE POLYNOMIAL
      *   N   = THE NUMBER OF ZEROES
      *   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
      *   DZ  = SCALED HERMITE DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
      *   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
      *   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
      *   QX  = VALUE OF THE POLYNOMIAL IN X
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1), QZ(1)
      IF (N .EQ. 0) RETURN

      QX = QZ(1)
      IF (N .EQ. 1) RETURN

      EPS = 1.D-14
      CALL VAHESF(N,X,Y,DY)
      QX = 0.DO
      DO 1 J=1,N
        ED = X-CS(J)
        IF(DABS(ED) .LT. EPS) THEN
          QX = QZ(J)
          RETURN
        ELSE
          PR = 1.DO
          DO 2 K=1,N/2
            DK = 4.DO*DFLOAT(K)
            PR = PR*(DK+X*X)/(DK+CS(J)**2)
          2 CONTINUE
          QX = QX+QZ(J)*Y*PR/(DZ(J)*ED)
        ENDIF
      1 CONTINUE

      RETURN
      END

```

INJAGL

For any $n \geq 1$, the routine computes the value at the point x of a polynomial q of degree at most n , determined by the values attained at the Jacobi Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$.

One has

$$q(x) = \sum_{j=0}^n q(\eta_j^{(n)}) \tilde{l}_j^{(n)}(x),$$

where

$$\tilde{l}_j^{(n)}(x) = \begin{cases} \frac{(\beta + 1)(x - 1)}{n(n + \alpha + \beta + 1) P_n^{(\alpha, \beta)}(\eta_0^{(n)})} \frac{d}{dx} P_n^{(\alpha, \beta)}(x) & j = 0, \\ \frac{(x^2 - 1)}{n(n + \alpha + \beta + 1) P_n^{(\alpha, \beta)}(\eta_j^{(n)}) (x - \eta_j^{(n)})} \frac{d}{dx} P_n^{(\alpha, \beta)}(x) & 1 \leq j \leq n - 1, \\ \frac{(\alpha + 1)(x + 1)}{n(n + \alpha + \beta + 1) P_n^{(\alpha, \beta)}(\eta_n^{(n)})} \frac{d}{dx} P_n^{(\alpha, \beta)}(x) & j = n. \end{cases}$$

The nodes and the values $P_n^{(\alpha, \beta)}(\eta_j^{(n)})$, $0 \leq j \leq n$, can be obtained by calling subroutine ZEJAGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n	QX , the value of q in x
A , the parameter α	
B , the parameter β	
ET , vector of the nodes $\eta_j^{(n)}$	
VN , the values $P_n^{(\alpha, \beta)}(\eta_j^{(n)})$	
QN , values of q at the nodes	
X , the point x	

Auxiliary routine: **VAJAPO**

```

      SUBROUTINE INJAGL(N,A,B,ET,VN,QN,X,QX)
      *****
      *   COMPUTES THE VALUE AT A GIVEN POINT OF A POLYNOMIAL INDIVIDUATED
      *   BY THE VALUES ATTAINED AT THE JACOBI GAUSS-LOBATTO NODES
      *   N   = THE DEGREE OF THE POLYNOMIAL
      *   A   = PARAMETER > -1
      *   B   = PARAMETER > -1
      *   ET  = VECTOR OF THE NODES, ET(I), I=0,N
      *   VN  = VALUES OF THE JACOBI POLYNOMIAL AT THE NODES, VN(I), I=0,N
      *   QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      *   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
      *   QX  = VALUE OF THE POLYNOMIAL IN X
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(O:*), VN(O:*), QN(O:*)
      IF (N .EQ. 0) RETURN

      EPS = 1.D-14
      CALL VAJAPO(N,A,B,X,Y,DY,D2Y)
      DN = DFLOAT(N)
      C  = 1.DO/(DN*(DN+A+B+1.DO))
      QX = QN(O)*C*DY*(B+1.DO)*(X-1.DO)/VN(O)
      QX = QX+QN(N)*C*DY*(A+1.DO)*(X+1.DO)/VN(N)
      IF (N .EQ. 1) RETURN

      DO 1 J=1,N-1
        ED = X-ET(J)
        IF(DABS(ED) .LT. EPS) THEN
          QX = QN(J)
          RETURN
        ELSE
          QX = QX+QN(J)*C*DY*(X*X-1.DO)/(VN(J)*ED)
        ENDIF
      1  CONTINUE

      RETURN
      END

```

INLEGL

For any $n \geq 1$, the routine computes the value at the point x of a polynomial q of degree at most n , determined by the values attained at the Legendre Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$.

One has

$$q(x) = \sum_{j=0}^n q(\eta_j^{(n)}) \tilde{l}_j^{(n)}(x),$$

where

$$\tilde{l}_j^{(n)}(x) = \begin{cases} \frac{(x-1) P_n'(x)}{n(n+1) P_n(\eta_0^{(n)})} & j = 0, \\ \frac{(x^2-1) P_n'(x)}{n(n+1) P_n(\eta_j^{(n)}) (x-\eta_j^{(n)})} & 1 \leq j \leq n-1, \\ \frac{(x+1) P_n'(x)}{n(n+1) P_n(\eta_n^{(n)})} & j = n. \end{cases}$$

The nodes and the values $P_n(\eta_j^{(n)})$, $0 \leq j \leq n$, can be obtained by calling subroutine ZELEGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n	QX , the value of q in x
ET , vector of the nodes $\eta_j^{(n)}$	
VN , the values $P_n(\eta_j^{(n)})$	
QN , values of q at the nodes	
X , the point x	

Auxiliary routine: **VALEPO**

```

      SUBROUTINE INLEGL(N,ET,VN,QN,X,QX)
      *****
      *   COMPUTES THE VALUE AT A GIVEN POINT OF A POLYNOMIAL INDIVIDUATED
      *   BY THE VALUES ATTAINED AT THE LEGENDRE GAUSS-LOBATTO NODES
      *   N   = THE DEGREE OF THE POLYNOMIAL
      *   ET  = VECTOR OF THE NODES, ET(I), I=0,N
      *   VN  = VALUES OF THE LEGENDRE POLYNOMIAL AT THE NODES, VN(I), I=0,N
      *   QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      *   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
      *   QX  = VALUE OF THE POLYNOMIAL IN X
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), VN(0:*), QN(0:*)
      IF (N .EQ. 0) RETURN

      EPS = 1.D-14
      CALL VALEPO(N,X,Y,DY,D2Y)
      DN = DFLOAT(N)
      C  = 1.DO/(DN*(DN+1.DO))
      QX = QN(0)*C*DY*(X-1.DO)/VN(0)
      QX = QX+QN(N)*C*DY*(X+1.DO)/VN(N)
      IF (N .EQ. 1) RETURN

      DO 1 J=1,N-1
        ED = X-ET(J)
        IF(DABS(ED) .LT. EPS) THEN
          QX = QN(J)
        RETURN
        ELSE
          QX = QX+QN(J)*C*DY*(X*X-1.DO)/(VN(J)*ED)
        ENDIF
      1 CONTINUE

      RETURN
      END

```

INCHGL

For any $n \geq 1$, the routine computes the value at the point x of a polynomial q of degree at most n , determined by the values attained at the Chebyshev Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$.

One has

$$q(x) = \sum_{j=0}^n q(\eta_j^{(n)}) \tilde{l}_j^{(n)}(x),$$

where

$$\tilde{l}_j^{(n)}(x) = \begin{cases} \frac{(-1)^n}{2n^2} (x-1) T_n'(x) & j = 0, \\ \frac{(-1)^{n+j}}{n^2} \frac{(x^2-1) T_n'(x)}{x - \eta_j^{(n)}} & 1 \leq j \leq n-1, \\ \frac{1}{2n^2} (x+1) T_n'(x) & j = n. \end{cases}$$

<i>Input variables</i>	<i>Output variable</i>
N , the degree n QN , values of q at the nodes X , the point x	QX , the value of q in x

Auxiliary routine: **VACHPO**

```

      SUBROUTINE INCHGL(N,QN,X,QX)
      *****
      *   COMPUTES THE VALUE AT A GIVEN POINT OF A POLYNOMIAL INDIVIDUATED
      *   BY THE VALUES ATTAINED AT THE CHEBYSHEV GAUSS-LOBATTO NODES
      *   N   = THE DEGREE OF THE POLYNOMIAL
      *   QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      *   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
      *   QX  = VALUE OF THE POLYNOMIAL IN X
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION QN(0:*)
      IF (N .EQ. 0) RETURN

      EPS = 1.D-14
      PI = 3.14159265358979323846D0
      CALL VACHPO(N,X,Y,DY,D2Y)
      DN = DFLOAT(N)
      SN = DFLOAT(1+4*(N/2)-2*N)
      PN = PI/DN
      C = 1.DO/(DN*DN)
      QX = .5DO*SN*QN(0)*C*DY*(X-1.DO)
      QX = QX+.5DO*QN(N)*C*DY*(X+1.DO)
      IF (N .EQ. 1) RETURN

      SJ = -1.DO
      DO 1 J=1,N-1
        ED = X+DCOS(PN*DFLOAT(J))
        IF(DABS(ED) .LT. EPS) THEN
          QX = QN(J)
        RETURN
      ELSE
        QX = QX+QN(J)*SN*SJ*C*DY*(X*X-1.DO)/ED
      ENDIF
      SJ = -SJ
1    CONTINUE

      RETURN
      END

```

INLAGR

For any $n \geq 1$, the routine computes the value at the point x of a polynomial q of degree at most $n - 1$, determined by the values attained at the Laguerre Gauss-Radau nodes $\eta_j^{(n)}$, $0 \leq j \leq n - 1$.

One has

$$q(x) = \sum_{j=0}^{n-1} q(\eta_j^{(n)}) \tilde{l}_j^{(n)}(x),$$

where the polynomial $\tilde{l}_j^{(n)}$ is given by

$$-\frac{\alpha + 1}{n} [\hat{L}_n^{(\alpha)}(\eta_0^{(n)})]^{-1} \left[\frac{d}{dx} \hat{L}_n^{(\alpha)}(x) + \left(\sum_{m=1}^n \frac{1}{4m + x} \right) \hat{L}_n^{(\alpha)}(x) \right] \prod_{k=1}^n \frac{4k + x}{4k} \quad \text{if } j = 0,$$

$$-\frac{x}{n(x - \eta_j^{(n)})} [\hat{L}_n^{(\alpha)}(\eta_j^{(n)})]^{-1} \left[\frac{d}{dx} \hat{L}_n^{(\alpha)}(x) + \left(\sum_{m=1}^n \frac{1}{4m + x} \right) \hat{L}_n^{(\alpha)}(x) \right] \prod_{k=1}^n \frac{4k + x}{4k + \eta_j^{(n)}} \quad \text{if } 1 \leq j \leq n - 1.$$

Here, $\hat{L}_n^{(\alpha)}$ is the scaled Laguerre function of degree n , introduced in the description of subroutine VALASF. The nodes and the values $\hat{L}_n^{(\alpha)}(\eta_j^{(n)})$, $0 \leq j \leq n$, can be obtained by calling subroutine ZELAGR.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n	QX , the value of q in x
A , the parameter α	
ET , vector of the nodes $\eta_j^{(n)}$	
VN , the values $\hat{L}_n^{(\alpha)}(\eta_j^{(n)})$	
QN , values of q at the nodes	
X , the point x	

Auxiliary routine: **VALASF**


```

      SUBROUTINE INLAGR(N,A,ET,VN,QN,X,QX)
      *****
      *   COMPUTES THE VALUE AT A GIVEN POINT OF A POLYNOMIAL INDIVIDUATED
      *   BY THE VALUES ATTAINED AT THE LAGUERRE GAUSS-RADAU NODES
      *   N   = THE NUMBER OF NODES
      *   A   = PARAMETER > -1
      *   ET  = VECTOR OF THE NODES, ET(I), I=0,N-1
      *   VN  = SCALED LAGUERRE FUNCTION AT THE NODES, VN(I), I=0,N-1
      *   QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N-1
      *   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
      *   QX  = VALUE OF THE POLYNOMIAL IN X
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), VN(0:*), QN(0:*)
      IF (N .EQ. 0) RETURN
      QX = QN(0)
      IF (N .EQ. 1) RETURN
      EPS = 1.D-14
      DN = DFLOAT(N)
      C  = -1.DO/DN
      PR = 1.DO
      SU = 0.DO
      DO 1 M=1,N
      DM = 4.DO*DFLOAT(M)
      PR = PR*(DM+X)/DM
      SU = SU+1.DO/(DM+X)
1    CONTINUE
      CALL VALASF(N,A,X,Y,DY)
      QX = QN(0)*C*(A+1.DO)*PR*(DY+SU*Y)/VN(0)
      DO 2 J=1,N-1
      ED = X-ET(J)
      IF(DABS(ED) .LT. EPS) THEN
      QX = QN(J)
      RETURN
      ELSE
      PR = 1.DO
      DO 3 K=1,N
      DK = 4.DO*DFLOAT(K)
      PR = PR*(DK+X)/(DK+ET(J))
3    CONTINUE
      QX = QX+QN(J)*C*PR*(DY+SU*Y)*X/(VN(J)*ED)
      ENDIF
2    CONTINUE
      RETURN
      END

```

NOLEGA

For any $n \geq 1$, the routine computes the quantities

$$\mathcal{I} = \left(\int_{-1}^1 q^2(x) dx \right)^{\frac{1}{2}} = \left(\sum_{j=1}^n q^2(\xi_j^{(n)}) w_j^{(n)} \right)^{\frac{1}{2}},$$

$$\mathcal{M} = \max_{1 \leq j \leq n} |q(\xi_j^{(n)})|,$$

where q is a polynomial of degree at most $n - 1$ individuated by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the Legendre polynomial P_n .

Here, $w_j^{(n)}$, $1 \leq j \leq n$, are the weights of the Legendre Gauss integration formula, which can be determined calling subroutine WELEGA.

<i>Input variables</i>	<i>Output variables</i>
N , the number of zeroes	QI , the norm \mathcal{I}
QZ , the values of q at the zeroes	QM , the norm \mathcal{M}
WE , Gauss-Legendre weights $w_j^{(n)}$	

```
      SUBROUTINE NOLEGA(N,QZ,WE,QI,QM)
*****
*   COMPUTES THE NORMS OF A POLYNOMIAL DEFINED AT THE LEGENDRE ZEROES
*   N   = THE NUMBER OF ZEROES
*   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
*   WE  = VECTOR OF THE LEGENDRE GAUSS WEIGHTS, WE(I), I=1,N
*   QI  = INTEGRAL NORM OF THE POLYNOMIAL
*   QM  = MAXIMUM VALUE OF THE POLYNOMIAL AT THE ZEROES
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION QZ(1), WE(1)
      IF (N .EQ. 0) RETURN

      EPS = 1.D-14
      SU  = 0.DO
      QM  = 0.DO
      DO 1 J=1,N
          Y = DABS(QZ(J))
          IF(Y .GT. QM) QM=Y
          IF(Y .LT. EPS) GOTO 1
          SU = SU+Y*Y*WE(J)
1      CONTINUE
          QI = DSQRT(SU)

      RETURN
      END
```

NOCHGA

For any $n \geq 1$, the routine computes the quantities

$$\mathcal{W} = \left(\int_{-1}^1 q^2(x) \frac{dx}{\sqrt{1-x^2}} \right)^{\frac{1}{2}} = \left(\frac{\pi}{n} \sum_{j=1}^n q^2(\xi_j^{(n)}) \right)^{\frac{1}{2}},$$

$$\mathcal{I} = \left(\int_{-1}^1 q^2(x) dx \right)^{\frac{1}{2}} = \left(\sum_{j=0}^{2n} q^2 \left(\cos \frac{j\pi}{2n} \right) \chi_j^{(2n)} \right)^{\frac{1}{2}},$$

$$\mathcal{M} = \max_{1 \leq j \leq n} |q(\xi_j^{(n)})|,$$

where q is a polynomial of degree at most $n-1$ individuated by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the Chebyshev polynomial T_n .

Here, $\chi_j^{(2n)}$, $0 \leq j \leq 2n$, are the weights of the Clenshaw-Curtis integration formula, which can be determined by calling subroutine WECHCC. The zeroes and the values $T_n'(\xi_j^{(n)})$, $1 \leq j \leq n$, can be obtained by calling subroutine ZECHGA.

<i>Input variables</i>	<i>Output variables</i>
N , the number of zeroes	QW , the norm \mathcal{W}
DZ , the values $T_n'(\xi_j^{(n)})$	QI , the norm \mathcal{I}
QZ , the values of q at the zeroes	QM , the norm \mathcal{M}
WK , Clenshaw-Curtis weights $\chi_j^{(2n)}$	

Auxiliary routines: **INCHGA**, **VACHPO**

```

      SUBROUTINE NOCHGA(N,DZ,QZ,WK,QW,QI,QM)
      *****
      *   COMPUTES THE NORMS OF A POLYNOMIAL DEFINED AT THE CHEBYSHEV ZEROES
      *   N   = THE NUMBER OF ZEROES
      *   DZ  = CHEBYSHEV POLYNOMIAL DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
      *   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
      *   WK  = VECTOR OF THE CLENSHAW-CURTIS WEIGHTS, WE(I), I=0,2*N
      *   QW  = WEIGHTED INTEGRAL NORM OF THE POLYNOMIAL
      *   QI  = INTEGRAL NORM OF THE POLYNOMIAL
      *   QM  = MAXIMUM VALUE OF THE POLYNOMIAL AT THE ZEROES
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION DZ(1), QZ(1), WK(0:*)
      IF (N .EQ. 0) RETURN

      PI = 3.14159265358979323846D0
      DN = DFLOAT(N)
      X  = -1.DO
      CALL INCHGA(N,DZ,QZ,X,QX)
      S1 = 0.DO
      S2 = QX*QX*WK(0)
      QM = 0.DO
      DO 1 J=1,N
      DJ = DFLOAT(J)
      J2 = 2*J
      X  = -DCOS(PI*DJ/DN)
      Y  = DABS(QZ(J))
      IF(Y .GT. QM) QM=Y
      CALL INCHGA(N,DZ,QZ,X,QX)
      S1 = S1+Y*Y
      S2 = S2+Y*Y*WK(J2-1)+QX*QX*WK(J2)
1     CONTINUE
      QW = DSQRT(S1*PI/DN)
      QI = DSQRT(S2)

      RETURN
      END

```

NOJAGL

For any $n \geq 1$, the routine computes the quantities

$$\mathcal{W} = \left(\int_{-1}^1 q^2(x) (1-x)^\alpha (1+x)^\beta dx \right)^{\frac{1}{2}},$$

$$\mathcal{S} = \left(\sum_{j=0}^n q^2(\eta_j^{(n)}) \tilde{w}_j^{(n)} \right)^{\frac{1}{2}},$$

$$\mathcal{M} = \max_{0 \leq j \leq n} |q(\eta_j^{(n)})|,$$

where q is a polynomial of degree at most n individuated by the values attained at the Jacobi Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$.

Here, $\tilde{w}_j^{(n)}$, $0 \leq j \leq n$, are the weights of the Jacobi Gauss-Lobatto formula, which can be determined by calling subroutine WEJAGL.

The nodes and the values $P_n^{(\alpha, \beta)}(\eta_j^{(n)})$, $0 \leq j \leq n$, can be obtained by calling subroutine ZEJAGL.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	QW , the norm \mathcal{W}
A , the parameter α	QS , the norm \mathcal{S}
B , the parameter β	QM , the norm \mathcal{M}
VN , the values $P_n^{(\alpha, \beta)}(\eta_j^{(n)})$	
QN , the values of q at the nodes	
WT , the weights $\tilde{w}_j^{(n)}$	

Auxiliary routine: **GAMMAF**

```

      SUBROUTINE NOJAGL(N,A,B,VN,QN,WT,QW,QS,QM)
*****
*   COMPUTES THE NORMS OF A POLYNOMIAL DEFINED AT THE JACOBI GAUSS-
*   LOBATTO NODES
*   N = THE DEGREE OF THE POLYNOMIAL
*   A = PARAMETER > -1
*   B = PARAMETER > -1
*   VN = VALUES OF THE JACOBI POLYNOMIAL AT THE NODES, VN(I), I=0,N
*   QN = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
*   WT = VECTOR OF THE JACOBI GAUSS-LOBATTO WEIGHTS, WT(I), I=0,N
*   QW = WEIGHTED INTEGRAL NORM OF THE POLYNOMIAL
*   QS = QUADRATURE NORM OF THE POLYNOMIAL
*   QM = MAXIMUM VALUE OF THE POLYNOMIAL AT THE NODES
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION VN(0:*), QN(0:*), WT(0:*)
      IF (N .EQ. 0) RETURN

      EPS = 1.D-14
      A1 = A+1.DO
      B1 = B+1.DO
      AB = A+B
      AB2 = AB+2.DO
      DN = DFLOAT(N)
      C = ((2.DO)**(AB+1.DO))*(DN+AB+1.DO)/(2.DO*DN+AB+1.DO)
      CALL GAMMAF(A1,GA1)
      CALL GAMMAF(B1,GB1)
      CALL GAMMAF(AB2,GAB2)
      C = C*GA1*GB1/GAB2
      DO 1 K=1,N
         DK = DFLOAT(K)
         C = C*(DK+A)*(DK+B)/(DK*(DK+AB+1.DO))
1      CONTINUE

      S1 = 0.DO
      S2 = 0.DO
      S3 = 0.DO
      QM = 0.DO
      DO 2 J=0,N
         Y1 = QN(J)
         YM = DABS(Y1)
         IF(YM .GT. QM) QM=YM
         Y2 = VN(J)
         S2 = S2+Y1*Y2*WT(J)
         S3 = S3+Y2*Y2*WT(J)
         IF(YM .LT. EPS) GOTO 2
         S1 = S1+Y1*Y1*WT(J)
2      CONTINUE
      QS = DSQRT(S1)
      QW = DSQRT(DABS(S1-(S3-C)*S2*S2/(S3*S3)))

      RETURN
      END

```

NOLEGL

For any $n \geq 1$, the routine computes the quantities

$$\mathcal{I} = \left(\int_{-1}^1 q^2(x) dx \right)^{\frac{1}{2}} = \left[\mathcal{S}^2 - \frac{n(n+1)}{2(2n+1)} \left(\sum_{j=0}^n (qP_n)(\eta_j^{(n)}) \tilde{w}_j^{(n)} \right)^2 \right]^{\frac{1}{2}},$$

$$\mathcal{S} = \left(\sum_{j=0}^n q^2(\eta_j^{(n)}) \tilde{w}_j^{(n)} \right)^{\frac{1}{2}},$$

$$\mathcal{M} = \max_{0 \leq j \leq n} |q(\eta_j^{(n)})|,$$

where q is a polynomial of degree at most n individuated by the values attained at the Legendre Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$.

Here, $\tilde{w}_j^{(n)}$, $0 \leq j \leq n$, are the weights of the Legendre Gauss-Lobatto formula, which can be determined by calling subroutine WELEGL. The nodes and the values $P_n(\eta_j^{(n)})$, $0 \leq j \leq n$, can be obtained by calling subroutine ZELEGL.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	QI , the norm \mathcal{I}
VN , the values $P_n(\eta_j^{(n)})$	QS , the norm \mathcal{S}
QN , the values of q at the nodes	QM , the norm \mathcal{M}
WT , the weights $\tilde{w}_j^{(n)}$	


```

      SUBROUTINE NOLEGL(N,VN,QN,WT,QI,QS,QM)
      *****
      *   COMPUTES THE NORMS OF A POLYNOMIAL DEFINED AT THE LEGENDRE GAUSS-
      *   LOBATTO NODES
      *   N   = THE DEGREE OF THE POLYNOMIAL
      *   VN  = VALUES OF THE LEGENDRE POLYNOMIAL AT THE NODES, VN(I), I=0,N
      *   QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      *   WT  = VECTOR OF THE LEGENDRE GAUSS-LOBATTO WEIGHTS, WT(I), I=0,N
      *   QW  = INTEGRAL NORM OF THE POLYNOMIAL
      *   QS  = QUADRATURE NORM OF THE POLYNOMIAL
      *   QM  = MAXIMUM VALUE OF THE POLYNOMIAL AT THE NODES
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION VN(O:*), QN(O:*), WT(O:*)
      IF (N .EQ. 0) RETURN

      EPS = 1.D-14
      DN = DFLOAT(N)
      C   = .5DO*DN*(DN+1.DO)/(2.DO*DN+1.DO)

      S1 = 0.DO
      S2 = 0.DO
      QM = 0.DO
      DO 1 J=0,N
         Y1 = QN(J)
         YM = DABS(Y1)
         IF(YM .GT. QM) QM=YM
         Y2 = VN(J)
         S2 = S2+Y1*Y2*WT(J)
         IF(YM .LT. EPS) GOTO 1
         S1 = S1+Y1*Y1*WT(J)
1      CONTINUE
         QS = DSQRT(S1)
         QI = DSQRT(DABS(S1-C*S2*S2))

      RETURN
      END

```

NOCHGL

For any $n \geq 1$, the routine computes the quantities

$$\mathcal{W} = \left(\int_{-1}^1 q^2(x) \frac{dx}{\sqrt{1-x^2}} \right)^{\frac{1}{2}} = \left[\mathcal{S}^2 - \frac{1}{2\pi} \left(\sum_{j=0}^n (-1)^{n+j} q(\eta_j^{(n)}) \tilde{w}_j^{(n)} \right)^2 \right]^{\frac{1}{2}},$$

$$\mathcal{I} = \left(\int_{-1}^1 q^2(x) dx \right)^{\frac{1}{2}} = \left[\sum_{j=0}^{2n} q^2 \left(\cos \frac{j\pi}{2n} \right) \chi_j^{(2n)} \right]^{\frac{1}{2}},$$

$$\mathcal{S} = \left(\sum_{j=0}^n q^2(\eta_j^{(n)}) \tilde{w}_j^{(n)} \right)^{\frac{1}{2}}, \quad \mathcal{M} = \max_{0 \leq j \leq n} |q(\eta_j^{(n)})|,$$

where q is a polynomial of degree at most n individuated by the values attained at the Chebyshev Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$.

Here, $\tilde{w}_j^{(n)}$, $0 \leq j \leq n$, are the weights of the Chebyshev Gauss-Lobatto formula and $\chi_j^{(2n)}$, $0 \leq j \leq 2n$, are the Clenshaw-Curtis weights. The first can be determined by calling subroutine WECHGL, the last by calling subroutine WECHCC.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	QW , the norm \mathcal{W}
QN , the values of q at the nodes	QI , the norm \mathcal{I}
WK , Clenshaw-Curtis weights $\chi_j^{(2n)}$	QS , the norm \mathcal{S}
	QM , the norm \mathcal{M}

Auxiliary routines: INCHGL, VACHPO

SUBROUTINE NOCHGL(N, QN, WK, QW, QI, QS, QM)

 * COMPUTES THE NORMS OF A POLYNOMIAL DEFINED AT THE CHEBYSHEV GAUSS-
 * LOBATTO NODES

```

*   N = THE DEGREE OF THE POLYNOMIAL
*   QN = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
*   WK = VECTOR OF THE CLENSHAW-CURTIS WEIGHTS, WE(I), I=0,2*N
*   QW = WEIGHTED INTEGRAL NORM OF THE POLYNOMIAL
*   QI = INTEGRAL NORM OF THE POLYNOMIAL
*   QS = QUADRATURE NORM OF THE POLYNOMIAL
*   QM = MAXIMUM VALUE OF THE POLYNOMIAL AT THE NODES
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION QN(0:*), WK(0:*)
      IF (N .EQ. 0) RETURN

      PI = 3.14159265358979323846D0
      PH = 1.57079632679489661923D0
      DN = DFLOAT(N)
      SN = DFLOAT(1+4*(N/2)-2*N)
      Y = QN(0)
      S1 = .5D0*Y*Y
      S2 = .5D0*Y*SN
      S3 = Y*Y*WK(0)
      QM = DABS(Y)

      SJ = -1.D0
      DO 1 J=1,N-1
        J2 = 2*J
        DJ = DFLOAT(J2-1)
        X = -DCOS(PH*DJ/DN)
        Y = QN(J)
        YM = DABS(Y)
        IF(YM .GT. QM) QM=YM
        CALL INCHGL(N,QN,X,QX)
        S1 = S1+Y*Y
        S2 = S2+Y*SN*SJ
        S3 = S3+QX*QX*WK(J2-1)+Y*Y*WK(J2)
        SJ = -SJ
1     CONTINUE
        N2 = 2*N
        DD = DFLOAT(N2-1)
        X = -DCOS(PH*DD/DN)
        Y = QN(N)
        YM = DABS(Y)
        IF(YM .GT. QM) QM=YM
        CALL INCHGL(N,QN,X,QX)
        S1 = S1+.5D0*Y*Y
        S2 = S2+.5D0*Y
        S3 = S3+QX*QX*WK(N2-1)+Y*Y*WK(N2)

      QW = DSQRT(DABS(PI*S1/DN-PH*S2*S2/(DN*DN)))
      QI = DSQRT(S3)
      QS = DSQRT(PI*S1/DN)

      RETURN
      END

```

COJAGA

For any $n \geq 1$, the routine computes the n Fourier coefficients c_k , $0 \leq k \leq n-1$, with respect to the Jacobi polynomial basis, of a polynomial q of degree at most $n-1$ determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of $P_n^{(\alpha,\beta)}$.

One has

$$\begin{aligned} c_k &= \gamma_k \int_{-1}^1 q(x) P_k^{(\alpha,\beta)}(x) (1-x)^\alpha (1+x)^\beta dx \\ &= \gamma_k \sum_{j=1}^n q(\xi_j^{(n)}) P_k^{(\alpha,\beta)}(\xi_j^{(n)}) w_j^{(n)} \quad 0 \leq k \leq n-1, \end{aligned}$$

where $\gamma_k = \left(\int_{-1}^1 [P_k^{(\alpha,\beta)}(x)]^2 (1-x)^\alpha (1+x)^\beta dx \right)^{-1}$, $0 \leq k \leq n-1$.

Here, $w_j^{(n)}$, $1 \leq j \leq n$, are the weights of the Jacobi Gauss integration formula, which can be determined by calling subroutine WEJAGA. The zeroes can be obtained by calling subroutine ZEJAGA.

<i>Input variables</i>	<i>Output variable</i>
N , the number of zeroes A , the parameter α B , the parameter β CS , vector of the zeroes $\xi_j^{(n)}$ QZ , the values of q at the zeroes WE , the weights $w_j^{(n)}$	CO , the Fourier coefficients of q

Auxiliary routine: **GAMMAF**

SUBROUTINE COJAGA(N,A,B,CS,QZ,WE,CO)

* COMPUTES THE JACOBI FOURIER COEFFICIENTS OF A POLYNOMIAL
 * INDIVIDUATED BY THE VALUES ATTAINED AT THE JACOBI ZEROES
 * N = THE NUMBER OF ZEROES
 * A = PARAMETER >-1
 * B = PARAMETER >-1

```

*   CS = VECTOR OF THE ZEROES, CS(I), I=1,N
*   QZ = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
*   WE = VECTOR OF THE WEIGHTS, WE(I), I=1,N
*   CO = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N-1
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), QZ(1), WE(1), CO(0:*)
      IF (N .EQ. 0) RETURN

      A1 = A+1.DO
      B1 = B+1.DO
      AB = A+B
      AB2 = AB+2.DO
      CALL GAMMAF(A1,GA1)
      CALL GAMMAF(B1,GB1)
      CALL GAMMAF(AB2,GAB2)
      C = ((2.DO)**(AB+1.DO))*GA1*GB1/GAB2

      SU = 0.DO
      DO 1 J=1,N
        SU = SU+QZ(J)*WE(J)
        CO(J-1) = 0.DO
1     CONTINUE
      CO(0) = SU/C
      IF (N .EQ. 1) RETURN

      DO 2 J=1,N
        X = CS(J)
        YP = QZ(J)*WE(J)
        Y = .5DO*YP*(AB2*X+A-B)
      DO 3 K=1,N-1
        CO(K) = CO(K)+Y
        DK = DFLOAT(K+1)
        CC = 2.DO*DK+AB
        C1 = 2.DO*DK*(DK+AB)*(CC-2.DO)
        C2 = (CC-1.DO)*(CC-2.DO)*CC
        C3 = (CC-1.DO)*(A-B)*AB
        C4 = 2.DO*(DK+A-1.DO)*CC*(DK+B-1.DO)
        YM = Y
        Y = ((C2*X+C3)*Y-C4*YP)/C1
        YP = YM
3     CONTINUE
2     CONTINUE

      DO 4 K=1,N-1
        DK = DFLOAT(K)
        C = C*(DK+A)*(DK+B)/DK
        CO(K) = CO(K)*(2.DO*DK+AB+1.DO)/C
        C = C/(DK+AB+1.DO)
4     CONTINUE

      RETURN
      END

```

COLEGA

For any $n \geq 1$, the routine computes the n Fourier coefficients c_k , $0 \leq k \leq n - 1$, with respect to the Legendre polynomial basis, of a polynomial q of degree at most $n - 1$ determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of P_n .

One has

$$c_k = \frac{2k+1}{2} \int_{-1}^1 q(x) P_k(x) dx = \frac{2k+1}{2} \sum_{j=1}^n q(\xi_j^{(n)}) P_k(\xi_j^{(n)}) w_j^{(n)} \quad 0 \leq k \leq n-1.$$

Here, $w_j^{(n)}$, $1 \leq j \leq n$, are the weights of the Legendre Gauss integration formula, which can be determined by calling subroutine WELEGA.

The zeroes can be obtained by calling subroutine ZELEGA.

<i>Input variables</i>	<i>Output variable</i>
N , the number of zeroes CS , vector of the zeroes $\xi_j^{(n)}$ QZ , the values of q at the zeroes WE , the weights $w_j^{(n)}$	CO , the Fourier coefficients of q

```

      SUBROUTINE COLEGA(N,CS,QZ,WE,CO)
*****
*   COMPUTES THE LEGENDRE FOURIER COEFFICIENTS OF A POLYNOMIAL
*   INDIVIDUATED BY THE VALUES ATTAINED AT THE LEGENDRE ZEROES
*   N   = THE NUMBER OF ZEROES
*   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
*   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
*   WE  = VECTOR OF THE WEIGHTS, WE(I), I=1,N
*   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N-1
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), QZ(1), WE(1), CO(0:*)
      IF (N .EQ. 0) RETURN

      SU = 0.DO
      DO 1 J=1,N
         SU = SU+QZ(J)*WE(J)
         CO(J-1) = 0.DO
1      CONTINUE
         CO(0) = .5DO*SU
      IF (N .EQ. 1) RETURN

      DO 2 J=1,N
         X = CS(J)
         YP = QZ(J)*WE(J)
         Y = X*YP
      DO 3 K=1,N-1
         CO(K) = CO(K)+Y
         DK = DFLOAT(K+1)
         C1 = 2.DO*DK-1.DO
         C2 = DK-1.DO
         YM = Y
         Y = (C1*X*Y-C2*YP)/DK
         YP = YM
3      CONTINUE
2      CONTINUE

      DO 4 K=1,N-1
         CO(K) = .5DO*CO(K)*(2.DO*DFLOAT(K)+1.DO)
4      CONTINUE

      RETURN
      END

```

COCHGA

For any $n \geq 1$, the routine computes the n Fourier coefficients c_k , $0 \leq k \leq n - 1$, with respect to the Chebyshev polynomial basis, of a polynomial q of degree at most $n - 1$ determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of T_n .

One has for $0 \leq k \leq n - 1$

$$c_k = \gamma_k \int_{-1}^1 q(x) T_k(x) \frac{dx}{\sqrt{1-x^2}} = \frac{(-1)^k \gamma_k \pi}{n} \sum_{j=1}^n q(\xi_j^{(n)}) \cos \frac{k(2j-1)\pi}{2n},$$

where $\gamma_0 = 1/\pi$ and $\gamma_k = 2/\pi$, $1 \leq k \leq n - 1$.

<i>Input variables</i>	<i>Output variable</i>
N , the number of zeroes QZ , the values of q at the zeroes	CO , the Fourier coefficients of q


```

      SUBROUTINE COCHGA(N,QZ,CO)
      *****
      *   COMPUTES THE CHEBYSHEV FOURIER COEFFICIENTS OF A POLYNOMIAL
      *   INDIVIDUATED BY THE VALUES ATTAINED AT THE CHEBYSHEV ZEROES
      *   N   = THE NUMBER OF ZEROES
      *   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
      *   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N-1
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION  QZ(1), CO(0:*)
      IF (N .EQ. 0) RETURN

      PH = 1.57079632679489661923D0
      DN = DFLOAT(N)
      SU = 0.DO
      DO 1 J=1,N
      SU = SU+QZ(J)
1     CONTINUE
      CO(0) = SU/DN
      IF (N .EQ. 1) RETURN

      SK = -2.DO
      DO 2 K=1,N-1
      DK = DFLOAT(K)
      SU = 0.DO
      DO 3 J=1,N
      DJ = 2.DO*DFLOAT(J)-1.DO
      SU = SU+QZ(J)*DCOS(DK*DJ*PH/DN)
3     CONTINUE
      CO(K) = SK*SU/DN
      SK = -SK
2     CONTINUE

      RETURN
      END

```

COLAGA

For any $n \geq 1$, the routine computes the n Fourier coefficients c_k , $0 \leq k \leq n-1$, with respect to the Laguerre polynomial basis, of a polynomial q of degree at most $n-1$ determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of $L_n^{(\alpha)}$.

One has

$$c_k = \frac{k!}{\Gamma(k + \alpha + 1)} \int_0^{+\infty} q(x) L_k^{(\alpha)}(x) x^\alpha e^{-x} dx$$

$$= \frac{k!}{\Gamma(k + \alpha + 1)} \sum_{j=1}^n q(\xi_j^{(n)}) L_k^{(\alpha)}(\xi_j^{(n)}) w_j^{(n)} \quad 0 \leq k \leq n-1.$$

Here, $w_j^{(n)}$, $1 \leq j \leq n$, are the weights of the Laguerre Gauss integration formula, which can be determined by calling subroutine WELAGA.

The zeroes can be obtained by calling subroutine ZELAGA.

<i>Input variables</i>	<i>Output variable</i>
N , the number of zeroes	CO , the Fourier coefficients of q
A , the parameter α	
CS , vector of the zeroes $\xi_j^{(n)}$	
QZ , the values of q at the zeroes	
WE , the weights $w_j^{(n)}$	

Auxiliary routine: **GAMMAF**

```

      SUBROUTINE COLAGA(N,A,CS,QZ,WE,CO)
*****
*   COMPUTES THE LAGUERRE FOURIER COEFFICIENTS OF A POLYNOMIAL
*   INDIVIDUATED BY THE VALUES ATTAINED AT THE LAGUERRE ZEROES
*   N   = THE NUMBER OF ZEROES
*   A   = PARAMETER >-1
*   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
*   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
*   WE  = VECTOR OF THE WEIGHTS, WE(I), I=1,N
*   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N-1
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), QZ(1), WE(1), CO(0:*)
      IF (N .EQ. 0) RETURN

      A1 = A+1.DO
      CALL GAMMAF(A1,C)
      SU = 0.DO
      DO 1 J=1,N
          SU = SU+QZ(J)*WE(J)
          CO(J-1) = 0.DO
1      CONTINUE
          CO(0) = SU/C
      IF (N .EQ. 1) RETURN

      DO 2 J=1,N
          X = CS(J)
          YP = QZ(J)*WE(J)
          Y = (A1-X)*YP
      DO 3 K=1,N-1
          CO(K) = CO(K)+Y
          DK = DFLOAT(K+1)
          B1 = (2.DO*DK+A-1.DO-X)/DK
          B2 = (DK+A-1.DO)/DK
          YM = Y
          Y = B1*Y-B2*YP
          YP = YM
3      CONTINUE
2      CONTINUE
      DO 4 K=1,N-1
          DK = DFLOAT(K)
          C = C*(DK+A)/DK
          CO(K) = CO(K)/C
4      CONTINUE

      RETURN
      END

```

COHEGA

For any $n \geq 1$, the routine computes the n Fourier coefficients c_k , $0 \leq k \leq n-1$, with respect to the Hermite polynomial basis, of a polynomial q of degree at most $n-1$ determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of H_n .

One has for $0 \leq k \leq n-1$

$$c_k = \frac{1}{2^k k! \sqrt{\pi}} \int_{-\infty}^{+\infty} q(x) H_k(x) e^{-x^2} dx$$

$$= \frac{1}{2^k k! \sqrt{\pi}} \sum_{j=1}^n q(\xi_j^{(n)}) H_k(\xi_j^{(n)}) w_j^{(n)} = \frac{1}{\sqrt{\pi}} \sum_{j=1}^n q(\xi_j^{(n)}) \tilde{H}_k(\xi_j^{(n)}) w_j^{(n)},$$

where $\tilde{H}_k = (2^k k!)^{-1} H_k$, $k \in \mathbf{N}$, is given by the recursion formula

$$\begin{cases} \tilde{H}_0(x) = 1 \\ \tilde{H}_1(x) = x \\ \tilde{H}_k(x) = \frac{x}{k} \tilde{H}_{k-1}(x) - \frac{1}{2k} \tilde{H}_{k-2}(x), \quad k \geq 2. \end{cases}$$

Here, $w_j^{(n)}$, $1 \leq j \leq n$, are the weights of the Hermite Gauss integration formula, which can be determined by calling subroutine WEHEGA.

The zeroes can be obtained by calling subroutine ZEHEGA.

<i>Input variables</i>	<i>Output variable</i>
N , the number of zeroes	CO , the Fourier coefficients of q
CS , vector of the zeroes $\xi_j^{(n)}$	
QZ , the values of q at the zeroes	
WE , the weights $w_j^{(n)}$	

```

      SUBROUTINE COHEGA(N,CS,QZ,WE,CO)
*****
*   COMPUTES THE HERMITE FOURIER COEFFICIENTS OF A POLYNOMIAL
*   INDIVIDUATED BY THE VALUES ATTAINED AT THE HERMITE ZEROES
*   N   = THE NUMBER OF ZEROES
*   CS  = VECTOR OF THE ZEROES, CS(I), I=1,N
*   QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
*   WE  = VECTOR OF THE WEIGHTS, WE(I), I=1,N
*   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N-1
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), QZ(1), WE(1), CO(0:*)
      IF (N .EQ. 0) RETURN

      PR = 1.77245385090551588D0
      SU = 0.DO
      DO 1 J=1,N
         SU = SU+QZ(J)*WE(J)
         CO(J-1) = 0.DO
1      CONTINUE
         CO(0) = SU/PR
      IF (N .EQ. 1) RETURN

      DO 2 J=1,N
         X = CS(J)
         YP = QZ(J)*WE(J)/PR
         Y = X*YP
      DO 3 K=1,N-1
         CO(K) = CO(K)+Y
         DK = DFLOAT(K+1)
         YM = Y
         Y = (X*Y-.5D0*YP)/DK
         YP = YM
3      CONTINUE
2      CONTINUE

      RETURN
      END

```

COJAGL

For any $n \in \mathbf{N}$, the routine computes the $n + 1$ Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Jacobi polynomial basis, of a polynomial q of degree at most n determined by the values attained at the Jacobi Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$.

One has

$$\begin{aligned} c_k &= \gamma_k \int_{-1}^1 q(x) P_k^{(\alpha, \beta)}(x) (1-x)^\alpha (1+x)^\beta dx \\ &= \delta_k \sum_{j=0}^n q(\eta_j^{(n)}) P_k^{(\alpha, \beta)}(\eta_j^{(n)}) \tilde{w}_j^{(n)} \quad 0 \leq k \leq n, \end{aligned}$$

where $\gamma_k = \left(\int_{-1}^1 [P_k^{(\alpha, \beta)}(x)]^2 (1-x)^\alpha (1+x)^\beta dx \right)^{-1}$, $0 \leq k \leq n$, and $\delta_k = \gamma_k$ if $0 \leq k \leq n-1$, while $\delta_n = \left(\sum_{j=0}^n [P_n^{(\alpha, \beta)}(\eta_j^{(n)})]^2 \tilde{w}_j^{(n)} \right)^{-1}$.

Here, $\tilde{w}_j^{(n)}$, $0 \leq j \leq n$, are the weights of the Jacobi Gauss-Lobatto formula, which can be determined by calling subroutine WEJAGL. The nodes and the values $P_n^{(\alpha, \beta)}(\eta_j^{(n)})$, $0 \leq j \leq n$, can be obtained by calling subroutine ZEJAGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n A , the parameter α B , the parameter β ET , vector of the nodes $\eta_j^{(n)}$ VN , the values $P_n^{(\alpha, \beta)}(\eta_j^{(n)})$ QN , the values of q at the nodes WT , the weights $\tilde{w}_j^{(n)}$	CO , the Fourier coefficients of q

Auxiliary routine: **GAMMAF**

```

      SUBROUTINE COJAGL(N,A,B,ET,VN,QN,WT,CO)
      *****
      *   COMPUTES THE JACOBI FOURIER COEFFICIENTS OF A POLYNOMIAL
      *   INDIVIDUATED BY ITS VALUES AT THE JACOBI GAUSS-LOBATTO NODES
      *   N   = THE DEGREE OF THE POLYNOMIAL
      *   A   = PARAMETER >-1
      *   B   = PARAMETER >-1
      *   ET  = VECTOR OF THE NODES, ET(I), I=0,N
      *   VN  = VALUES OF THE JACOBI POLYNOMIAL AT THE NODES, VN(I), I=0,N
      *   QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      *   WT  = VECTOR OF THE WEIGHTS, WT(I), I=0,N
      *   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(O:*), VN(O:*), QN(O:*), WT(O:*), CO(O:*)
      CO(O) = QN(O)
      IF (N .EQ. 0) RETURN

      A1 = A+1.DO
      B1 = B+1.DO
      AB = A+B
      AB2 = AB+2.DO
      CO(1) = (QN(1)-QN(O))/AB2
      CO(O) = .5DO*(QN(O)+QN(1)-(A-B)*CO(1))
      IF (N .EQ. 1) RETURN

      CALL GAMMAF(A1,GA1)
      CALL GAMMAF(B1,GB1)
      CALL GAMMAF(AB2,GAB2)
      C = ((2.DO)**(AB+1.DO))*GA1*GB1/GAB2
      SU = 0.DO
      DO 1 J=0,N
      SU = SU+QN(J)*WT(J)
      CO(J) = 0.DO
1  CONTINUE
      CO(O) = SU/C
      CN = 0.DO
      DO 2 J=0,N
      X = ET(J)
      YP = QN(J)*WT(J)
      Y = .5DO*YP*(AB2*X+A-B)
      CN = CN+VN(J)*VN(J)*WT(J)
      DO 3 K=1,N
      CO(K) = CO(K)+Y
      DK = DFLOAT(K+1)
      CC = 2.DO*DK+AB
      C1 = 2.DO*DK*(DK+AB)*(CC-2.DO)

```

```
      C2 = (CC-1.DO)*(CC-2.DO)*CC
      C3 = (CC-1.DO)*(A-B)*AB
      C4 = 2.DO*(DK+A-1.DO)*CC*(DK+B-1.DO)
      YM = Y
      Y = ((C2*X+C3)*Y-C4*YP)/C1
      YP = YM
3     CONTINUE
2     CONTINUE

      DO 4 K=1,N-1
          DK = DFLOAT(K)
          C = C*(DK+A)*(DK+B)/DK
          CO(K) = CO(K)*(2.DO*DK+AB+1.DO)/C
          C = C/(DK+AB+1.DO)
4     CONTINUE
      CO(N) = CO(N)/CN

      RETURN
      END
```


COLEGL

For any $n \in \mathbf{N}$, the routine computes the $n + 1$ Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Legendre polynomial basis, of a polynomial q of degree at most n determined by the values attained at the Legendre Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$.

One has

$$c_k = \gamma_k \int_{-1}^1 q(x) P_k(x) dx = \gamma_k \sum_{j=0}^n q(\eta_j^{(n)}) P_k(\eta_j^{(n)}) \tilde{w}_j^{(n)} \quad 0 \leq k \leq n,$$

where $\gamma_k = (2k + 1)/2$ if $0 \leq k \leq n - 1$, while $\gamma_n = n/2$.

Here, $\tilde{w}_j^{(n)}$, $0 \leq j \leq n$, are the weights of the Legendre Gauss-Lobatto formula, which can be determined by calling subroutine WELEGL.

The nodes can be obtained by calling subroutine ZELEGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n ET , vector of the nodes $\eta_j^{(n)}$ QN , the values of q at the nodes WT , the weights $\tilde{w}_j^{(n)}$	CO , the Fourier coefficients of q

```

      SUBROUTINE COLEGL(N,ET,QN,WT,CO)
      *****
      *   COMPUTES THE LEGENDRE FOURIER COEFFICIENTS OF A POLYNOMIAL
      *   INDIVIDUATED BY ITS VALUES AT THE LEGENDRE GAUSS-LOBATTO NODES
      *   N = THE DEGREE OF THE POLYNOMIAL
      *   ET = VECTOR OF THE NODES, ET(I), I=0,N
      *   QN = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      *   WT = VECTOR OF THE WEIGHTS, WT(I), I=0,N
      *   CO = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), QN(0:*), WT(0:*), CO(0:*)
      CO(0) = QN(0)
      IF (N .EQ. 0) RETURN

      CO(0) = .5D0*(QN(0)+QN(1))
      CO(1) = .5D0*(QN(1)-QN(0))
      IF (N .EQ. 1) RETURN

      SU = 0.D0
      DO 1 J=0,N
      SU = SU+QN(J)*WT(J)
      CO(J) = 0.D0
1     CONTINUE
      CO(0) = .5D0*SU

      DO 2 J=0,N
      X = ET(J)
      YP = QN(J)*WT(J)
      Y = X*YP
      DO 3 K=1,N
      CO(K) = CO(K)+Y
      DK = DFLOAT(K+1)
      C1 = 2.D0*DK-1.D0
      C2 = DK-1.D0
      YM = Y
      Y = (C1*X*Y-C2*YP)/DK
      YP = YM
3     CONTINUE
2     CONTINUE

      DN = DFLOAT(N)
      CO(N) = .5D0*DN*CO(N)
      IF (N .EQ. 1) RETURN

      DO 4 K=1,N-1
      CO(K) = .5D0*CO(K)*(2.D0*DFLOAT(K)+1.D0)
4     CONTINUE

      RETURN
      END

```

COCHGL

For any $n \in \mathbf{N}$, the routine computes the $n + 1$ Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Chebyshev polynomial basis, of a polynomial q of degree at most n determined by the values attained at the Chebyshev Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$.

One has

$$c_k = \gamma_k \int_{-1}^1 q(x) T_k(x) \frac{dx}{\sqrt{1-x^2}}$$

$$= \frac{(-1)^k \gamma_k \pi}{n} \left(\frac{1}{2} q(\eta_0^{(n)}) + \frac{1}{2} (-1)^k q(\eta_n^{(n)}) + \sum_{j=1}^{n-1} q(\eta_j^{(n)}) \cos \frac{kj\pi}{n} \right) \quad 0 \leq k \leq n,$$

where $\gamma_0 = \gamma_n = 1/\pi$ and $\gamma_k = 2/\pi$, $1 \leq k \leq n - 1$.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n QN , the values of q at the nodes	CO , the Fourier coefficients of q

```

      SUBROUTINE COCHGL(N,QN,CO)
      *****
      *   COMPUTES THE CHEBYSHEV FOURIER COEFFICIENTS OF A POLYNOMIAL
      *   INDIVIDUATED BY ITS VALUES AT THE CHEBYSHEV GAUSS-LOBATTO NODES
      *   N   = THE DEGREE OF THE POLYNOMIAL
      *   QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      *   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION QN(0:*), CO(0:*)
      CO(0) = QN(0)
      IF (N .EQ. 0) RETURN

      PI = 3.14159265358979323846D0
      DN = DFLOAT(N)
      DD = DFLOAT(1+4*(N/2)-2*N)
      CO(0) = .5D0*(QN(0)+QN(N))
      CO(N) = .5D0*(QN(0)+DD*QN(N))
      IF (N .EQ. 1) RETURN

      SO = CO(0)
      SN = CO(N)
      SJ = -1.D0
      DO 1 J=1,N-1
      SO = SO+QN(J)
      SN = SN+QN(J)*SJ
      SJ = -SJ
1    CONTINUE
      CO(0) = SO/DN
      CO(N) = DD*SN/DN

      SK = -1.D0
      DO 2 K=1,N-1
      DK = DFLOAT(K)
      SU = .5D0*(QN(0)+QN(N)*SK)
      DO 3 J=1,N-1
      DJ = DFLOAT(J)
      SU = SU+QN(J)*DCOS(DK*DJ*PI/DN)
3    CONTINUE
      CO(K) = 2.D0*SK*SU/DN
      SK = -SK
2    CONTINUE

      RETURN
      END

```

COLAGR

For any $n \in \mathbf{N}$, the routine computes the n Fourier coefficients c_k , $0 \leq k \leq n-1$, with respect to the Laguerre polynomial basis, of a polynomial q of degree at most $n-1$ determined by the values attained at the Laguerre Gauss-Radau nodes $\eta_j^{(n)}$, $0 \leq j \leq n-1$.

One has

$$\begin{aligned} c_k &= \frac{k!}{\Gamma(k + \alpha + 1)} \int_0^{+\infty} q(x) L_k^{(\alpha)}(x) x^\alpha e^{-x} dx \\ &= \frac{k!}{\Gamma(k + \alpha + 1)} \sum_{j=0}^{n-1} q(\eta_j^{(n)}) L_k^{(\alpha)}(\eta_j^{(n)}) \tilde{w}_j^{(n)} \quad 0 \leq k \leq n-1. \end{aligned}$$

Here, $\tilde{w}_j^{(n)}$, $0 \leq j \leq n-1$, are the weights of the Laguerre Gauss-Radau formula, which can be determined by calling subroutine WELAGR.

The nodes can be obtained by calling subroutine ZELAGR.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n A , the parameter α B , the parameter β ET , vector of the nodes $\eta_j^{(n)}$ QN , the values of q at the nodes WT , the weights $\tilde{w}_j^{(n)}$	CO , the Fourier coefficients of q

Auxiliary routine: **GAMMAF**

```

      SUBROUTINE COLAGR(N,A,ET,QN,WT,CO)
      *****
      *   COMPUTES THE LAGUERRE FOURIER COEFFICIENTS OF A POLYNOMIAL
      *   INDIVIDUATED BY ITS VALUES AT THE LAGUERRE GAUSS-RADAU NODES
      *   N   = THE NUMBER OF NODES
      *   A   = PARAMETER >-1
      *   ET  = VECTOR OF THE NODES, ET(I), I=0,N-1
      *   QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N-1
      *   WT  = VECTOR OF THE WEIGHTS, WT(I), I=0,N-1
      *   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N-1
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(O:*), QN(O:*), WT(O:*), CO(O:*)
      IF (N .EQ. 0) RETURN
      A1 = A+1.DO
      CALL GAMMAF(A1,C)
      SU = 0.DO
      DO 1 J=0,N-1
      SU = SU+QN(J)*WT(J)
      CO(J) = 0.DO
1     CONTINUE
      CO(0) = SU/C
      IF (N .EQ. 1) RETURN

      DO 2 J=0,N-1
      X = ET(J)
      YP = QN(J)*WT(J)
      Y = (A1-X)*YP
      DO 3 K=1,N-1
      CO(K) = CO(K)+Y
      DK = DFLOAT(K+1)
      B1 = (2.DO*DK+A-1.DO-X)/DK
      B2 = (DK+A-1.DO)/DK
      YM = Y
      Y = B1*Y-B2*YP
      YP = YM
3     CONTINUE
2     CONTINUE

      DO 4 K=1,N-1
      DK = DFLOAT(K)
      C = C*(DK+A)/DK
      CO(K) = CO(K)/C
4     CONTINUE

      RETURN
      END

```

PVJAEX

For any $n \in \mathbf{N}$, the routine computes the value of a polynomial q of degree n and its first and second derivatives at a point x , from the Fourier coefficients of the expansion of q with respect to the Jacobi polynomial basis. More precisely, one has

$$q(x) = \sum_{k=0}^n c_k P_k^{(\alpha,\beta)}(x).$$

The values $q'(x)$ and $q''(x)$ are obtained by the recursion formula relating the Jacobi polynomials (see subroutine VAJAPO).

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of q in x
A , the parameter α	DY , the value of q' in x
B , the parameter β	$D2Y$, the value of q'' in x
X , the argument x	
CO , the Fourier coefficients of q	

SUBROUTINE PVJAEX(N,A,B,X,CO,Y,DY,D2Y)

```
*****
*   COMPUTES THE VALUE OF A POLYNOMIAL OF DEGREE N AND ITS FIRST AND
*   SECOND DERIVATIVES BY KNOWING THE JACOBI FOURIER COEFFICIENTS
*   N   = THE DEGREE OF THE POLYNOMIAL
*   A   = PARAMETER >-1
*   B   = PARAMETER >-1
*   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
*   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   Y   = VALUE OF THE POLYNOMIAL IN X
*   DY  = VALUE OF THE FIRST DERIVATIVE IN X
*   D2Y= VALUE OF THE SECOND DERIVATIVE IN X
*****
```



```

IMPLICIT DOUBLE PRECISION (A-H,O-Z)
DIMENSION CO(0:*)
  Y  = CO(0)
  DY = 0.DO
  D2Y = 0.DO
IF (N .EQ. 0) RETURN

  AB = A+B
  P  = .5D0*((AB+2.DO)*X+A-B)
  DP = .5D0*(AB+2.DO)
  D2P = 0.DO
  Y  = CO(0)+P*CO(1)
  DY = DP*CO(1)
  D2Y = 0.DO
IF (N .EQ. 1) RETURN

  PP = 1.DO
  DPP = 0.DO
  D2PP = 0.DO
DO 1 K=2,N
  DK = DFLOAT(K)
  CC = 2.DO*DK+AB
  C1 = 2.DO*DK*(DK+AB)*(CC-2.DO)
  C2 = (CC-1.DO)*(CC-2.DO)*CC
  C3 = (CC-1.DO)*(A-B)*AB
  C4 = 2.DO*(DK+A-1.DO)*CC*(DK+B-1.DO)
  PM = P
  P  = ((C2*X+C3)*P-C4*PP)/C1
  Y  = Y+P*CO(K)
  PP = PM
  DPM = DP
  DP  = ((C2*X+C3)*DP-C4*DPP+C2*PP)/C1
  DY  = DY+DP*CO(K)
  DPP = DPM
  D2PM = D2P
  D2P  = ((C2*X+C3)*D2P-C4*D2PP+2.DO*C2*DPP)/C1
  D2Y  = D2Y+D2P*CO(K)
  D2PP = D2PM
1  CONTINUE

RETURN
END

```

PVLEEX

For any $n \in \mathbf{N}$, the routine computes the value of a polynomial q of degree n and its first and second derivatives at a point x , from the Fourier coefficients of the expansion of q with respect to the Legendre polynomial basis. More precisely, one has

$$q(x) = \sum_{k=0}^n c_k P_k(x).$$

The values $q'(x)$ and $q''(x)$ are obtained by the recursion formula relating the Legendre polynomials (see subroutine VALEPO).

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of q in x
X , the argument x	DY , the value of q' in x
CO , the Fourier coefficients of q	$D2Y$, the value of q'' in x

```

      SUBROUTINE PVLEEX(N,X,CO,Y,DY,D2Y)
*****
*   COMPUTES THE VALUE OF A POLYNOMIAL OF DEGREE N AND ITS FIRST AND
*   SECOND DERIVATIVES BY KNOWING THE LEGENDRE FOURIER COEFFICIENTS
*   N   = THE DEGREE OF THE POLYNOMIAL
*   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
*   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   Y   = VALUE OF THE POLYNOMIAL IN X
*   DY  = VALUE OF THE FIRST DERIVATIVE IN X
*   D2Y = VALUE OF THE SECOND DERIVATIVE IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*)
      Y   = CO(0)
      DY  = 0.DO
      D2Y = 0.DO
      IF (N .EQ. 0) RETURN

      P   = X
      DP  = 1.DO
      D2P = 0.DO
      Y   = CO(0)+P*CO(1)
      DY  = DP*CO(1)
      D2Y = 0.DO
      IF (N .EQ. 1) RETURN

      PP  = 1.DO
      DPP = 0.DO
      D2PP = 0.DO
      DO 1 K=2,N
      DK = DFLOAT(K)
      C2 = 2.DO*DK-1.DO
      C4 = DK-1.DO
      PM = P
      P  = (C2*X*P-C4*PP)/DK
      Y  = Y+P*CO(K)
      PP = PM
      DPM = DP
      DP  = (C2*X*DP-C4*DPP+C2*PP)/DK
      DY  = DY+DP*CO(K)
      DPP = DPM
      D2PM = D2P
      D2P  = (C2*X*D2P-C4*D2PP+2.DO*C2*DPP)/DK
      D2Y  = D2Y+D2P*CO(K)
      D2PP = D2PM
1      CONTINUE

      RETURN
      END

```

PVCHEX

For any $n \in \mathbf{N}$, the routine computes the value of a polynomial q of degree n and its first and second derivatives at a point x , from the Fourier coefficients of the expansion of q with respect to the Chebyshev polynomial basis. More precisely, one has

$$q(x) = \sum_{k=0}^n c_k T_k(x).$$

The values $q'(x)$ and $q''(x)$ are obtained by the recursion formula relating the Chebyshev polynomials (see subroutine VACHPO).

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of q in x
X , the argument x	DY , the value of q' in x
CO , the Fourier coefficients of q	$D2Y$, the value of q'' in x

```

      SUBROUTINE PVCHEX(N,X,CO,Y,DY,D2Y)
*****
*   COMPUTES THE VALUE OF A POLYNOMIAL OF DEGREE N AND ITS FIRST AND
*   SECOND DERIVATIVES BY KNOWING THE CHEBYSHEV FOURIER COEFFICIENTS
*   N   = THE DEGREE OF THE POLYNOMIAL
*   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
*   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   Y   = VALUE OF THE POLYNOMIAL IN X
*   DY  = VALUE OF THE FIRST DERIVATIVE IN X
*   D2Y = VALUE OF THE SECOND DERIVATIVE IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*)
          Y   = CO(0)
          DY  = 0.DO
          D2Y = 0.DO
      IF (N .EQ. 0) RETURN

          P   = X
          DP  = 1.DO
          D2P = 0.DO
          Y   = CO(0)+P*CO(1)
          DY  = DP*CO(1)
          D2Y = 0.DO
      IF (N .EQ. 1) RETURN

          PP  = 1.DO
          DPP = 0.DO
          D2PP = 0.DO
      DO 1 K=2,N
          PM = P
          P  = 2.DO*X*P-PP
          Y  = Y+P*CO(K)
          PP = PM
          DPM = DP
          DP  = 2.DO*X*DP+2.DO*PP-DPP
          DY  = DY+DP*CO(K)
          DPP = DPM
          D2PM = D2P
          D2P  = 2.DO*X*D2P+4.DO*DPP-D2PP
          D2Y  = D2Y+D2P*CO(K)
          D2PP = D2PM
1      CONTINUE

      RETURN
      END

```

PVLAEX

For any $n \in \mathbf{N}$, the routine computes the value of a polynomial q of degree n and its first and second derivatives at a point x , from the Fourier coefficients of the expansion of q with respect to the Laguerre polynomial basis. More precisely, one has

$$q(x) = \sum_{k=0}^n c_k L_k^{(\alpha)}(x).$$

The values $q'(x)$ and $q''(x)$ are obtained by the recursion formula relating the Laguerre polynomials (see subroutine VALAPO).

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of q in x
A , the parameter α	DY , the value of q' in x
X , the argument x	$D2Y$, the value of q'' in x
CO , the Fourier coefficients of q	

```

      SUBROUTINE PVLAEX(N,A,X,CO,Y,DY,D2Y)
*****
*   COMPUTES THE VALUE OF A POLYNOMIAL OF DEGREE N AND ITS FIRST AND
*   SECOND DERIVATIVES BY KNOWING THE LAGUERRE FOURIER COEFFICIENTS
*   N = THE DEGREE OF THE POLYNOMIAL
*   A = PARAMETER >-1
*   X = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
*   CO = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   Y = VALUE OF THE POLYNOMIAL IN X
*   DY = VALUE OF THE FIRST DERIVATIVE IN X
*   D2Y= VALUE OF THE SECOND DERIVATIVE IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*)
      Y = CO(0)
      DY = 0.D0
      D2Y = 0.D0
      IF (N .EQ. 0) RETURN

      P = 1.D0+A-X
      DP = -1.D0
      D2P = 0.D0
      Y = CO(0)+P*CO(1)
      DY = DP*CO(1)
      D2Y = 0.D0
      IF (N .EQ. 1) RETURN

      PP = 1.D0
      DPP = 0.D0
      D2PP = 0.D0
      DO 1 K=2,N
      DK = DFLOAT(K)
      B1 = (2.D0*DK+A-1.D0-X)/DK
      B2 = (DK+A-1.D0)/DK
      PM = P
      P = B1*P-B2*PP
      Y = Y+P*CO(K)
      PP = PM
      DPM = DP
      DP = B1*DP-PP/DK-B2*DPP
      DY = DY+DP*CO(K)
      DPP = DPM
      D2PM = D2P
      D2P = B1*D2P-2.D0*DPP/DK-B2*D2PP
      D2Y = D2Y+D2P*CO(K)
      D2PP = D2PM
1      CONTINUE

      RETURN
      END

```

PVHEEX

For any $n \in \mathbf{N}$, the routine computes the value of a polynomial q of degree n and its first and second derivatives at a point x , from the Fourier coefficients of the expansion of q with respect to the Hermite polynomial basis. More precisely, one has

$$q(x) = \sum_{k=0}^n c_k H_k(x).$$

The values $q'(x)$ and $q''(x)$ are obtained by the recursion formula relating the Hermite polynomials (see subroutine VAHEPO).

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	Y , the value of q in x
X , the argument x	DY , the value of q' in x
CO , the Fourier coefficients of q	$D2Y$, the value of q'' in x


```

      SUBROUTINE PVHEEX(N,X,CO,Y,DY,D2Y)
*****
*   COMPUTES THE VALUE OF A POLYNOMIAL OF DEGREE N AND ITS FIRST AND
*   SECOND DERIVATIVES BY KNOWING THE HERMITE FOURIER COEFFICIENTS
*   N   = THE DEGREE OF THE POLYNOMIAL
*   X   = THE POINT IN WHICH THE COMPUTATION IS PERFORMED
*   CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   Y   = VALUE OF THE POLYNOMIAL IN X
*   DY  = VALUE OF THE FIRST DERIVATIVE IN X
*   D2Y = VALUE OF THE SECOND DERIVATIVE IN X
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*)
          Y   = CO(0)
          DY  = 0.DO
          D2Y = 0.DO
      IF (N .EQ. 0) RETURN

          P   = 2.DO*X
          DP  = 2.DO
          D2P = 0.DO
          Y   = CO(0)+P*CO(1)
          DY  = DP*CO(1)
          D2Y = 0.DO
      IF (N .EQ. 1) RETURN

          PP  = 1.DO
          DPP = 0.DO
          D2PP = 0.DO
      DO 1 K=2,N
          DK = DFLOAT(K)
          PM = P
          P   = 2.DO*X*P-2.DO*PP*(DK-1.DO)
          Y   = Y+P*CO(K)
          DY  = DY+2.DO*DK*PM*CO(K)
          D2Y = D2Y+4.DO*DK*(DK-1.DO)*PP*CO(K)
          PP  = PM
1      CONTINUE

      RETURN
      END

```

NOJAEX

For any $n \in \mathbf{N}$, the routine computes the quantity

$$\begin{aligned} \mathcal{W} &= \left(\int_{-1}^1 q^2(x) (1-x)^\alpha (1+x)^\beta dx \right)^{\frac{1}{2}} \\ &= \left(\sum_{k=0}^n c_k^2 \int_{-1}^1 [P_n^{(\alpha,\beta)}(x)]^2 (1-x)^\alpha (1+x)^\beta dx \right)^{\frac{1}{2}}, \end{aligned}$$

where q is a polynomial of degree at most n individuated by the Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Jacobi polynomial basis.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n A , the parameter α B , the parameter β CO , the Fourier coefficients of q	QW , the norm \mathcal{W}

Auxiliary routine: **GAMMAF**

```

      SUBROUTINE NOJAEX(N,A,B,CO,QW)
*****
*   COMPUTES THE INTEGRAL NORM OF A POLYNOMIAL BY KNOWING THE
*   FOURIER COEFFICIENTS WITH RESPECT TO THE JACOBI BASIS
*   N   = THE DEGREE OF THE POLYNOMIAL
*   A   = PARAMETER >-1
*   B   = PARAMETER >-1
*   CO  = COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   QW  = WEIGHTED INTEGRAL NORM OF THE POLYNOMIAL
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*)

      EPS = 1.D-14
      A1  = A+1.DO
      B1  = B+1.DO
      AB  = A+B
      AB2 = AB+2.DO
      DN  = DFLOAT(N)
      CALL GAMMAF(A1,GA1)
      CALL GAMMAF(B1,GB1)
      CALL GAMMAF(AB2,GAB2)
      C   = ((2.DO)**(AB+1.DO))*GA1*GB1/GAB2
      V   = DABS(CO(0))
      QW  = V*DSQRT(C)
      IF (N .EQ. 0) RETURN

      SU = 0.DO
      IF (V .LT. EPS) GOTO 1
      SU = C*V*V
1     DO 2 K=1,N
          DK = DFLOAT(K)
          C  = C*(DK+A)*(DK+B)/DK
          V  = DABS(CO(K))
          IF (V .LT. EPS) GOTO 3
          SU = SU+C*V*V/(2.DO*DK+AB+1.DO)
3     C  = C/(DK+AB+1.DO)
2     CONTINUE
      QW = DSQRT(SU)

      RETURN
      END

```

NOLEEX

For any $n \in \mathbf{N}$, the routine computes the quantity

$$\mathcal{I} = \left(\int_{-1}^1 q^2(x) dx \right)^{\frac{1}{2}} = \left(\sum_{k=0}^n \frac{2c_k^2}{2k+1} \right)^{\frac{1}{2}},$$

where q is a polynomial of degree at most n individuated by the Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Legendre polynomial basis.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n CO , the Fourier coefficients of q	QI , the norm \mathcal{I}

```
      SUBROUTINE NOLEEX(N,CO,QI)
*****
*   COMPUTES THE INTEGRAL NORM OF A POLYNOMIAL BY KNOWING THE
*   FOURIER COEFFICIENTS WITH RESPECT TO THE LEGENDRE BASIS
*   N   = THE DEGREE OF THE POLYNOMIAL
*   CO  = COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   QI  = INTEGRAL NORM OF THE POLYNOMIAL
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*)

      EPS = 1.D-14
      SU  = 0.DO
      DO 1 K=0,N
          DK = DFLOAT(K)
          V  = DABS(CO(K))
          IF (V .LT. EPS) GOTO 1
          SU = SU+V*V/(2.DO*DK+1.DO)
1      CONTINUE
          QI = DSQRT(2.DO*SU)

      RETURN
      END
```

NOCHEX

For any $n \in \mathbf{N}$, the routine computes the quantities

$$\mathcal{W} = \left(\int_{-1}^1 q^2(x) \frac{dx}{\sqrt{1-x^2}} \right)^{\frac{1}{2}} = \left(\pi c_0^2 + \frac{\pi}{2} \sum_{k=1}^n c_k^2 \right)^{\frac{1}{2}},$$

$$\mathcal{I} = \left(\int_{-1}^1 q^2(x) dx \right)^{\frac{1}{2}} = \left(\sum_{k=0}^n \sum_{m=0}^n c_k c_m I_{km} \right)^{\frac{1}{2}},$$

$$\text{with } I_{km} = \begin{cases} 0 & \text{if } k+m \text{ is odd,} \\ \frac{1}{1-(k+m)^2} + \frac{1}{1-(k-m)^2} & \text{if } k+m \text{ is even,} \end{cases}$$

where q is a polynomial of degree at most n individuated by the Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Chebyshev polynomial basis.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	QW , the norm \mathcal{W}
CO , the Fourier coefficients of q	QI , the norm \mathcal{I}

```

      SUBROUTINE NOCHEX(N,CO,QW,QI)
*****
*   COMPUTES THE INTEGRAL NORMS OF A POLYNOMIAL BY KNOWING THE
*   FOURIER COEFFICIENTS WITH RESPECT TO THE CHEBYSHEV BASIS
*   N   = THE DEGREE OF THE POLYNOMIAL
*   CO  = COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   QW  = WEIGHTED INTEGRAL NORM OF THE POLYNOMIAL
*   QI  = INTEGRAL NORM OF THE POLYNOMIAL
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*)

      EPS = 1.D-14
      PR = 1.77245385090551588D0
      R2 = 1.41421356237309515D0
      V = DABS(CO(0))
      QW = PR*V
      QI = R2*V
      IF (N .EQ. 0) RETURN

      SU = 0.D0
      IF (V .LT. EPS) GOTO 1
      SU = 2.D0*V*V
1     DO 2 K=1,N
          V = DABS(CO(K))
      IF (V .LT. EPS) GOTO 2
      SU = SU+V*V
2     CONTINUE
      QW = PR*DSQRT(.5D0*SU)

      SU = 0.D0
      DO 3 K=0,N,2
          V = CO(K)
      DO 3 M=0,N,2
          D1 = 1.D0-DFLOAT((K+M)*(K+M))
          D2 = 1.D0-DFLOAT((K-M)*(K-M))
          C = 1.D0/D1+1.D0/D2
          SU = SU+C*V*CO(M)
3     CONTINUE
      DO 4 K=1,N,2
          V = CO(K)
      DO 4 M=1,N,2
          D1 = 1.D0-DFLOAT((K+M)*(K+M))
          D2 = 1.D0-DFLOAT((K-M)*(K-M))
          C = 1.D0/D1+1.D0/D2
          SU = SU+C*V*CO(M)
4     CONTINUE
      QI = DSQRT(SU)

      RETURN
      END

```

NOLAEX

For any $n \in \mathbf{N}$, the routine computes the quantity

$$\mathcal{W} = \left(\int_0^{+\infty} q^2(x) x^\alpha e^{-x} dx \right)^{\frac{1}{2}} = \left(\sum_{k=0}^n c_k^2 \frac{\Gamma(k + \alpha + 1)}{k!} \right)^{\frac{1}{2}},$$

where q is a polynomial of degree at most n individuated by the Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Laguerre polynomial basis.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n A , the parameter α CO , the Fourier coefficients of q	QW , the norm \mathcal{W}

Auxiliary routine: **GAMMAF**


```
      SUBROUTINE NOLAEX(N,A,CO,QW)
*****
*   COMPUTES THE INTEGRAL NORM OF A POLYNOMIAL BY KNOWING THE
*   FOURIER COEFFICIENTS WITH RESPECT TO THE LAGUERRE BASIS
*   N   = THE DEGREE OF THE POLYNOMIAL
*   A   = PARAMETER >-1
*   CO  = COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   QW  = WEIGHTED INTEGRAL NORM OF THE POLYNOMIAL
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*)

      EPS = 1.D-14
      A1  = A+1.DO
      CALL GAMMAF(A1,C)
      V   = DABS(CO(0))
      QW  = V*DSQRT(C)
      IF (N .EQ. 0) RETURN

      SU = 0.DO
      IF (V .LT. EPS) GOTO 1
      SU = C*V*V
1     DO 2 K=1,N
          DK = DFLOAT(K)
          C  = C*(DK+A)/DK
          V  = DABS(CO(K))
          IF (V .LT. EPS) GOTO 2
          SU = SU+C*V*V
2     CONTINUE
      QW = DSQRT(SU)

      RETURN
      END
```

NOHEEX

For any $n \in \mathbf{N}$, the routine computes the quantity

$$\mathcal{W} = \left(\int_{-\infty}^{+\infty} q^2(x) e^{-x^2} dx \right)^{\frac{1}{2}} = \left(\sqrt{\pi} \sum_{k=0}^n c_k^2 2^k k! \right)^{\frac{1}{2}},$$

where q is a polynomial of degree at most n individuated by the Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Hermite polynomial basis.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n CO , the Fourier coefficients of q	QW , the norm \mathcal{W}

```
      SUBROUTINE NOHEEX(N,CO,QW)
*****
*   COMPUTES THE INTEGRAL NORM OF A POLYNOMIAL BY KNOWING THE
*   FOURIER COEFFICIENTS WITH RESPECT TO THE HERMITE BASIS
*   N   = THE DEGREE OF THE POLYNOMIAL
*   CO  = COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   QW  = WEIGHTED INTEGRAL NORM OF THE POLYNOMIAL
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*)

      EPS = 1.D-14
      PR  = 1.33133536380038953D0
      R2  = 1.41421356237309515D0
      V   = DABS(CO(0))
      QW  = V*PR
      IF (N .EQ. 0) RETURN

      SU = 0.DO
      IF (V .LT. EPS) GOTO 1
      SU = V*V
1      C = 1.DO
      DO 2 K=1,N
          DK = DFLOAT(K)
          C  = C*R2*DSQRT(DK)
          V  = DABS(CO(K))
          IF (V .LT. EPS) GOTO 2
          SU = SU+(C*V)*(C*V)
2      CONTINUE
      QW = PR*DSQRT(SU)

      RETURN
      END
```

COJADE

For any $n \in \mathbf{N}$, the routine computes the Fourier coefficients of the first derivative ($c_k^{(1)}$, $0 \leq k \leq n$) and the second derivative ($c_k^{(2)}$, $0 \leq k \leq n$) of a polynomial q of degree at most n determined by the Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Jacobi polynomial basis in the case $\alpha = \beta = \gamma$.

The first derivative coefficients are given by

$$c_k^{(1)} = \begin{cases} 0 & \text{if } k = n, \\ \frac{(2n+2\gamma-1)(n+\gamma)}{n+2\gamma} c_n & \text{if } k = n-1, \\ \frac{(2k+2\gamma+1)(k+\gamma+1)}{k+2\gamma+1} \left[\frac{k+\gamma+2}{(2k+2\gamma+5)(k+2\gamma+2)} c_{k+2}^{(1)} + c_{k+1} \right] & \text{if } 1 \leq k \leq n-2, \\ \frac{\gamma+2}{4\gamma+10} c_2^{(1)} + (\gamma+1)c_1 & \text{if } k = 0. \end{cases}$$

The second derivative coefficients are obtained by assuming $c_k^{(1)}$ in place of c_k in the above formula.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	CD , the Fourier coefficients of q'
G , the parameter γ	$CD2$, the Fourier coefficients of q''
CO , the Fourier coefficients of q	

```

      SUBROUTINE COJADE(N,G,CO,CD,CD2)
      *****
      *   COMPUTES THE FOURIER COEFFICIENTS OF THE DERIVATIVES OF A POLYNOMIAL
      *   FROM ITS FOURIER COEFFICIENTS WITH RESPECT TO THE JACOBI BASIS
      *   N = THE DEGREE OF THE POLYNOMIAL
      *   G = PARAMETER >-1
      *   CO = COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
      *   CD = COEFFICIENTS OF THE FIRST DERIVATIVE, CD(I), I=0,N
      *   CD2 = COEFFICIENTS OF THE SECOND DERIVATIVE, CD2(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*), CD(0:*), CD2(0:*)

      CD(N) = 0.DO
      CD2(N) = 0.DO
      IF (N .EQ. 0) RETURN

      CD(0) = (G+1.DO)*CO(1)
      CD2(N-1) = 0.DO
      IF (N .EQ. 1) RETURN

      DN = DFLOAT(N)
      G2 = 2.DO*G
      CD(N-1) = (2.DO*DN+G2-1.DO)*(DN+G)*CO(N)/(DN+G2)
      DO 1 K=0,N-2
      KR = N-K-2
      IF (KR .NE. 0) THEN
      DK = DFLOAT(KR)
      C1 = (2.DO*DK+G2+1.DO)*(DK+G+1.DO)/(DK+G2+1.DO)
      C2 = (DK+G+2.DO)/((2.DO*DK+G2+5.DO)*(DK+G2+2.DO))
      CD(KR) = C1*(C2*CD(KR+2)+CO(KR+1))
      CD2(KR) = C1*(C2*CD2(KR+2)+CD(KR+1))
      ELSE
      CD(0) = .25D0*(G+2.DO)*CD(2)/(G+2.5D0)+(G+1.DO)*CO(1)
      CD2(0) = .25D0*(G+2.DO)*CD2(2)/(G+2.5D0)+(G+1.DO)*CD(1)
      ENDIF
1      CONTINUE

      RETURN
      END

```

COLEDE

For any $n \in \mathbf{N}$, the routine computes the Fourier coefficients of the first derivative ($c_k^{(1)}$, $0 \leq k \leq n$) and the second derivative ($c_k^{(2)}$, $0 \leq k \leq n$) of a polynomial q of degree at most n determined by the Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Legendre polynomial basis.

The first derivative coefficients are given by

$$c_k^{(1)} = \begin{cases} 0 & \text{if } k = n, \\ (2n - 1)c_n & \text{if } k = n - 1, \\ (2k + 1) \left(\frac{1}{2k+5} c_{k+2}^{(1)} + c_{k+1} \right) & \text{if } 0 \leq k \leq n - 2. \end{cases}$$

The second derivative coefficients are obtained by assuming $c_k^{(1)}$ in place of c_k in the above formula.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	CD , the Fourier coefficients of q'
CO , the Fourier coefficients of q	$CD2$, the Fourier coefficients of q''

```

      SUBROUTINE COLEDE(N,CO,CD,CD2)
*****
*   COMPUTES THE FOURIER COEFFICIENTS OF THE DERIVATIVES OF A POLYNOMIAL
*   FROM ITS FOURIER COEFFICIENTS WITH RESPECT TO THE LEGENDRE BASIS
*   N = THE DEGREE OF THE POLYNOMIAL
*   CO = COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   CD = COEFFICIENTS OF THE FIRST DERIVATIVE, CD(I), I=0,N
*   CD2 = COEFFICIENTS OF THE SECOND DERIVATIVE, CD2(I), I=0,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*), CD(0:*), CD2(0:*)

      CD(N) = 0.DO
      CD2(N) = 0.DO
      IF (N .EQ. 0) RETURN

      DN = DFLOAT(N)
      CD(N-1) = (2.DO*DN-1.DO)*CO(N)
      CD2(N-1) = 0.DO
      IF (N .EQ. 1) RETURN

      DO 1 K=0,N-2
      KR = N-K-2
      DK = 2.DO*DFLOAT(KR)+1.DO
      CD(KR) = DK*(CD(KR+2)/(DK+4.DO)+CO(KR+1))
      CD2(KR) = DK*(CD2(KR+2)/(DK+4.DO)+CD(KR+1))
1      CONTINUE

      RETURN
      END

```

COCHDE

For any $n \in \mathbf{N}$, the routine computes the Fourier coefficients of the first derivative ($c_k^{(1)}$, $0 \leq k \leq n$) and the second derivative ($c_k^{(2)}$, $0 \leq k \leq n$) of a polynomial q of degree at most n determined by the Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Chebyshev polynomial basis.

The first derivative coefficients are given by

$$c_k^{(1)} = \begin{cases} 0 & \text{if } k = n, \\ 2nc_n & \text{if } k = n - 1, \\ c_{k+2}^{(1)} + 2(k+1)c_{k+1} & \text{if } 1 \leq k \leq n - 2, \\ \frac{1}{2}c_2^{(1)} + c_1 & \text{if } k = 0. \end{cases}$$

The second derivative coefficients are obtained by assuming $c_k^{(1)}$ in place of c_k in the above formula.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	CD , the Fourier coefficients of q'
CO , the Fourier coefficients of q	$CD2$, the Fourier coefficients of q''


```

      SUBROUTINE COCHDE(N,CO,CD,CD2)
      *****
      *   COMPUTES THE FOURIER COEFFICIENTS OF THE DERIVATIVES OF A POLYNOMIAL
      *   FROM ITS FOURIER COEFFICIENTS WITH RESPECT TO THE CHEBYSHEV BASIS
      *   N = THE DEGREE OF THE POLYNOMIAL
      *   CO = COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
      *   CD = COEFFICIENTS OF THE FIRST DERIVATIVE, CD(I), I=0,N
      *   CD2 = COEFFICIENTS OF THE SECOND DERIVATIVE, CD2(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*), CD(0:*), CD2(0:*)

      CD(N) = 0.DO
      CD2(N) = 0.DO
      IF (N .EQ. 0) RETURN

      CD(0) = CO(1)
      CD2(N-1) = 0.DO
      IF (N .EQ. 1) RETURN

      DN = DFLOAT(N)
      CD(N-1) = 2.DO*DN*CO(N)
      DO 1 K=0,N-2
      KR = N-K-2
      IF (KR .NE. 0) THEN
      DK = 2.DO*(DFLOAT(KR)+1.DO)
      CD(KR) = CD(KR+2)+DK*CO(KR+1)
      CD2(KR) = CD2(KR+2)+DK*CD(KR+1)
      ELSE
      CD(0) = .5DO*CD(2)+CO(1)
      CD2(0) = .5DO*CD2(2)+CD(1)
      ENDIF
1      CONTINUE

      RETURN
      END

```

COLADE

For any $n \in \mathbf{N}$, the routine computes the Fourier coefficients of the first derivative ($c_k^{(1)}$, $0 \leq k \leq n$) and the second derivative ($c_k^{(2)}$, $0 \leq k \leq n$) of a polynomial q of degree at most n determined by the Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Laguerre polynomial basis.

The first derivative coefficients are given by

$$c_k^{(1)} = \begin{cases} 0 & \text{if } k = n, \\ c_{k+1}^{(1)} - c_{k+1} & \text{if } 0 \leq k \leq n - 1. \end{cases}$$

The second derivative coefficients are obtained by assuming $c_k^{(1)}$ in place of c_k in the above formula.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	CD , the Fourier coefficients of q'
CO , the Fourier coefficients of q	$CD2$, the Fourier coefficients of q''

```
      SUBROUTINE COLADE(N,CO,CD,CD2)
*****
*   COMPUTES THE FOURIER COEFFICIENTS OF THE DERIVATIVES OF A POLYNOMIAL
*   FROM ITS FOURIER COEFFICIENTS WITH RESPECT TO THE LAGUERRE BASIS
*   N = THE DEGREE OF THE POLYNOMIAL
*   CO = COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   CD = COEFFICIENTS OF THE FIRST DERIVATIVE, CD(I), I=0,N
*   CD2 = COEFFICIENTS OF THE SECOND DERIVATIVE, CD2(I), I=0,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*), CD(0:*), CD2(0:*)

      CD(N) = 0.DO
      CD2(N) = 0.DO
      IF (N .EQ. 0) RETURN

      CD(N-1) = -CO(N)
      CD2(N-1) = 0.DO
      IF (N .EQ. 1) RETURN

      DO 1 K=0,N-2
      KR = N-K-2
      CD(KR) = CD(KR+1)-CO(KR+1)
      CD2(KR) = CD2(KR+2)-CD(KR+1)
1      CONTINUE

      RETURN
      END
```

COHEDE

For any $n \in \mathbf{N}$, the routine computes the Fourier coefficients of the first derivative ($c_k^{(1)}$, $0 \leq k \leq n$) and the second derivative ($c_k^{(2)}$, $0 \leq k \leq n$) of a polynomial q of degree at most n determined by the Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Hermite polynomial basis.

The first derivative coefficients are given by

$$c_k^{(1)} = \begin{cases} 0 & \text{if } k = n, \\ 2(k+1)c_{k+1} & \text{if } 0 \leq k \leq n-1. \end{cases}$$

The second derivative coefficients are obtained by assuming $c_k^{(1)}$ in place of c_k in the above formula.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	CD , the Fourier coefficients of q'
CO , the Fourier coefficients of q	$CD2$, the Fourier coefficients of q''

```
      SUBROUTINE COHEDE(N,CO,CD,CD2)
*****
*   COMPUTES THE FOURIER COEFFICIENTS OF THE DERIVATIVES OF A POLYNOMIAL
*   FROM ITS FOURIER COEFFICIENTS WITH RESPECT TO THE HERMITE BASIS
*   N = THE DEGREE OF THE POLYNOMIAL
*   CO = COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
*   CD = COEFFICIENTS OF THE FIRST DERIVATIVE, CD(I), I=0,N
*   CD2 = COEFFICIENTS OF THE SECOND DERIVATIVE, CD2(I), I=0,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CO(0:*), CD(0:*), CD2(0:*)

      CD(N) = 0.DO
      CD2(N) = 0.DO
      IF (N .EQ. 0) RETURN

      DN = DFLOAT(N)
      CD(N-1) = 2.DO*DN*CO(N)
      CD2(N-1) = 0.DO
      IF (N .EQ. 1) RETURN

      DO 1 K=0,N-2
      KR = N-K-2
      DK = 2.DO*DFLOAT(KR)+2.DO
      CD(KR) = DK*CO(KR+1)
      CD2(KR) = DK*CD(KR+1)
1      CONTINUE

      RETURN
      END
```

DEJAGA

For any $n \geq 1$, the routine computes the first derivative of a polynomial q of degree at most $n - 1$ at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the Jacobi polynomial of degree n . The polynomial q is individuated by the values attained at the zeroes.

We have the formula

$$q'(\xi_i^{(n)}) = \sum_{j=1}^n d_{ij}^{(1)} q(\xi_j^{(n)}),$$

where

$$d_{ij}^{(1)} = \begin{cases} \frac{\frac{d}{dx} P_n^{(\alpha, \beta)}(\xi_i^{(n)})}{\frac{d}{dx} P_n^{(\alpha, \beta)}(\xi_j^{(n)})} \frac{1}{\xi_i^{(n)} - \xi_j^{(n)}} & \text{if } i \neq j, \\ \frac{(\alpha + \beta + 2)\xi_i^{(n)} + \alpha - \beta}{2(1 - [\xi_i^{(n)}]^2)} & \text{if } i = j. \end{cases}$$

The Legendre case is obtained by setting $\alpha = \beta = 0$. Similarly, the Chebyshev case is obtained by setting $\alpha = \beta = -\frac{1}{2}$. The zeroes and the values $\frac{d}{dx} P_n^{\alpha, \beta}(\xi_j^{(n)})$, $1 \leq j \leq n$, can be determined by calling subroutine ZEJAGA.

<i>Input variables</i>	<i>Output variable</i>
N , the number of zeroes	DQZ , the values of q' at the zeroes
A , the parameter α	
B , the parameter β	
CZ , vector of the zeroes $\xi_j^{(n)}$	
DZ , the values $\frac{d}{dx} P_n^{(\alpha, \beta)}(\xi_j^{(n)})$	
QZ , the values of q at the zeroes	

```

      SUBROUTINE DEJAGA(N,A,B,CS,DZ,QZ,DQZ)
*****
*   COMPUTES THE DERIVATIVE OF A POLYNOMIAL AT THE JACOBI ZEROES FROM
*   THE VALUES OF THE POLYNOMIAL ATTAINED AT THE SAME POINTS
*   N = THE NUMBER OF ZEROES
*   A = PARAMETER >-1
*   B = PARAMETER >-1
*   CS = ZEROES OF THE JACOBI POLYNOMIAL, CS(I), I=1,N
*   DZ = IACOBI DERIVATIVES AT THE ZEROES, DZ(I), I=1,N
*   QZ = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
*   DQZ = DERIVATIVES OF THE POLYNOMIAL AT THE ZEROES, DQZ(I), I=1,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), DZ(1), QZ(1), DQZ(1)
      IF (N .EQ. 0) RETURN

      DO 1 I=1,N
      SU = 0.DO
      CI = CS(I)
      DI = DZ(I)
      DO 2 J=1,N
      IF (I .NE. J) THEN
      CJ = CS(J)
      DJ = DZ(J)
      SU = SU+QZ(J)/(DJ*(CI-CJ))
      ELSE
      SU = SU+.5DO*QZ(I)*((A+B+2.DO)*CI+A-B)/(DI*(1.DO-CI*CI))
      ENDIF
2      CONTINUE
      DQZ(I) = DI*SU
1      CONTINUE

      RETURN
      END

```

DELAGA

For any $n \geq 1$, the routine computes the first derivative of a polynomial q of degree at most $n - 1$ at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the Laguerre polynomial of degree n . The polynomial q is individuated by the values attained at the zeroes.

We have the formula

$$q'(\xi_i^{(n)}) = \sum_{j=1}^n d_{ij}^{(1)} q(\xi_j^{(n)}),$$

where

$$d_{ij}^{(1)} = \begin{cases} \frac{\frac{d}{dx} L_n^{(\alpha)}(\xi_i^{(n)})}{\frac{d}{dx} L_n^{(\alpha)}(\xi_j^{(n)})} \frac{1}{\xi_i^{(n)} - \xi_j^{(n)}} & \text{if } i \neq j, \\ \frac{\xi_i^{(n)} - \alpha - 1}{2\xi_i^{(n)}} & \text{if } i = j. \end{cases}$$

The zeroes can be determined by calling subroutine ZELAGA.

<i>Input variables</i>	<i>Output variable</i>
N , the number of zeroes A , the parameter α CZ , vector of the zeroes $\xi_j^{(n)}$ QZ , the values of q at the zeroes	DQZ , the values of q' at the zeroes

Auxiliary routine: **VALAPO**


```

      SUBROUTINE DELAGA(N,A,CS,QZ,DQZ)
      *****
      *   COMPUTES THE DERIVATIVE OF A POLYNOMIAL AT THE LAGUERRE ZEROES FROM
      *   THE VALUES OF THE POLYNOMIAL ATTAINED AT THE SAME POINTS
      *   N = THE NUMBER OF ZEROES
      *   A = PARAMETER >-1
      *   CS = ZEROES OF THE LAGUERRE POLYNOMIAL, CS(I), I=1,N
      *   QZ = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
      *   DQZ = DERIVATIVES OF THE POLYNOMIAL AT THE ZEROES, DQZ(I), I=1,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), QZ(1), DQZ(1)
      IF (N .EQ. 0) RETURN

      DO 1 J=1,N
      CJ = CS(J)
      CALL VALAPO(N,A,CJ,Y,DY,D2Y)
      QZ(J) = QZ(J)/DY
1     CONTINUE

      DO 2 I=1,N
      SU = 0.DO
      CI = CS(I)
      CALL VALAPO(N,A,CI,Y,DI,D2Y)
      DO 3 J=1,N
      IF (I .NE. J) THEN
      CJ = CS(J)
      SU = SU+QZ(J)/(CI-CJ)
      ELSE
      SU = SU+.5D0*QZ(I)*(CI-A-1.DO)/CI
      ENDIF
3     CONTINUE
      DQZ(I) = DI*SU
2     CONTINUE

      DO 4 I=1,N
      CI = CS(I)
      CALL VALAPO(N,A,CI,Y,DY,D2Y)
      QZ(I) = DY*QZ(I)
4     CONTINUE

      RETURN
      END

```

DEHEGA

For any $n \geq 1$, the routine computes the first derivative of a polynomial q of degree at most $n - 1$ at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the Hermite polynomial of degree n . The polynomial q is individuated by the values attained at the zeroes.

We have the formula

$$q'(\xi_i^{(n)}) = \sum_{j=1}^n d_{ij}^{(1)} q(\xi_j^{(n)}),$$

where

$$d_{ij}^{(1)} = \begin{cases} \frac{H'_n(\xi_i^{(n)})}{H'_n(\xi_j^{(n)})} \frac{1}{\xi_i^{(n)} - \xi_j^{(n)}} & \text{if } i \neq j, \\ \xi_i^{(n)} & \text{if } i = j. \end{cases}$$

The zeroes can be determined by calling subroutine ZEHEGA.

<i>Input variables</i>	<i>Output variable</i>
N , the number of zeroes	DQZ , the values of q' at the zeroes
CZ , vector of the zeroes $\xi_j^{(n)}$	
QZ , the values of q at the zeroes	

Auxiliary routine: **VAHEPO**

```

      SUBROUTINE DEHEGA(N,CS,QZ,DQZ)
*****
*   COMPUTES THE DERIVATIVE OF A POLYNOMIAL AT THE HERMITE ZEROES FROM
*   THE VALUES OF THE POLYNOMIAL ATTAINED AT THE SAME POINTS
*   N = THE NUMBER OF ZEROES
*   CS = ZEROES OF THE HERMITE POLYNOMIAL, CS(I), I=1,N
*   QZ = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
*   DQZ = DERIVATIVES OF THE POLYNOMIAL AT THE ZEROES, DQZ(I), I=1,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION CS(1), QZ(1), DQZ(1)
      IF (N .EQ. 0) RETURN

      DO 1 J=1,N
      CJ = CS(J)
      CALL VAHEPO(N,CJ,Y,DY,D2Y)
      QZ(J) = QZ(J)/DY
1     CONTINUE

      DO 2 I=1,N
      SU = 0.DO
      CI = CS(I)
      CALL VAHEPO(N,CI,Y,DI,D2Y)
      DO 3 J=1,N
      IF (I .NE. J) THEN
      CJ = CS(J)
      SU = SU+QZ(J)/(CI-CJ)
      ELSE
      SU = SU+CI*QZ(I)
      ENDIF
3     CONTINUE
      DQZ(I) = DI*SU
2     CONTINUE

      DO 4 I=1,N
      CI = CS(I)
      CALL VAHEPO(N,CI,Y,DY,D2Y)
      QZ(I) = DY*QZ(I)
4     CONTINUE

      RETURN
      END

```

DEJAGL

For any $n \in \mathbf{N}$, the routine computes the first derivative of a polynomial q of degree at most n at the Jacobi Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$. The polynomial q is individuated by the values attained at these nodes.

We have the formula

$$q'(\eta_i^{(n)}) = \sum_{j=0}^n \tilde{d}_{ij}^{(1)} q(\eta_j^{(n)}),$$

where

$$\tilde{d}_{ij}^{(1)} = \left\{ \begin{array}{ll} \frac{\alpha - n(n + \alpha + \beta + 1)}{2(\beta + 2)} & \text{if } i = j = 0, \\ \frac{(\beta + 1)P_n^{(\alpha, \beta)}(\eta_i^{(n)})}{(1 + \eta_i^{(n)})P_n^{(\alpha, \beta)}(\eta_0^{(n)})} & \text{if } 1 \leq i \leq n - 1, j = 0, \\ \frac{(\beta + 1)P_n^{(\alpha, \beta)}(\eta_n^{(n)})}{2(\alpha + 1)P_n^{(\alpha, \beta)}(\eta_0^{(n)})} & \text{if } i = n, j = 0, \\ \frac{-P_n^{(\alpha, \beta)}(\eta_0^{(n)})}{(\beta + 1)(1 + \eta_j^{(n)})P_n^{(\alpha, \beta)}(\eta_j^{(n)})} & \text{if } i = 0, 1 \leq j \leq n - 1, \\ \frac{P_n^{(\alpha, \beta)}(\eta_i^{(n)})}{P_n^{(\alpha, \beta)}(\eta_j^{(n)})} \frac{1}{\eta_i^{(n)} - \eta_j^{(n)}} & \text{if } 1 \leq i \leq n - 1, 1 \leq j \leq n - 1, i \neq j, \\ \frac{(\alpha + \beta)\eta_i^{(n)} + \alpha - \beta}{2(1 - [\eta_i^{(n)}]^2)} & \text{if } 1 \leq i = j \leq n - 1, \\ \frac{P_n^{(\alpha, \beta)}(\eta_n^{(n)})}{(\alpha + 1)(1 - \eta_j^{(n)})P_n^{(\alpha, \beta)}(\eta_j^{(n)})} & \text{if } i = n, 1 \leq j \leq n - 1, \\ \frac{-(\alpha + 1)P_n^{(\alpha, \beta)}(\eta_0^{(n)})}{2(\beta + 1)P_n^{(\alpha, \beta)}(\eta_n^{(n)})} & \text{if } i = 0, j = n, \\ \frac{-(\alpha + 1)P_n^{(\alpha, \beta)}(\eta_i^{(n)})}{(1 - \eta_i^{(n)})P_n^{(\alpha, \beta)}(\eta_n^{(n)})} & \text{if } 1 \leq i \leq n - 1, j = n, \\ \frac{n(n + \alpha + \beta + 1) - \beta}{2(\alpha + 2)} & \text{if } i = j = n. \end{array} \right.$$

The nodes and the values $P_n^{(\alpha, \beta)}(\eta_j^{(n)})$, $0 \leq j \leq n$, can be determined by calling subroutine ZEJAGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n A , the parameter α B , the parameter β ET , vector of the nodes $\eta_j^{(n)}$ VN , the values $P_n^{(\alpha,\beta)}(\eta_j^{(n)})$ QN , the values of q at the nodes	DQN , the values of q' at the nodes

```

      SUBROUTINE DEJAGL(N,A,B,ET,VN,QN,DQN)
      *****
      *   COMPUTES THE DERIVATIVE OF A POLYNOMIAL AT THE JACOBI GAUSS-LOBATTO
      *   NODES FROM THE VALUES OF THE POLYNOMIAL ATTAINED AT THE SAME POINTS
      *   N = THE DEGREE OF THE POLYNOMIAL
      *   A = PARAMETER >-1
      *   B = PARAMETER >-1
      *   ET = VECTOR OF THE NODES, ET(I), I=0,N
      *   VN = VALUES OF THE IACOBI POLYNOMIAL AT THE NODES, VN(I), I=0,N
      *   QN = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      *   DQN = DERIVATIVES OF THE POLYNOMIAL AT THE NODES, DQZ(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(O:*), VN(O:*), QN(O:*), DQN(O:*)
      DQN(O) = 0.DO
      IF (N .EQ. 0) RETURN

      A1 = A+1.DO
      B1 = B+1.DO
      AB = A+B
      DN = DFLOAT(N)
      C1 = DN*(DN+AB+1.DO)
      C2 = A1*VN(O)/(B1*VN(N))
      DQN(O) = .5DO*((A-C1)*QN(O)/(B+2.DO)-C2*QN(N))
      DQN(N) = .5DO*(QN(O)/C2+(C1-B)*QN(N)/(A+2.DO))
      IF (N .EQ. 1) RETURN

      S1 = DQN(O)
      S2 = DQN(N)

```

```

C3 = -VN(0)/B1
C4 = VN(N)/A1
  DO 1 J=1,N-1
VJ = QN(J)/VN(J)
EJ = ET(J)
S1 = S1+C3*VJ/(1.DO+EJ)
S2 = S2+C4*VJ/(1.DO-EJ)
1  CONTINUE
DQN(0) = S1
DQN(N) = S2

  DO 2 I=1,N-1
VI = VN(I)
EI = ET(I)
SU = B1*QN(0)/((1.DO+EI)*VN(0))-A1*QN(N)/((1.DO-EI)*VN(N))
  DO 3 J=1,N-1
  IF (I .NE. J) THEN
VJ = VN(J)
EJ = ET(J)
SU = SU+QN(J)/(VJ*(EI-EJ))
  ELSE
SU = SU+.5DO*QN(I)*(AB*EI+A-B)/(VI*(1.DO-EI*EI))
  ENDIF
3  CONTINUE
DQN(I) = VI*SU
2  CONTINUE

RETURN
END

```

.

DELEGL

For any $n \in \mathbf{N}$, the routine computes the first derivative of a polynomial q of degree at most n at the Legendre Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$. The polynomial q is individuated by the values attained at these nodes.

We have the formula

$$q'(\eta_i^{(n)}) = \sum_{j=0}^n \tilde{d}_{ij}^{(1)} q(\eta_j^{(n)}),$$

where

$$\tilde{d}_{ij}^{(1)} = \begin{cases} -\frac{1}{4}n(n+1) & \text{if } i = j = 0, \\ \frac{P_n(\eta_i^{(n)})}{P_n(\eta_j^{(n)})} \frac{1}{\eta_i^{(n)} - \eta_j^{(n)}} & \text{if } 0 \leq i \leq n, 0 \leq j \leq n, i \neq j, \\ 0 & \text{if } 1 \leq i = j \leq n-1, \\ \frac{1}{4}n(n+1) & \text{if } i = j = n. \end{cases}$$

The nodes and the values $P_n(\eta_j^{(n)})$, $0 \leq j \leq n$, can be determined by calling subroutine ZELEGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n	DQN , the values of q' at the nodes
ET , vector of the nodes $\eta_j^{(n)}$	
VN , the values $P_n(\eta_j^{(n)})$	
QN , the values of q at the nodes	


```

      SUBROUTINE DELEGL(N,ET,VN,QN,DQN)
      *****
      * COMPUTES THE DERIVATIVE OF A POLYNOMIAL AT THE LEGENDRE GAUSS-LOBATTO
      * NODES FROM THE VALUES OF THE POLYNOMIAL ATTAINED AT THE SAME POINTS
      * N = THE DEGREE OF THE POLYNOMIAL
      * ET = VECTOR OF THE NODES, ET(I), I=0,N
      * VN = VALUES OF THE LEGENDRE POLYNOMIAL AT THE NODES, VN(I), I=0,N
      * QN = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      * DQN = DERIVATIVES OF THE POLYNOMIAL AT THE NODES, DQZ(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(O:*), VN(O:*), QN(O:*), DQN(O:*)
      DQN(O) = 0.DO
      IF (N .EQ. 0) RETURN

      DO 1 I=0,N
      SU = 0.DO
      VI = VN(I)
      EI = ET(I)
      DO 2 J=0,N
      IF (I .EQ. J) GOTO 2
      VJ = VN(J)
      EJ = ET(J)
      SU = SU+QN(J)/(VJ*(EI-EJ))
2      CONTINUE
      DQN(I) = VI*SU
1      CONTINUE

      DN = DFLOAT(N)
      C = .25DO*DN*(DN+1.DO)
      DQN(O) = DQN(O)-C*QN(O)
      DQN(N) = DQN(N)+C*QN(N)

      RETURN
      END

```

DECHGL

For any $n \in \mathbf{N}$, the routine computes the first derivative of a polynomial q of degree at most n at the Chebyshev Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$. The polynomial q is individuated by the values attained at these nodes.

We have the formula

$$q'(\eta_i^{(n)}) = \sum_{j=0}^n \tilde{d}_{ij}^{(1)} q(\eta_j^{(n)}),$$

where

$$\tilde{d}_{ij}^{(1)} = \begin{cases} -\frac{1}{6}(2n^2 + 1) & \text{if } i = j = 0, \\ \frac{1}{2}(-1)^i / (1 + \eta_i^{(n)}) & \text{if } 1 \leq i \leq n-1, j = 0, \\ \frac{1}{2}(-1)^n & \text{if } i = n, j = 0, \\ -2(-1)^j / (1 + \eta_j^{(n)}) & \text{if } i = 0, 1 \leq j \leq n-1, \\ \frac{(-1)^{i+j}}{\eta_i^{(n)} - \eta_j^{(n)}} & \text{if } 1 \leq i \leq n-1, 1 \leq j \leq n-1, i \neq j, \\ \frac{-\eta_i^{(n)}}{2(1 - [\eta_i^{(n)}]^2)} & \text{if } 1 \leq i = j \leq n-1, \\ 2(-1)^{n+j} / (1 - \eta_j^{(n)}) & \text{if } i = n, 1 \leq j \leq n-1, \\ -\frac{1}{2}(-1)^n & \text{if } i = 0, j = n, \\ -\frac{1}{2}(-1)^{n+i} / (1 - \eta_i^{(n)}) & \text{if } 1 \leq i \leq n-1, j = n, \\ \frac{1}{6}(2n^2 + 1) & \text{if } i = j = n. \end{cases}$$

The nodes can be determined by calling subroutine ZECHGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n ET , vector of the nodes $\eta_j^{(n)}$ QN , the values of q at the nodes	DQN , the values of q' at the nodes

```

      SUBROUTINE DECHGL(N,ET,QN,DQN)
      *****
      * COMPUTES THE DERIVATIVE OF A POLYNOMIAL AT THE CHEBYSHEV GAUSS-LOBATTO
      * NODES FROM THE VALUES OF THE POLYNOMIAL ATTAINED AT THE SAME POINTS
      *   N = THE DEGREE OF THE POLYNOMIAL
      *   ET = VECTOR OF THE NODES, ET(I), I=0,N
      *   QN = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      *   DQN = DERIVATIVES OF THE POLYNOMIAL AT THE NODES, DQZ(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), QN(0:*), DQN(0:*)
      DQN(0) = 0.DO
      IF (N .EQ. 0) RETURN

      DN = DFLOAT(N)
      CN = (2.DO*DN*DN+1.DO)/6.DO
      SN = DFLOAT(1+4*(N/2)-2*N)
      DQN(0) = -CN*QN(0)-.5DO*SN*QN(N)
      DQN(N) = .5DO*SN*QN(0)+CN*QN(N)
      IF (N .EQ. 1) RETURN

      S1 = DQN(0)
      S2 = DQN(N)
      SGN = -1.DO
      DO 1 J=1,N-1
      EJ = ET(J)
      QJ = 2.DO*SGN*QN(J)
      S1 = S1-QJ/(1.DO+EJ)
      S2 = S2+QJ*SN/(1.DO-EJ)
      SGN = -SGN
1      CONTINUE
      DQN(0) = S1
      DQN(N) = S2

```

```
SGNI = -1.DO
  DO 2 I=1,N-1
EI = ET(I)
SU = .5DO*SGNI*(QN(O)/(1.DO+EI)-SN*QN(N)/(1.DO-EI))
SGNJ = -1.DO
  DO 3 J=1,N-1
    IF (I .NE. J) THEN
EJ = ET(J)
SU = SU+SGNI*SGNJ*QN(J)/(EI-EJ)
    ELSE
SU = SU-.5DO*EI*QN(I)/(1.DO-EI*EI)
    ENDIF
SGNJ = -SGNJ
3  CONTINUE
DQN(I) = SU
SGNI = -SGNI
2  CONTINUE

RETURN
END
```

.

DELAGR

For any $n \geq 1$, the routine computes the first derivative of a polynomial q of degree at most $n - 1$ at the Laguerre Gauss-Radau nodes $\eta_j^{(n)}$, $0 \leq j \leq n - 1$. The polynomial q is individuated by the values attained at these nodes.

We have the formula

$$q'(\eta_i^{(n)}) = \sum_{j=0}^{n-1} \tilde{d}_{ij}^{(1)} q(\eta_j^{(n)}),$$

where

$$\tilde{d}_{ij}^{(1)} = \begin{cases} -(n+1)/(\alpha+2) & \text{if } i = j = 0, \\ \frac{(\alpha+1)L_n^{(\alpha)}(\eta_i^{(n)})}{\eta_i^{(n)} L_n^{(\alpha)}(\eta_0^{(n)})} & \text{if } 1 \leq i \leq n-1, j = 0, \\ \frac{-L_n^{(\alpha)}(\eta_0^{(n)})}{(\alpha+1)\eta_j^{(n)} L_n^{(\alpha)}(\eta_j^{(n)})} & \text{if } i = 0, 1 \leq j \leq n-1, \\ \frac{L_n^{(\alpha)}(\eta_i^{(n)})}{L_n^{(\alpha)}(\eta_j^{(n)})} \frac{1}{\eta_i^{(n)} - \eta_j^{(n)}} & \text{if } 1 \leq i \leq n-1, 1 \leq j \leq n-1, i \neq j, \\ (\eta_i^{(n)} - \alpha)/2\eta_i^{(n)} & \text{if } 1 \leq i = j \leq n-1. \end{cases}$$

The nodes can be determined by calling subroutine ZELAGR.

<i>Input variables</i>	<i>Output variable</i>
N , the number of nodes	DQN , the values of q' at the nodes
A , the parameter α	
ET , vector of the nodes $\eta_j^{(n)}$	
QN , the values of q at the nodes	

Auxiliary routine: **VALAPO**

```

      SUBROUTINE DELAGR(N,A,ET,QN,DQN)
      *****
      *   COMPUTES THE DERIVATIVE OF A POLYNOMIAL AT THE LAGUERRE GAUSS-RADAU
      *   NODES FROM THE VALUES OF THE POLYNOMIAL ATTAINED AT THE SAME POINTS
      *   N = THE NUMBER OF NODES
      *   A = PARAMETER >-1
      *   ET = VECTOR OF THE NODES, ET(I), I=0,N-1
      *   QN = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N-1
      *   DQN = DERIVATIVES OF THE POLYNOMIAL AT THE NODES, DQZ(I), I=0,N-1
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), QN(0:*), DQN(0:*)
      DN = DFLOAT(N)
      DQN(0) = (1.DO-DN)*QN(0)/(A+2.DO)
      IF (N .EQ. 1) RETURN

      A1 = A+1.DO
      SU = DQN(0)
      X = 0.DO
      CALL VALAPO(N,A,X,Y,DY,D2Y)
      C = Y
      DO 1 J=1,N-1
      EJ = ET(J)
      CALL VALAPO(N,A,EJ,Y,DY,D2Y)
      QN(J) = QN(J)/Y
      SU = SU-C*QN(J)/(A1+EJ)
1     CONTINUE
      DQN(0) = SU

      DO 2 I=1,N-1
      EI = ET(I)
      CALL VALAPO(N,A,EI,Y,DY,D2Y)
      SU = A1*QN(0)/(C*EI)
      DO 3 J=1,N-1
      IF (I .NE. J) THEN
      EJ = ET(J)
      SU = SU+QN(J)/(EI-EJ)
      ELSE
      SU = SU+.5D0*QN(I)*(EI-A)/EI
      ENDIF
3     CONTINUE
      DQN(I) = Y*SU
2     CONTINUE

      DO 4 J=1,N-1
      EJ = ET(J)
      CALL VALAPO(N,A,EJ,Y,DY,D2Y)
      QN(J) = Y*QN(J)
4     CONTINUE

      RETURN
      END

```

DMJAGL

For any $n \in \mathbf{N}$, the routine gives the entries of the $(n+1) \times (n+1)$ matrix $\tilde{d}_{ij}^{(1)}$, $0 \leq i \leq n$, $0 \leq j \leq n$, relative to the derivative operator in the space of polynomials of degree at most n , with respect to the Jacobi Gauss-Lobatto nodes. The expression of the entries is provided in the description of subroutine DEJAGL. The nodes and the values $P_n^{(\alpha,\beta)}(\eta_j^{(n)})$, $0 \leq j \leq n$, can be determined by calling subroutine ZEJAGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n NM , actual dimension of the matrix A , the parameter α B , the parameter β ET , vector of the nodes $\eta_j^{(n)}$ VN , the values $P_n^{(\alpha,\beta)}(\eta_j^{(n)})$	DMA , derivative matrix

SUBROUTINE DMJAGL(N,NM,A,B,ET,VN,DMA)

```

*****
* COMPUTES THE ENTRIES OF THE DERIVATIVE MATRIX RELATIVE TO THE
* IACOBI GAUSS-LOBATTO NODES
* N   = PARAMETER RELATIVE TO THE DIMENSION OF THE MATRIX
* NM  = ORDER OF THE MATRIX AS DECLARED IN THE MAIN DIMENSION STATEMENT
* A   = PARAMETER >-1
* B   = PARAMETER >-1
* ET  = VECTOR OF THE NODES, ET(I), I=0,N
* VN  = VALUES OF THE IACOBI POLYNOMIAL AT THE NODES, VN(I), I=0,N
* DMA = DERIVATIVE MATRIX, DMA(I,J), I=0,N J=0,N
*****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), VN(0:*), DMA(0:NM,0:*)
      DMA(0,0) = 0.DO
      IF (N .EQ. 0) RETURN

```



```
A1 = A+1.DO
B1 = B+1.DO
AB = A+B
DN = DFLOAT(N)
C1 = DN*(DN+AB+1.DO)
C2 = A1*VN(O)/(B1*VN(N))
DMA(O,O) = .5DO*(A-C1)/(B+2.DO)
DMA(N,N) = .5DO*(C1-B)/(A+2.DO)
DMA(O,N) = -.5DO*C2
DMA(N,O) = .5DO/C2
    IF (N .EQ. 1) RETURN

C3 = VN(O)/B1
C4 = VN(N)/A1
    DO 1 J=1,N-1
VJ = VN(J)
EJ = ET(J)
DMA(O,J) = -C3/(VJ*(1.DO+EJ))
DMA(N,J) = C4/(VJ*(1.DO-EJ))
DMA(J,O) = VJ/(C3*(1.DO+EJ))
DMA(J,N) = -VJ/(C4*(1.DO-EJ))
1    CONTINUE

    DO 2 I=1,N-1
VI = VN(I)
EI = ET(I)
    DO 3 J=1,N-1
    IF (I .NE. J) THEN
VJ = VN(J)
EJ = ET(J)
DMA(I,J) = VI/(VJ*(EI-EJ))
    ELSE
DMA(I,I) = .5DO*(AB*EI+A-B)/(1.DO-EI*EI)
    ENDIF
3    CONTINUE
2    CONTINUE

RETURN
END
```

DMLEGL

For any $n \in \mathbf{N}$, the routine gives the entries of the $(n+1) \times (n+1)$ matrix $\tilde{d}_{ij}^{(1)}$, $0 \leq i \leq n$, $0 \leq j \leq n$, relative to the derivative operator in the space of polynomials of degree at most n , with respect to the Legendre Gauss-Lobatto nodes. The expression of the entries is provided in the description of subroutine DELEGL. The nodes and the values $P_n(\eta_j^{(n)})$, $0 \leq j \leq n$, can be determined by calling subroutine ZELEGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n	DMA , derivative matrix
NM , actual dimension of the matrix	
ET , vector of the nodes $\eta_j^{(n)}$	
VN , the values $P_n(\eta_j^{(n)})$	

```

      SUBROUTINE DMLEGL(N,NM,ET,VN,DMA)
      *****
      * COMPUTES THE ENTRIES OF THE DERIVATIVE MATRIX RELATIVE TO THE
      * LEGENDRE GAUSS-LOBATTO NODES
      * N   = PARAMETER RELATIVE TO THE DIMENSION OF THE MATRIX
      * NM  = ORDER OF THE MATRIX AS DECLARED IN THE MAIN DIMENSION STATEMENT
      * ET  = VECTOR OF THE NODES, ET(I), I=0,N
      * VN  = VALUES OF THE LEGENDRE POLYNOMIAL AT THE NODES, VN(I), I=0,N
      * DMA = DERIVATIVE MATRIX, DMA(I,J), I=0,N J=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(0:*), VN(0:*), DMA(0:NM,0:*)
      DMA(0,0) = 0.DO
      IF (N .EQ. 0) RETURN

      DO 1 I=0,N
      VI = VN(I)
      EI = ET(I)
      DO 2 J=0,N
      IF (I .NE. J) THEN
      VJ = VN(J)
      EJ = ET(J)
      DMA(I,J) = VI/(VJ*(EI-EJ))
      ELSE
      DMA(I,I) = 0.DO
      ENDIF
2      CONTINUE
1      CONTINUE

      DN = DFLOAT(N)
      C = .25DO*DN*(DN+1.DO)
      DMA(0,0) = -C
      DMA(N,N) = C

      RETURN
      END

```

DMCHGL

For any $n \in \mathbf{N}$, the routine gives the entries of the $(n+1) \times (n+1)$ matrix $\tilde{d}_{ij}^{(1)}$, $0 \leq i \leq n$, $0 \leq j \leq n$, relative to the derivative operator in the space of polynomials of degree at most n , with respect to the Chebyshev Gauss-Lobatto nodes. The expression of the entries is provided in the description of subroutine DECHGL. The nodes can be determined by calling subroutine ZECHGL.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n	DMA , derivative matrix
NM , actual dimension of the matrix	
ET , vector of the nodes $\eta_j^{(n)}$	

```
SUBROUTINE DMCHGL(N,NM,ET,DMA)
```

```
*****
* COMPUTES THE ENTRIES OF THE DERIVATIVE MATRIX RELATIVE TO THE
* CHEBYSHEV GAUSS-LOBATTO NODES
* N   = PARAMETER RELATIVE TO THE DIMENSION OF THE MATRIX
* NM  = ORDER OF THE MATRIX AS DECLARED IN THE MAIN DIMENSION STATEMENT
* ET  = VECTOR OF THE NODES, ET(I), I=0,N
* DMA = DERIVATIVE MATRIX, DMA(I,J), I=0,N J=0,N
*****
  IMPLICIT DOUBLE PRECISION (A-H,O-Z)
  DIMENSION ET(0:*), DMA(0:NM,0:*)
  DMA(0,0) = 0.DO
  IF (N .EQ. 0) RETURN
```

```
DN = DFLOAT(N)
CN = (2.DO*DN*DN+1.DO)/6.DO
SN = DFLOAT(1+4*(N/2)-2*N)
DMA(O,O) = -CN
DMA(N,N) = CN
DMA(O,N) = -.5DO*SN
DMA(N,O) = .5DO*SN
    IF (N .EQ. 1) RETURN

SGN = -1.DO
    DO 1 J=1,N-1
EJ = ET(J)
DMA(O,J) = -2.DO*SGN/(1.DO+EJ)
DMA(N,J) = 2.DO*SGN*SN/(1.DO-EJ)
DMA(J,O) = .5DO*SGN/(1.DO+EJ)
DMA(J,N) = -.5DO*SGN*SN/(1.DO-EJ)
SGN = -SGN
1    CONTINUE

SGNI = -1.DO
    DO 2 I=1,N-1
EI = ET(I)
SGNJ = -1.DO
    DO 3 J=1,N-1
    IF (I .NE. J) THEN
EJ = ET(J)
DMA(I,J) = SGNI*SGNJ/(EI-EJ)
    ELSE
DMA(I,I) = -.5DO*EI/(1.DO-EI*EI)
    ENDIF
SGNJ = -SGNJ
3    CONTINUE
SGNI = -SGNI
2    CONTINUE

RETURN
END
```

DMLAGR

For any $n \geq 1$, the routine gives the entries of the $n \times n$ matrix $\tilde{d}_{ij}^{(1)}$, $0 \leq i \leq n-1$, $0 \leq j \leq n-1$, relative to the derivative operator in the space of polynomials of degree at most n , with respect to the Laguerre Gauss-Radau nodes. The expression of the entries is provided in the description of subroutine DELAGR. The nodes can be determined by calling subroutine ZELAGR.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n NM , actual dimension of the matrix A , the parameter α ET , vector of the nodes $\eta_j^{(n)}$	DMA , derivative matrix

Auxiliary routine: **VALAPO**

```

      SUBROUTINE DMLAGR(N,NM,A,ET,DMA)
      *****
      * COMPUTES THE ENTRIES OF THE DERIVATIVE MATRIX RELATIVE TO THE
      * LAGUERRE GAUSS-RADAU NODES
      * N   = DIMENSION OF THE MATRIX
      * NM  = ORDER OF THE MATRIX AS DECLARED IN THE MAIN DIMENSION STATEMENT
      * A   = PARAMETER >-1
      * ET  = VECTOR OF THE NODES, ET(I), I=0,N-1
      * DMA = DERIVATIVE MATRIX, DMA(I,J), I=0,N-1  J=0,N-1
      *****

```

```
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION ET(O:*), DMA(O:NM,O:*)
      DN = DFLOAT(N)
      DMA(O,O) = (1.DO-DN)/(A+2.DO)
      IF (N .EQ. 1) RETURN

      A1 = A+1.DO
      X = O.DO
      CALL VALAPO(N,A,X,Y,DY,D2Y)
      C = Y
      DO 1 J=1,N-1
      EJ = ET(J)
      CALL VALAPO(N,A,EJ,Y,DY,D2Y)
      DMA(O,J) = -C/(A1*EJ*Y)
      DMA(J,O) = A1*Y/(C*EJ)
1     CONTINUE

      DO 2 I=1,N-1
      EI = ET(I)
      CALL VALAPO(N,A,EI,Y,DY,D2Y)
      DO 3 J=1,N-1
      IF (I .NE. J) THEN
      EJ = ET(J)
      DMA(I,J) = Y/(EI-EJ)
      ELSE
      DMA(I,I) = .5DO*(EI-A)/EI
      ENDIF
3     CONTINUE
2     CONTINUE

      DO 4 J=1,N-1
      EJ = ET(J)
      CALL VALAPO(N,A,EJ,Y,DY,D2Y)
      DO 5 I=1,N-1
      IF (I .EQ. J) GOTO 5
      DMA(I,J) = DMA(I,J)/Y
5     CONTINUE
4     CONTINUE

      RETURN
      END
```

USING THE FAST FOURIER TRANSFORM

FCCHGA

For any $n \geq 1$, the routine computes the n Fourier coefficients c_k , $0 \leq k \leq n-1$, with respect to the Chebyshev polynomial basis, of a polynomial q of degree at most $n-1$ determined by the values attained at the zeroes $\xi_j^{(n)}$, $1 \leq j \leq n$, of the polynomial T_n . One has

$$c_0 = \frac{\gamma_0}{n}, \quad c_k = \frac{(-1)^k}{n} \left[\gamma_k e^{ik\pi/2n} + \gamma_{n-k} e^{-ik\pi/2n} \right] \quad 1 \leq k \leq n-1,$$

where $\mathbf{i} = \sqrt{-1}$ is the imaginary unity, and

$$\gamma_k = \sum_{j=0}^{n-1} \delta_j e^{2ikj\pi/n}, \quad 0 \leq k \leq n-1,$$

with

$$\delta_j = \begin{cases} q(\xi_{2j+1}^{(n)}) & 0 \leq j \leq n - [\frac{n}{2}] - 1, \\ q(\xi_{2n-2j}^{(n)}) & n - [\frac{n}{2}] \leq j \leq n-1. \end{cases}$$

The NAG FORTRAN Library routine C06EAF is used for the computation of the coefficients γ_k , $0 \leq k \leq n-1$, with the help of the Fast Fourier Transform algorithm.

<i>Input variables</i>	<i>Output variable</i>
N , the number of zeroes QZ , the values of q at the zeroes	CO , the Fourier coefficients of q

Auxiliary NAG Library routines are called by C06EAF.

```

      SUBROUTINE FCCHGA(N,QZ,CO)
      *****
      * COMPUTES USING FFT THE CHEBYSHEV FOURIER COEFFICIENTS OF A POLYNOMIAL
      * INDIVIDUATED BY THE VALUES ATTAINED AT THE CHEBYSHEV ZEROES
      * N   = THE NUMBER OF ZEROES
      * QZ  = VALUES OF THE POLYNOMIAL AT THE ZEROES, QZ(I), I=1,N
      * CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N-1
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION QZ(1), CO(0:*)
      IF(N .EQ. 0) RETURN

      PH = 1.57079632679489661923D0
      R2 = 1.41421356237309515D0
      DN = DFLOAT(N)
      CO(0) = QZ(1)/DN
      IF(N .EQ. 1) RETURN

      N2 = N/2
      C  = 2.DO/DSQRT(DN)
      SN = DFLOAT(1+4*N2-2*N)
      CO(N-N2-1) = QZ(N)
      DO 1 I=1,N2
        CO(I-1) = QZ(2*I-1)
        CO(N-I) = QZ(2*I)
1     CONTINUE

      CALL C06EAF(CO,N,IFAIL)
      IF (IFAIL .NE. 0) THEN
        WRITE(*,*) 'IFAIL IS NOT ZERO IN SUBROUTINE C06EAF'
      ENDIF

      CO(0) = .5D0*C*CO(0)
      IF (2*N2 .EQ. N) THEN
        CO(N2) = C*((-1.DO)**N2)*CO(N2)/R2
      ENDIF
      IF (N .EQ. 2) RETURN
      SM = -1.DO
      DO 2 M=1,N-N2-1
        AR = PH*DFLOAT(M)/DN
        CS = DCOS(AR)
        SI = DSIN(AR)
        V1 = C*SM*(CO(M)*CS+CO(N-M)*SI)
        V2 = C*SM*SN*(CO(M)*SI-CO(N-M)*CS)
        CO(M) = V1
        CO(N-M) = V2
        SM = -SM
2     CONTINUE

      RETURN
      END

```

FCCHGL

For any $n \geq 1$, the routine computes the $n + 1$ Fourier coefficients c_k , $0 \leq k \leq n$, with respect to the Chebyshev polynomial basis, of a polynomial q of degree at most n determined by the values attained at the Chebyshev Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$. One has

$$c_0 = \frac{1}{n} \left[\frac{1}{2} q(\eta_0^{(n)}) + \frac{1}{2} q(\eta_n^{(n)}) + \sum_{j=1}^{n-1} q(\eta_j^{(n)}) \right],$$

$$c_n = \frac{(-1)^n}{n} \left[\frac{1}{2} q(\eta_0^{(n)}) + \frac{(-1)^n}{2} q(\eta_n^{(n)}) + \sum_{j=1}^{n-1} (-1)^j q(\eta_j^{(n)}) \right],$$

$$c_k = \frac{(-1)^k}{2n} \left[\gamma_k \left(1 + \frac{1}{2 \sin(\pi k/n)} \right) + \bar{\gamma}_{n-k} \left(1 - \frac{1}{2 \sin(\pi k/n)} \right) \right], \quad 1 \leq k \leq n-1,$$

where the overbar denotes complex conjugate, and

$$\gamma_k = \sum_{j=0}^{n-1} \delta_j e^{2ikj\pi/n}, \quad 1 \leq k \leq n-1,$$

with

$$\delta_j = \begin{cases} q(\eta_0^{(n)}) & j = 0, \\ q(\eta_{2j}^{(n)}) + \mathbf{i}[q(\eta_{2j+1}^{(n)}) - q(\eta_{2j-1}^{(n)})] & 1 \leq j \leq n - [\frac{n}{2}] - 1, \\ q(\eta_n^{(n)}) & j = \frac{n}{2} \text{ if } n \text{ is even,} \\ q(\eta_{2n-2j}^{(n)}) + \mathbf{i}[q(\eta_{2n-2j-1}^{(n)}) - q(\eta_{2n-2j+1}^{(n)})] & [\frac{n}{2}] + 1 \leq j \leq n-1, \end{cases}$$

being $\mathbf{i} = \sqrt{-1}$ the imaginary unity.

The NAG FORTRAN Library routine C06EBF is used for the computation of the coefficients γ_k , $1 \leq k \leq n-1$, with the help of the Fast Fourier Transform algorithm.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n QZ , the values of q at the nodes	CO , the Fourier coefficients of q

Auxiliary NAG Library routines are called by C06EBF.

```

      SUBROUTINE FCCHGL(N,QN,CO)
      *****
      * COMPUTES USING FFT THE CHEBYSHEV FOURIER COEFFICIENTS OF A POLYNOMIAL
      * INDIVIDUATED BY ITS VALUES AT THE CHEBYSHEV GAUSS-LOBATTO NODES
      * N   = THE DEGREE OF THE POLYNOMIAL
      * QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      * CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION QN(0:*), CO(0:*)
      IF(N .EQ. 0) RETURN

      PI = 3.14159265358979323846D0
      DN = DFLOAT(N)
      N2 = N/2
      SN = DFLOAT(1+4*N2-2*N)
      S1 = .5D0*(QN(0)+QN(N))
      S2 = .5D0*(SN*QN(0)+QN(N))
      CO(0) = S1
      CO(N) = S2
      IF(N .EQ. 1) RETURN

      CO(0) = QN(0)
      IF (2*N2 .EQ. N) THEN
        CO(N2) = QN(N)
      ENDIF
      IF(N .EQ. 2) GOTO 2

      DO 1 I=1,N-N2-1
        I2 = 2*I
        CO(I) = QN(I2)
        CO(N-I) = QN(I2-1)-QN(I2+1)
1     CONTINUE

```

```

2  CALL C06EBF(CO,N,IFAIL)
   IF (IFAIL .NE. 0) THEN
   WRITE(*,*) 'IFAIL IS NOT ZERO IN SUBROUTINE C06EBF'
   ENDIF

      SJ = -1.DO
DO 3 J=1,N-1
      S1 = S1+QN(J)
      S2 = S2+SJ*SN*QN(J)
      SJ = -SJ
3  CONTINUE
      CO(0) = S1/DN
      CO(N) = S2/DN
      C = .5DO/DSQRT(DN)
   IF (2*N2 .EQ. N) THEN
      CO(N2) = 2.DO*C*((-1.DO)**N2)*CO(N2)
   ENDIF
   IF(N .EQ. 2) RETURN

      SM = -1.DO
DO 4 M=1,N-N2-1
      AR = PI*DFLOAT(M)/DN
      SI = .5DO/DSIN(AR)
      V1 = CO(M)*(1.DO+SI)+CO(N-M)*(1.DO-SI)
      V2 = CO(M)*(1.DO-SI)+CO(N-M)*(1.DO+SI)
      CO(M) = C*SM*V1
      CO(N-M) = C*SM*SN*V2
      SM = -SM
4  CONTINUE

RETURN
END

```


FVCHGL

For any $n \geq 1$, the routine computes the values of a polynomial q of degree at most n at the Chebyshev Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$, by knowing the $n + 1$ Fourier coefficients c_k , $0 \leq k \leq n$, of q with respect to the Chebyshev polynomial basis. One has

$$q(\eta_0^{(n)}) = \sum_{k=0}^n (-1)^k c_k, \quad q(\eta_n^{(n)}) = \sum_{k=0}^n c_k,$$

$$q(\eta_j^{(n)}) = \frac{1}{4} \left[\gamma_j \left(1 + \frac{1}{2 \sin(\pi j/n)} \right) + \bar{\gamma}_{n-j} \left(1 - \frac{1}{2 \sin(\pi j/n)} \right) \right], \quad 1 \leq j \leq n-1,$$

where the overbar denotes complex conjugate, and

$$\gamma_j = \sum_{k=0}^{n-1} \delta_k e^{2ikj\pi/n}, \quad 1 \leq j \leq n-1,$$

with

$$\delta_k = \begin{cases} 2c_0 & k = 0 \\ c_{2k} + \mathbf{i}(c_{2k-1} - c_{2k+1}) & 1 \leq k \leq \frac{n}{2} - 1 \\ 2c_n & k = \frac{n}{2} \\ c_{2n-2k} + \mathbf{i}(c_{2n-2k-1} - c_{2n-2k+1}) & \frac{n}{2} + 1 \leq k \leq n-1 \end{cases} \quad \text{for } n \text{ even;}$$

$$\delta_k = \begin{cases} 2c_0 & k = 0 \\ c_{2k} + \mathbf{i}(c_{2k-1} - c_{2k+1}) & 1 \leq k \leq \frac{n-3}{2} \\ c_{n-1} + \mathbf{i}(c_{n-2} - 2c_n) & k = \frac{n-1}{2} \\ c_{n-1} + \mathbf{i}(2c_n - c_{n-2}) & k = \frac{n+1}{2} \\ c_{2n-2k} + \mathbf{i}(c_{2n-2k-1} - c_{2n-2k+1}) & \frac{n+3}{2} \leq k \leq n-1 \end{cases} \quad \text{for } n \text{ odd;}$$

being $\mathbf{i} = \sqrt{-1}$ the imaginary unity.

The NAG FORTRAN Library routine C06EBF is used for the computation of the coefficients γ_j , $1 \leq j \leq n-1$, with the help of the Fast Fourier Transform algorithm.

<i>Input variables</i>	<i>Output variable</i>
N , the degree n CO , the Fourier coefficients of q	QN , the values of q at the nodes

Auxiliary NAG Library routines are called by C06EBF.

```

      SUBROUTINE FVCHGL(N,CO,QN)
      *****
      * COMPUTES USING FFT THE VALUES OF A POLYNOMIAL AT THE CHEBYSHEV
      * GAUSS-LOBATTO NODES FROM ITS CHEBYSHEV FOURIER COEFFICIENTS
      * N   = THE DEGREE OF THE POLYNOMIAL
      * CO  = FOURIER COEFFICIENTS OF THE POLYNOMIAL, CO(I), I=0,N
      * QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION QN(0:*), CO(0:*)
      IF(N .EQ. 0) RETURN

      PI = 3.14159265358979323846D0
      DN = DFLOAT(N)
      N2 = N/2
      SN = DFLOAT(1+4*N2-2*N)
      S1 = CO(0)+SN*CO(N)
      S2 = CO(0)+CO(N)
      QN(0) = S1
      QN(N) = S2
      IF(N .EQ. 1) RETURN

      QN(0) = 2.DO*CO(0)
      IF (2*N2 .EQ. N) THEN
        QN(N2) = 2.DO*CO(N)
      ENDIF
      IF(N .EQ. 2) GOTO 2

```

```

DO 1 I=1,N-N2-1
  I2 = 2*I
  QN(N-I) = CO(I2+1)-CO(I2-1)
  QN(I) = CO(I2)
1 CONTINUE
IF (2*N2 .NE. N) THEN
  QN(N2+1) = 2.DO*CO(N)-CO(N-2)
ENDIF

2 CALL CO6EBF(QN,N,IFAIL)
IF (IFAIL .NE. 0) THEN
WRITE(*,*) 'IFAIL IS NOT ZERO IN SUBROUTINE CO6EBF'
ENDIF

  SJ = -1.DO
DO 3 J=1,N-1
  S1 = S1+SJ*CO(J)
  S2 = S2+CO(J)
  SJ = -SJ
3 CONTINUE
  QN(0) = S1
  QN(N) = S2
  C = .25DO*DSQRT(DN)
IF (2*N2 .EQ. N) THEN
  QN(N2) = 2.DO*C*QN(N2)
ENDIF
IF(N .EQ. 2) RETURN

DO 4 M=1,N-N2-1
  AR = PI*DFLOAT(M)/DN
  SI = .5DO/DSIN(AR)
  V1 = QN(M)*(1.DO+SI)+QN(N-M)*(1.DO-SI)
  V2 = QN(M)*(1.DO-SI)+QN(N-M)*(1.DO+SI)
  QN(M) = C*V1
  QN(N-M) = C*V2
4 CONTINUE

RETURN
END

```


FDCHGL

For any $n \geq 1$, the routine computes the Fourier coefficients $c_k^{(1)}$, $0 \leq k \leq n$, and the values at the Chebyshev Gauss-Lobatto nodes $\eta_j^{(n)}$, $0 \leq j \leq n$, of the derivative of a polynomial q of degree at most n , by knowing the values of q at the nodes.

The computation is performed by first computing the Fourier coefficients c_k , $0 \leq k \leq n$, of q using FCCHGL. Then, with the recursion formula given in the description of subroutine COCHDE, one obtains the coefficients $c_k^{(1)}$, $0 \leq k \leq n$. Finally, the quantities $q'(\eta_j^{(n)})$, $0 \leq j \leq n$, are evaluated using subroutine FVCHGL.

Note that the NAG FORTRAN Library routine C06EBF is called both in subroutines FCCHGL and FVCHGL.

<i>Input variables</i>	<i>Output variables</i>
N , the degree n	CD , the Fourier coefficients of q'
QN , the values of q at the nodes	DQN , the values of q' at the nodes

Auxiliary routines: **FCCHGL**, **FVCHGL**

```

      SUBROUTINE FDCHGL(N,QN,CD,DQN)
      *****
      * COMPUTES USING FFT THE FOURIER COEFFICIENTS AND THE VALUES AT THE
      * CHEBYSHEV GAUSS-LOBATTO NODES OF THE DERIVATIVE OF A POLYNOMIAL
      * FROM THE VALUES ATTAINED BY THE POLYNOMIAL AT THE NODES
      * N   = THE DEGREE OF THE POLYNOMIAL
      * QN  = VALUES OF THE POLYNOMIAL AT THE NODES, QN(I), I=0,N
      * CD  = FOURIER COEFFICIENTS OF THE DERIVATIVE, CD(I), I=0,N
      * DQN = VALUES OF THE DERIVATIVE AT THE NODES, DQN(I), I=0,N
      *****
      IMPLICIT DOUBLE PRECISION (A-H,O-Z)
      DIMENSION QN(0:*), CD(0:*), DQN(0:*)
      IF(N .EQ. 0) RETURN

      CALL FCCHGL(N,QN,DQN)
      CD(N) = 0.DO
      CD(0) = DQN(1)
      IF(N .EQ. 1) GOTO 2

      DN = DFLOAT(N)
      CD(N-1) = 2.DO*DN*DQN(N)
      DO 1 K=0,N-2
      KR = N-K-2
      IF(KR .NE. 0) THEN
      DK = 2.DO*(DFLOAT(KR)+1.DO)
      CD(KR) = CD(KR+2)+DK*DQN(KR+1)
      ELSE
      CD(0) = .5DO*CD(2)+DQN(1)
      ENDIF
1      CONTINUE

2      CALL FVCHGL(N,CD,DQN)

      RETURN
      END

```

AN EXAMPLE

For any $n \geq 1$, we would like to compute the approximated solution to the boundary-value problem

$$\begin{cases} -U'' + \tau U = f & \text{in }]-1, 1[, \quad \tau > 0, \\ U(-1) = \sigma_1, \\ U(1) = \sigma_2, \end{cases}$$

by the collocation method at the Legendre Gauss-Lobatto points.

Thus, the approximating polynomial s_n of degree n satisfies the set of equations

$$\begin{cases} -s_n''(\eta_i^{(n)}) + \tau s_n(\eta_i^{(n)}) = f(\eta_i^{(n)}) & 1 \leq i \leq n-1, \\ s_n(\eta_0^{(n)}) = \sigma_1, \\ s_n(\eta_n^{(n)}) = \sigma_2. \end{cases}$$

In order to show how some of the routines described in this book can be called by a main program, we provide a simple example. The program SAMPLE evaluates the approximating polynomial s_n for $2 \leq n \leq 10$, when for instance $\tau = 1$, $\sigma_1 = \sigma_2 = 0$, and $f(x) \equiv 1$, $x \in]-1, 1[$.

```

PROGRAM SAMPLE
IMPLICIT DOUBLE PRECISION (A-H,O-Z)
DIMENSION ET(0:10), VN(0:10), WT(0:10)
DIMENSION SO(0:10), FU(0:10), DMA(0:10,0:10)
DIMENSION DMA2(0:10,0:10), QN(0:10), WKSPCE(0:10)

C SET THE PARAMETERS
SIGMA1 = 0.DO
SIGMA2 = 0.DO
TAU = 1.DO
EF = 1.DO
EXP = 2.71828182845904509D0
COST = EXP/(1.DO+EXP*EXP)

DO 1 N=2,10

C COMPUTATION OF THE NODES, WEIGHTS AND DERIVATIVE MATRIX
CALL ZELEGL(N,ET,VN)
CALL WELEGL(N,ET,VN,WT)
CALL DMLEGL(N,10,ET,VN,DMA)

```



```
C CONSTRUCTION OF THE MATRIX CORRESPONDING TO THE
C DIFFERENTIAL OPERATOR
  DO 2 I=0,N
  DO 2 J=0,N
  SUM = 0.DO
  DO 3 K=0,N
  SUM = SUM + DMA(I,K)*DMA(K,J)
3  CONTINUE
  OPER = -SUM
  IF(I.EQ.J) OPER = -SUM + TAU
  DMA2(I,J) = OPER
2  CONTINUE

C CHANGE OF THE ENTRIES OF THE MATRIX ACCORDING TO THE
C BOUNDARY CONDITIONS
  DO 4 J=0,N
  DMA2(0,J) = 0.DO
  DMA2(N,J) = 0.DO
4  CONTINUE
  DMA2(0,0) = 1.DO
  DMA2(N,N) = 1.DO

C CONSTRUCTION OF THE RIGHT-HAND SIDE VECTOR
  DO 5 I=1,N-1
  FU(I) = EF
5  CONTINUE
  FU(0) = SIGMA1
  FU(N) = SIGMA2

C SOLUTION OF THE LINEAR SYSTEM
  N1 = N + 1
  CALL F04AAF(DMA2,11,FU,11,N1,1,SO,11,WKSPCE,IFAIL)

C COMPARISON WITH THE EXACT SOLUTION
  DO 6 I=0,N
  EI = DEXP(ET(I))
  QN(I) = SO(I)-1.DO+COST*(EI+1.DO/EI)
6  CONTINUE

C NORM OF THE ERROR VECTOR
  CALL NOLEGL(N,VN,QN,WT,QI,QS,QM)
  WRITE(*,*) N, QI

1  CONTINUE
  RETURN
  END
```

The program allows the computation of the error

$$E_n = \left(\int_{-1}^1 (s_n - I_n U)(x) dx \right)^{\frac{1}{2}},$$

where $I_n U$ is the n -degree interpolant polynomial at the nodes, of the exact solution given by

$$U(x) = 1 - \frac{e}{1 + e^2}(e^x + e^{-x}).$$

The output results are shown in the following table.

n	E_n
2	0.192227968354237912E - 01
3	0.437255201408398857E - 03
4	0.544146318231265948E - 04
5	0.585754215000575047E - 06
6	0.142928575377323905E - 06
7	0.908588864973060870E - 09
8	0.278096559701610084E - 09
9	0.117001702046232747E - 11
10	0.403084195089551988E - 12

To solve the system of $n + 1$ linear equations we used the NAG routine F04AAF.

TIMING

In this section, we provide an asymptotic estimate, with respect to the parameter n , of the CPU time needed to run each routine. The results, measured in microseconds, are only qualitative since they strongly depend on the computer architecture. The times of the routines using the Fast Fourier Transform are related to the case in which n is a power of two.

VAJAPO	$104 n$
VALEPO	$66 n$
VACHPO	$39 n$
VALAPO	$68 n$
VAHEPO	$24 n$
VALASF	$94 n$
VAHESF	$50 n$
ZEJAGA	$790 n^2$
ZELEGA	$260 n^2$
ZECHGA	$83 n$
ZELAGA	$580 n^2$
ZEHEGA	$150 n^2$

ZEJAGL	$740 n^2$
ZELEGL	$250 n^2$
ZECHGL	$22 n$
ZELAGR	$630 n^2$
WEJAGA	$70 n$
WELEGA	$15 n$
WECHGA	$6 n$
WELAGA	$75 n^2$
WEHEGA	$47 n^2$
WEJAGL	$200 n^2$
WELEGL	$15 n$
WECHGL	$8 n$

WELAGR	$130 n^2$
WECHCC	$57 n^2$
INJAGA	$130 n$
INLEGA	$95 n$
INCHGA	$99 n$
INLAGA	$22 n^2$
INHEGA	$12 n^2$
INJAGL	$140 n$
INLEGL	$100 n$
INCHGL	$130 n$
INLAGR	$23 n^2$
NOLEGA	$13 n$
NOCHGA	$100 n^2$
NOJAGL	$100 n$
NOLEGL	$30 n$
NOCHGL	$130 n^2$
COJAGA	$71 n^2$
COLEGA	$33 n^2$
COCHGA	$48 n^2$
COLAGA	$40 n^2$

COHEGA	$26 n^2$
COJAGL	$80 n^2$
COLEGL	$35 n^2$
COCHGL	$50 n^2$
COLAGR	$40 n^2$
PVJAEX	$120 n$
PVLEEX	$82 n$
PVCHEX	$53 n$
PVLAEX	$86 n$
PVHEEX	$52 n$
NOJAEX	$50 n$
NOLEEX	$13 n$
NOCHEX	$20 n^2$
NOLAEX	$36 n$
NOHEEX	$57 n$
COJADE	$73 n$
COLEDE	$46 n$
COCHDE	$40 n$
COLADE	$27 n$
COHEDE	$27 n$

DEJAGA	$23 n^2$
DELAGA	$220 n^2$
DEHEGA	$92 n^2$
DEJAGL	$25 n^2$
DELEGL	$23 n^2$
DECHGL	$27 n^2$
DELAGR	$220 n^2$
DMJAGL	$29 n^2$

DMLEGL	$30 n^2$
DMCHGL	$30 n^2$
DMLAGR	$230 n^2$
FCCHGA	$13 n \log_2 n$
FCCHGL	$13 n \log_2 n$
FVCHGL	$13 n \log_2 n$
FDCHGL	$30 n \log_2 n$

DIMENSIONING

Zeroes of the orthogonal polynomials of degree n	$\xi_i^{(n)} \quad 1 \leq i \leq n$	CS(I), I=1,N
Derivatives of the Jacobi polynomial of degree n at the zeroes	$\frac{d}{dx} P_n^{(\alpha,\beta)}(\xi_i^{(n)}) \quad 1 \leq i \leq n$	DZ(I), I=1,N
Derivatives of the Legendre polynomial of degree n at the zeroes	$P'_n(\xi_i^{(n)}) \quad 1 \leq i \leq n$	DZ(I), I=1,N
Derivatives of the Chebyshev polynomial of degree n at the zeroes	$T'_n(\xi_i^{(n)}) \quad 1 \leq i \leq n$	DZ(I), I=1,N
Derivatives of the scaled Laguerre function of degree n at the zeroes	$\frac{d}{dx} \hat{L}_n^{(\alpha)}(\xi_i^{(n)}) \quad 1 \leq i \leq n$	DZ(I), I=1,N
Derivatives of the scaled Hermite function of degree n at the zeroes	$\hat{H}'_n(\xi_i^{(n)}) \quad 1 \leq i \leq n$	DZ(I), I=1,N
Nodes of the Gauss-Lobatto formula of order n	$\eta_i^{(n)} \quad 0 \leq i \leq n$	ET(I), I=0,N
Nodes of the Laguerre Gauss-Radau formula of order n	$\eta_i^{(n)} \quad 0 \leq i \leq n-1$	ET(I), I=0,N-1
Values of the Jacobi polynomial of degree n at the nodes	$P_n^{(\alpha,\beta)}(\eta_i^{(n)}) \quad 0 \leq i \leq n$	VN(I), I=0,N
Values of the Legendre polynomial of degree n at the nodes	$P_n(\eta_i^{(n)}) \quad 0 \leq i \leq n$	VN(I), I=0,N
Values of the scaled Laguerre function of degree n at the nodes	$\hat{L}_n^{(\alpha)}(\eta_i^{(n)}) \quad 0 \leq i \leq n-1$	VN(I), I=0,N-1
Weights of the Gauss formula of order n	$w_i^{(n)} \quad 1 \leq i \leq n$	WE(I), I=1,N
Weights of the Gauss-Lobatto formula of order n	$\tilde{w}_i^{(n)} \quad 0 \leq i \leq n$	WT(I), I=0,N
Weights of the Laguerre Gauss-Radau formula of order n	$\tilde{w}_i^{(n)} \quad 0 \leq i \leq n-1$	WT(I), I=0,N-1
Weights of the Clenshaw-Curtis formula of order $2n$	$\chi_i^{(2n)} \quad 0 \leq i \leq 2n$	WK(I), I=0,2*N

Values of a polynomial of degree $n - 1$ at the zeroes	$q(\xi_i^{(n)}) \quad 1 \leq i \leq n$	$\text{QZ}(\mathbf{I}), \quad \mathbf{I}=1, \mathbf{N}$
Values of a polynomial of degree n at the Gauss-Lobatto nodes	$q(\eta_i^{(n)}) \quad 0 \leq i \leq n$	$\text{QN}(\mathbf{I}), \quad \mathbf{I}=0, \mathbf{N}$
Values of a polynomial of degree $n - 1$ at the Laguerre Gauss-Radau nodes	$q(\eta_i^{(n)}) \quad 0 \leq i \leq n - 1$	$\text{QN}(\mathbf{I}), \quad \mathbf{I}=0, \mathbf{N}-1$
Fourier coefficients of a polynomial of degree n	$c_k \quad 0 \leq k \leq n$	$\text{CO}(\mathbf{K}), \quad \mathbf{K}=0, \mathbf{N}$
Fourier coefficients of the derivative of a polynomial of degree n	$c_k^{(1)} \quad 0 \leq k \leq n$	$\text{CD}(\mathbf{K}), \quad \mathbf{K}=0, \mathbf{N}$
Fourier coefficients of the second derivative of a polynomial of degree n	$c_k^{(2)} \quad 0 \leq k \leq n$	$\text{CD2}(\mathbf{K}), \quad \mathbf{K}=0, \mathbf{N}$
Derivative of a polynomial of degree $n - 1$ at the zeroes	$q'(\xi_i^{(n)}) \quad 1 \leq i \leq n$	$\text{DQZ}(\mathbf{I}), \quad \mathbf{I}=1, \mathbf{N}$
Derivative of a polynomial of degree n at the Gauss-Lobatto nodes	$q'(\eta_i^{(n)}) \quad 0 \leq i \leq n$	$\text{DQN}(\mathbf{I}), \quad \mathbf{I}=0, \mathbf{N}$
Derivative of a polynomial of degree $n - 1$ at the Laguerre Gauss-Radau nodes	$q'(\eta_i^{(n)}) \quad 0 \leq i \leq n - 1$	$\text{DQN}(\mathbf{I}), \quad \mathbf{I}=0, \mathbf{N}-1$
Derivative matrix relative to the Gauss-Lobatto nodes	$d_{ij}^{(1)} \quad \begin{matrix} 0 \leq i \leq n \\ 0 \leq j \leq n \end{matrix}$	$\text{DMA}(\mathbf{I}, \mathbf{J}), \quad \begin{matrix} \mathbf{I} = 0, \mathbf{N} \\ \mathbf{J} = 0, \mathbf{N} \end{matrix}$
Derivative matrix relative to the Laguerre Gauss-Radau nodes	$d_{ij}^{(1)} \quad \begin{matrix} 0 \leq i \leq n - 1 \\ 0 \leq j \leq n - 1 \end{matrix}$	$\text{DMA}(\mathbf{I}, \mathbf{J}), \quad \begin{matrix} \mathbf{I} = 0, \mathbf{N} - 1 \\ \mathbf{J} = 0, \mathbf{N} - 1 \end{matrix}$

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